

Binary Search Trees



EECS2011 X:
Fundamentals of Data Structures
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Learning Outcomes of this Lecture

This module is designed to help you understand:

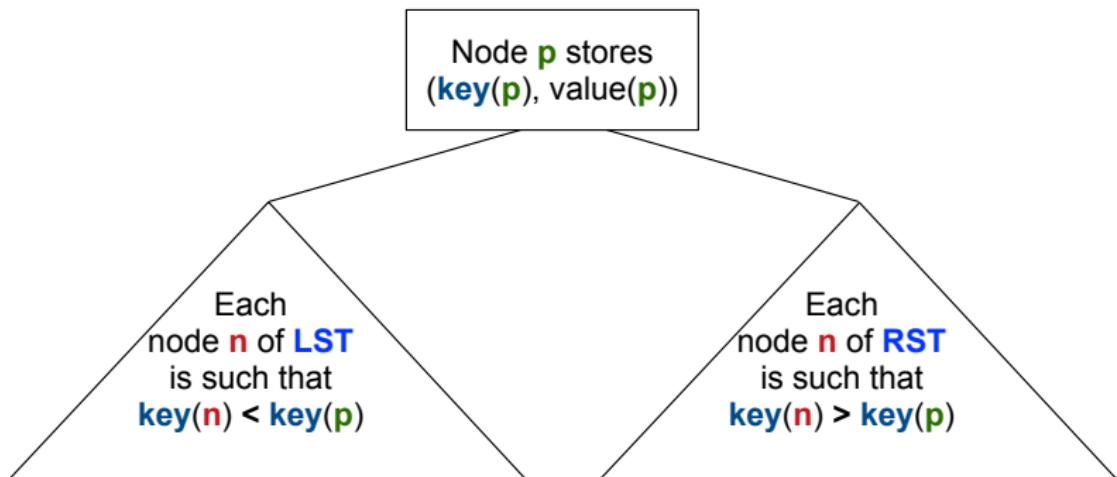
- ***Binary Search Trees (BSTs)*** = ***BTs*** + ***Search Property***
- Implementing a ***Generic*** BST in Java
- BST Operations:
 - ***Searching***: Implementation, Visualization, RT
 - ***Insertion***: (Sketch of) Implementation, Visualization, RT
 - ***Deletion***: (Sketch of) Implementation, Visualization, RT

Binary Search Tree: Recursive Definition

A **Binary Search Tree** (*BST*) is a **BT** satisfying the **search property**:

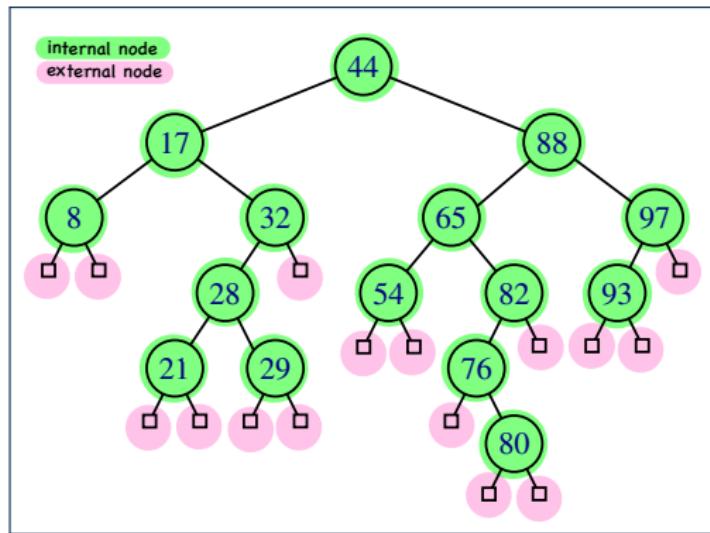
Each **internal node** p stores an **entry**, a key-value pair (k, v) , such that:

- For each node n in the **LST** of p : $\text{key}(n) < \text{key}(p)$
- For each node n in the **RST** of p : $\text{key}(n) > \text{key}(p)$



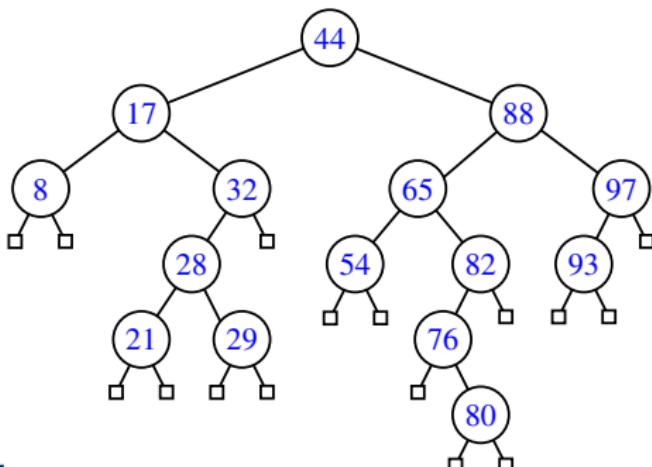
BST: Internal Nodes vs. External Nodes

- We store key-value pairs only in **internal nodes**.
- Recall how we treat **header** and **trailer** in a DLL.
- We treat **external nodes** as **sentinels**, in order to simplify the **coding logic** of BST algorithms.



BST: Sorting Property

- An ***in-order traversal*** of a **BST** will result in a sequence of nodes whose **keys** are arranged in an **ascending order**.
- Unless necessary, we may only show **keys** in BST nodes:



Justification:

- **In-Order Traversal**: Visit **LST**, then **root**, then **RST**.
- **Search Property of BST**: keys in **LST/RST** **</>** **root's key**

Implementation: Generic BST Nodes

```
public class BSTNode<E> {
    private int key; /* key */
    private E value; /* value */
    private BSTNode<E> parent; /* unique parent node */
    private BSTNode<E> left; /* left child node */
    private BSTNode<E> right; /* right child node */

    public BSTNode() { ... }
    public BSTNode(int key, E value) { ... }

    public boolean isExternal() {
        return this.getLeft() == null && this.getRight() == null;
    }
    public boolean isInternal() {
        return !this.isExternal();
    }
    public int getKey() { ... }
    public void setKey(int key) { ... }
    public E getValue() { ... }
    public void setValue(E value) { ... }
    public BSTNode<E> getParent() { ... }
    public void setParent(BSTNode<E> parent) { ... }
    public BSTNode<E> getLeft() { ... }
    public void setLeft(BSTNode<E> left) { ... }
    public BSTNode<E> getRight() { ... }
    public void setRight(BSTNode<E> right) { ... }
}
```

Implementation: BST Utilities – Traversal

```
import java.util.ArrayList;
public class BSTUtilities<E> {
    public ArrayList<BSTNode<E>> inOrderTraversal(BSTNode<E> root) {
        ArrayList<BSTNode<E>> result = null;
        if(root.isInternal()) {
            result = new ArrayList<>();

            if(root.getLeft().isInternal) {
                result.addAll(inOrderTraversal(root.getLeft()));
            }

            result.add(root);

            if(root.getRight().isInternal) {
                result.addAll(inOrderTraversal(root.getRight()));
            }
        }
        return result;
    }
}
```

Testing: Connected BST Nodes

Constructing a **BST** is similar to constructing a **General Tree**:

```
@Test
public void test_binary_search_trees_construction() {
    BSTNode<String> n28 = new BSTNode<>(28, "alan");
    BSTNode<String> n21 = new BSTNode<>(21, "mark");
    BSTNode<String> n35 = new BSTNode<>(35, "tom");
    BSTNode<String> extN1 = new BSTNode<>();
    BSTNode<String> extN2 = new BSTNode<>();
    BSTNode<String> extN3 = new BSTNode<>();
    BSTNode<String> extN4 = new BSTNode<>();

    n28.setLeft(n21); n21.setParent(n28);
    n28.setRight(n35); n35.setParent(n28);
    n21.setLeft(extN1); extN1.setParent(n21);
    n21.setRight(extN2); extN2.setParent(n21);
    n35.setLeft(extN3); extN3.setParent(n35);
    n35.setRight(extN4); extN4.setParent(n35);
    BSTUtilities<String> u = new BSTUtilities<>();
    ArrayList<BSTNode<String>> inOrderList = u.inOrderTraversal(n28);
    assertTrue(inOrderList.size() == 3);
    assertEquals(21, inOrderList.get(0).getKey());
    assertEquals("mark", inOrderList.get(0).getValue());
    assertEquals(28, inOrderList.get(1).getKey());
    assertEquals("alan", inOrderList.get(1).getValue());
    assertEquals(35, inOrderList.get(2).getKey());
    assertEquals("tom", inOrderList.get(2).getValue());
}
```

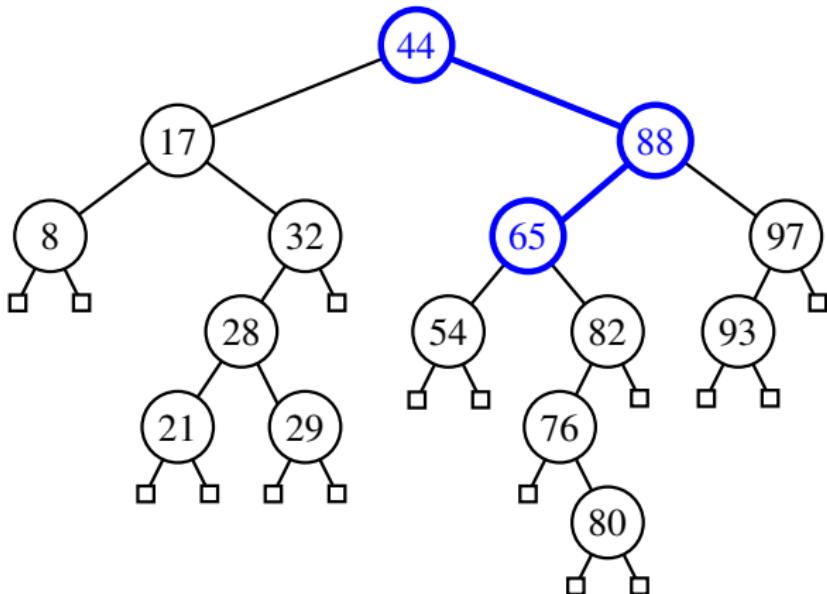
Implementing BST Operation: Searching

Given a **BST** rooted at node **p**, to locate a particular **node** whose **key** matches **k**, we may view it as a **decision tree**.

```
public BSTNode<E> search(BSTNode<E> p, int k) {  
    BSTNode<E> result = null;  
    if(p.isExternal()) {  
        result = p; /* unsuccessful search */  
    }  
    else if(p.getKey() == k) {  
        result = p; /* successful search */  
    }  
    else if(k < p.getKey()) {  
        result = search(p.getLeft(), k); /* recur on LST */  
    }  
    else if(k > p.getKey()) {  
        result = search(p.getRight(), k); /* recur on RST */  
    }  
    return result;  
}
```

Visualizing BST Operation: Searching (1)

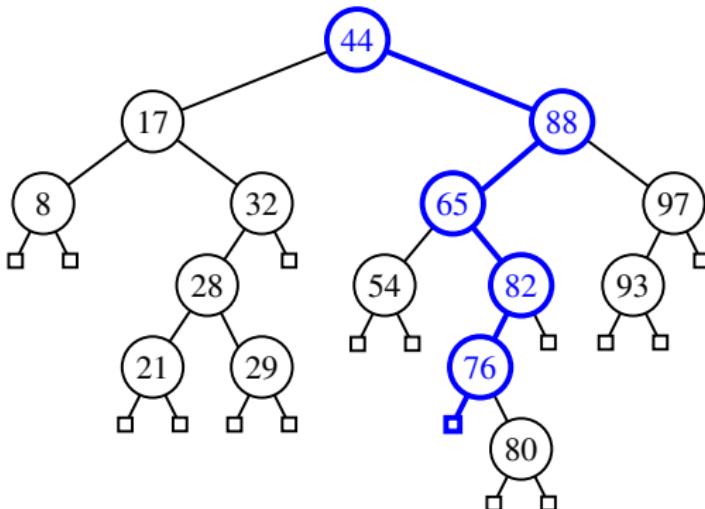
A **successful** search for **key 65**:



The **internal node** storing key 65 is returned.

Visualizing BST Operation: Searching (2)

- An *unsuccessful* search for **key 68**:



The **external, left child node** of the **internal node** storing **key 76** is returned.

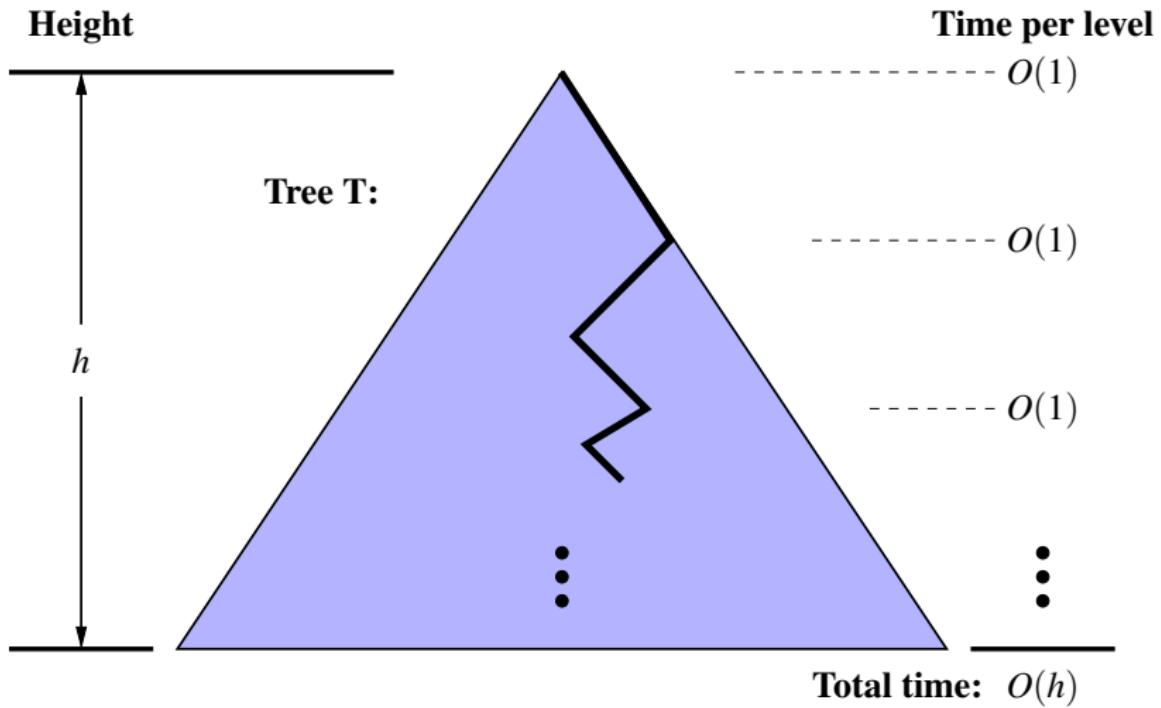
- Exercise :** Provide **keys** for different **external nodes** to be returned.

Testing BST Operation: Searching

```
@Test
public void test_binary_search_trees_search() {
    BSTNode<String> n28 = new BSTNode<>(28, "alan");
    BSTNode<String> n21 = new BSTNode<>(21, "mark");
    BSTNode<String> n35 = new BSTNode<>(35, "tom");
    BSTNode<String> extN1 = new BSTNode<>();
    BSTNode<String> extN2 = new BSTNode<>();
    BSTNode<String> extN3 = new BSTNode<>();
    BSTNode<String> extN4 = new BSTNode<>();
    n28.setLeft(n21); n21.setParent(n28);
    n28.setRight(n35); n35.setParent(n28);
    n21.setLeft(extN1); extN1.setParent(n21);
    n21.setRight(extN2); extN2.setParent(n21);
    n35.setLeft(extN3); extN3.setParent(n35);
    n35.setRight(extN4); extN4.setParent(n35);

    BSTUtilities<String> u = new BSTUtilities<>();
    /* search existing keys */
    assertTrue(n28 == u.search(n28, 28));
    assertTrue(n21 == u.search(n28, 21));
    assertTrue(n35 == u.search(n28, 35));
    /* search non-existing keys */
    assertTrue(extN1 == u.search(n28, 17)); /* *17* < 21 */
    assertTrue(extN2 == u.search(n28, 23)); /* 21 < *23* < 28 */
    assertTrue(extN3 == u.search(n28, 33)); /* 28 < *33* < 35 */
    assertTrue(extN4 == u.search(n28, 38)); /* 35 < *38* */
}
```

RT of BST Operation: Searching (1)



RT of BST Operation: Searching (2)

- Recursive calls of search are made on a **path** which
 - Starts from the **root**
 - Goes down one **level** at a time
 - RT of deciding from each node to go to LST or RST? [**O(1)**]
 - Stops when the key is found or when a **leaf** is reached
 - Maximum** number of nodes visited by the search? [**h + 1**]
 - ∴ RT of **search on a BST** is **O(h)**
- Recall: Given a BT with **n** nodes, the **height h** is bounded as:

$$\log(n+1) - 1 \leq h \leq n - 1$$
 - **Best** RT of a **binary search** is **O(log(n))** [**balanced** BST]
 - **Worst** RT of a **binary search** is **O(n)** [**ill-balanced** BST]
- **Binary search** on non-linear vs. linear structures:

	Search on a BST	Binary Search on a Sorted Array
START	Root of BST	Middle of Array
PROGRESS	LST or RST	Left Half or Right Half of Array
BEST RT	O(log(n))	
WORST RT	O(n)	O(log(n))

Sketch of BST Operation: Insertion

To **insert** an **entry** (with **key** k & **value** v) into a BST rooted at **node** n :

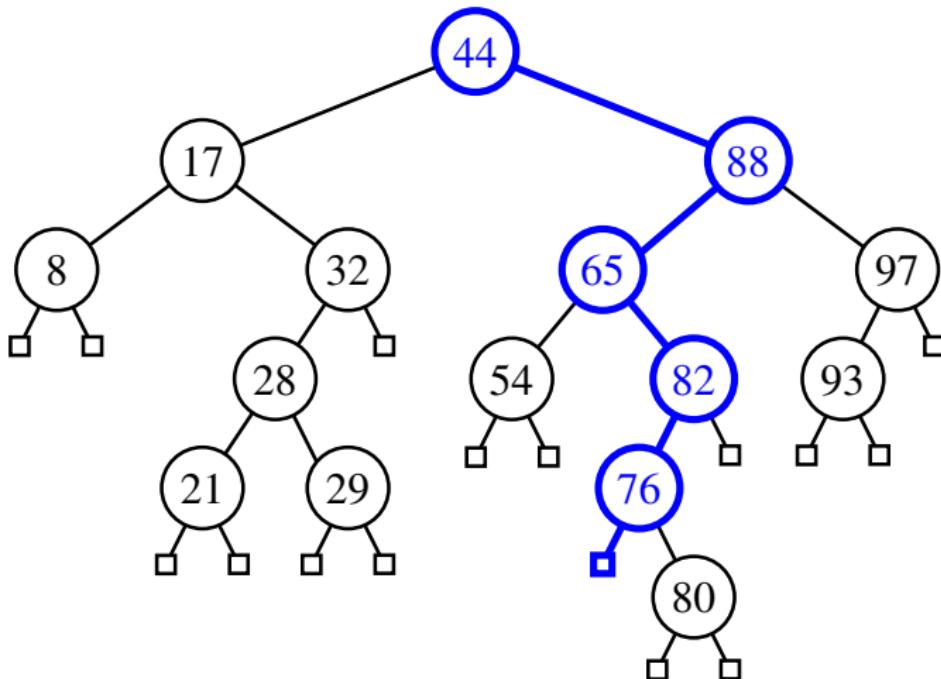
- Let node p be the return value from `search(n, k)`.
- If p is an **internal node**
 - ⇒ Key k exists in the BST.
 - ⇒ Set p 's value to v .
- If p is an **external node**
 - ⇒ Key k does not exist in the BST.
 - ⇒ Set p 's key and value to k and v .

Running time?

[$O(h)$]

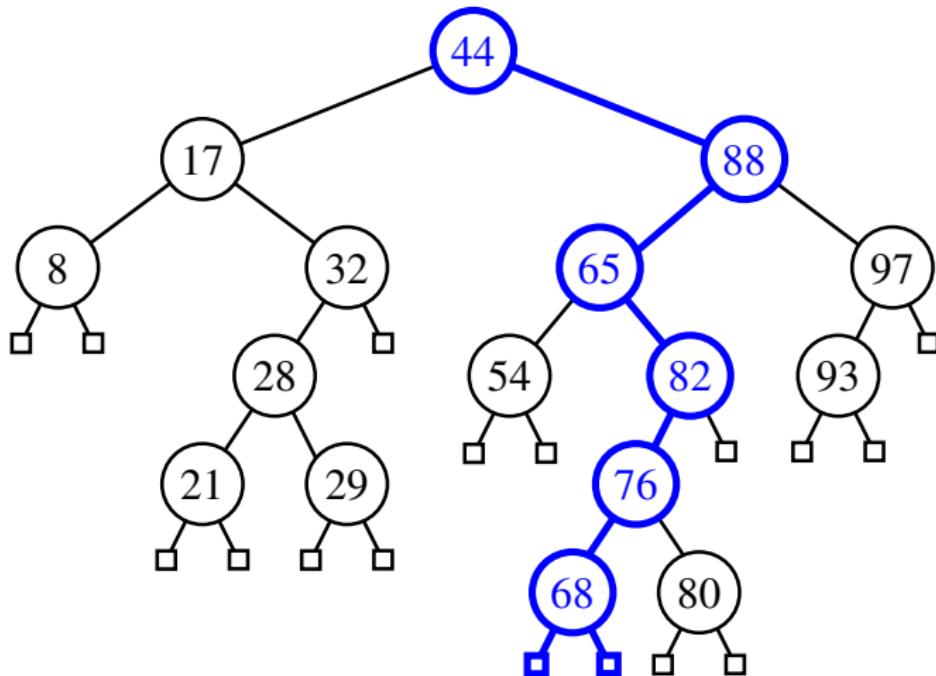
Visualizing BST Operation: Insertion (1)

Before *inserting* an entry with **key 68** into the following BST:



Visualizing BST Operation: Insertion (2)

After *inserting* an entry with **key 68** into the following BST:



Exercise on BST Operation: Insertion

Exercise: In BSTUtilities class, **implement** and **test** the
`void insert(BSTNode<E> p, int k, E v)` method.

Sketch of BST Operation: Deletion

To **delete** an **entry** (with key k) from a BST rooted at **node n** :

Let node p be the return value from $\text{search}(n, k)$.

- **Case 1:** Node p is **external**.

k is not an existing key \Rightarrow Nothing to remove

- **Case 2:** Both of node p 's child nodes are **external**.

No “orphan” subtrees to be handled \Rightarrow Remove p

[Still BST?]

- **Case 3:** One of the node p 's children, say r , is **internal**.

- r 's sibling is **external** \Rightarrow Replace node p by node r

[Still BST?]

- **Case 4:** Both of node p 's children are **internal**.

- Let r be the **right-most internal node** p 's **LST**.

$\Rightarrow r$ contains the **largest key s.t. $\text{key}(r) < \text{key}(p)$** .

Exercise: Can r contain the **smallest key s.t. $\text{key}(r) > \text{key}(p)$** ?

- Overwrite node p 's entry by node r 's entry.

[Still BST?]

- r being the **right-most internal node** may have:

- Two **external child nodes** \Rightarrow Remove r as in **Case 2**.

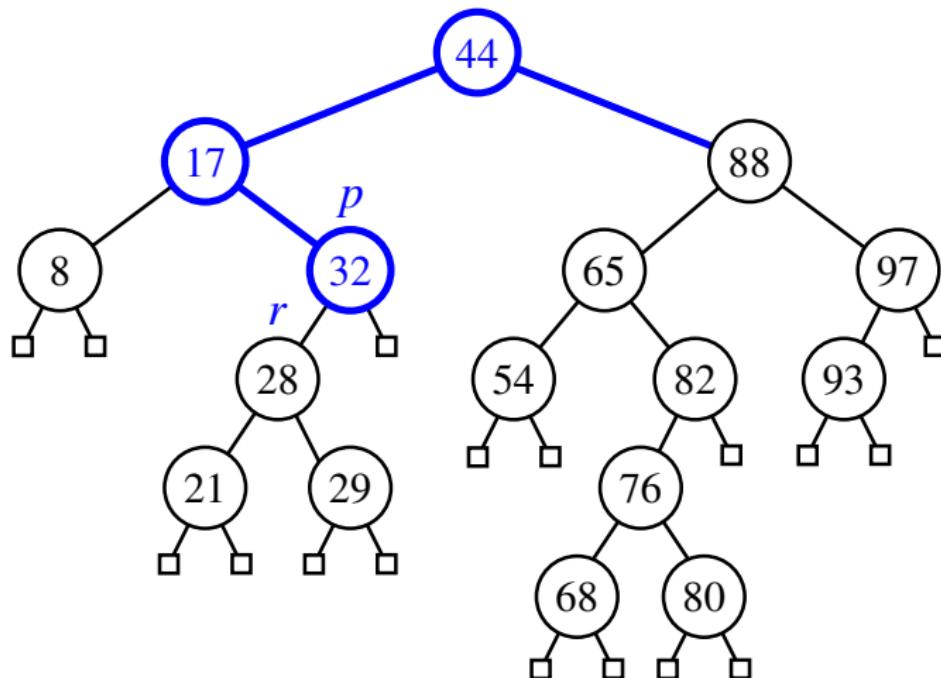
- An **external, RC** & an **internal LC** \Rightarrow Remove r as in **Case 3**.

Running time?

[$O(h)$]

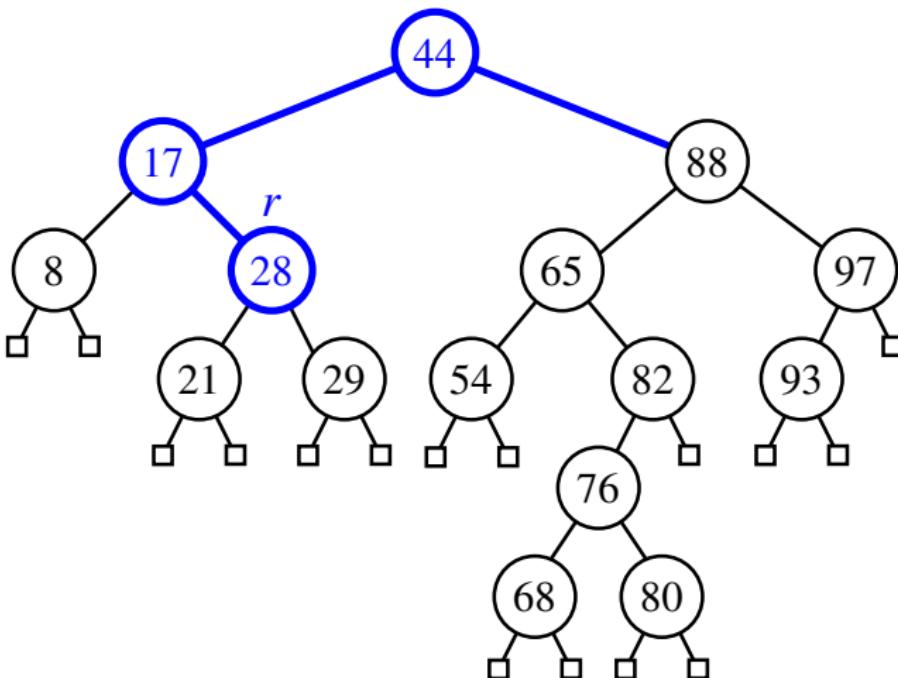
Visualizing BST Operation: Deletion (1.1)

(Case 3) Before *deleting* the node storing **key 32**:



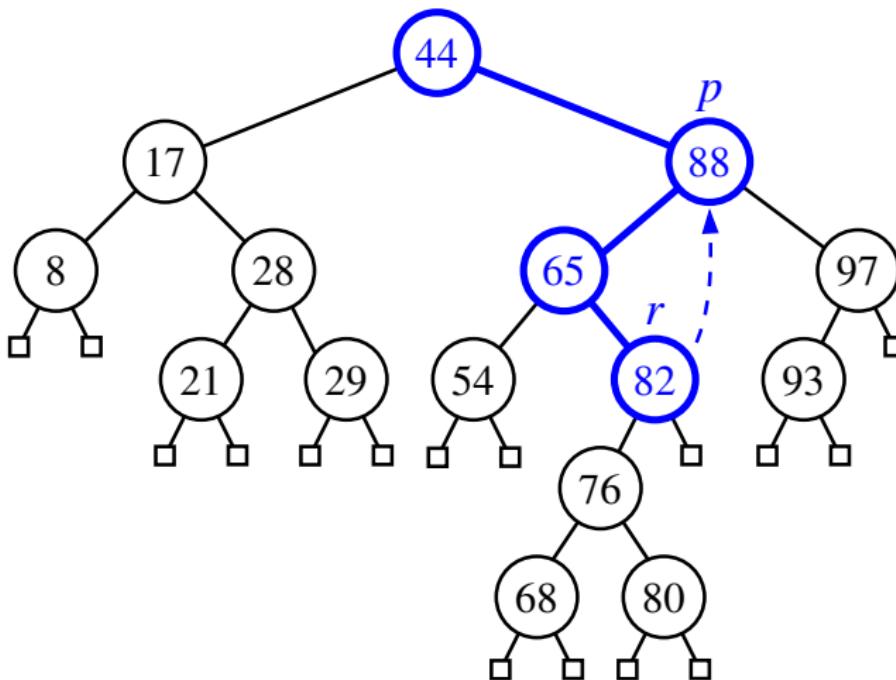
Visualizing BST Operation: Deletion (1.2)

(Case 3) After *deleting* the node storing **key 32**:



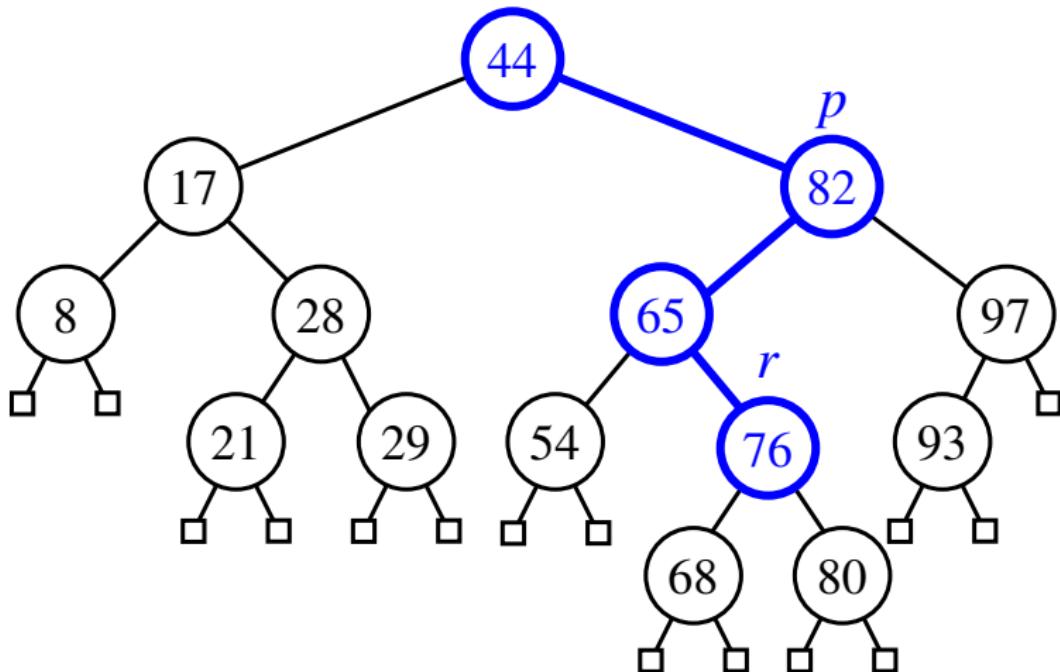
Visualizing BST Operation: Deletion (2.1)

(Case 4) Before *deleting* the node storing **key 88**:



Visualizing BST Operation: Deletion (2.2)

(Case 4) After *deleting* the node storing **key 88**:



Exercise on BST Operation: Deletion

Exercise : In BSTUtilities class, **implement** and **test** the
`void delete(BSTNode<E> p, int k)` method.

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