

Design-by-Contract (DbC)

Readings: OOSC2 Chapter 11



EECS3311 A: Software Design
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What This Course Is About

- Focus is **design**
 - **Architecture**: (many) **inter-related** modules
 - **Specification**: **precise** (functional) interface of each module
- For this course, having a prototypical, **working** implementation for your design suffices.
- A later **refinement** into more efficient data structures and algorithms is beyond the scope of this course.

[assumed from EECS2011, EECS3101]

∴ Having a suitable language for **design** matters the most.

Q: Is Java also a “good” **design** language?

A: Let’s first understand what a “good” **design** is.

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Motivation: Catching Defects – When?



- To minimize **development costs**, minimize **software defects**.

- Software Development Cycle:

Requirements → **Design** → **Implementation** → Release

Q. Design or Implementation Phase?

Catch defects **as early as possible**.

| Design and architecture | Implementation | Integration testing | Customer beta test | Postproduct release |
|-------------------------|----------------|---------------------|--------------------|---------------------|
| 1X* | 5X | 10X | 15X | 30X |

∴ The cost of fixing defects **increases exponentially** as software progresses through the development lifecycle.

- Discovering **defects** after **release** costs up to 30 times more than catching them in the **design** phase.
- Choice of **design language** for your project is therefore of paramount importance.

Source: Minimizing code defects to improve software quality and lower development costs.

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Terminology: Contract, Client, Supplier



- A **supplier** implements/provides a service (e.g., microwave).
- A **client** uses a service provided by some supplier.
 - The client is required to follow certain instructions to obtain the service (e.g., supplier **assumes** that client powers on, closes door, and heats something that is not explosive).
 - If instructions are followed, the client would **expect** that the service does what is guaranteed (e.g., a lunch box is heated).
 - The client does not care how the supplier implements it.
- What then are the **benefits** and **obligations** of the two parties?

| | benefits | obligations |
|----------|------------------------------|---------------------|
| CLIENT | obtain a service | follow instructions |
| SUPPLIER | assume instructions followed | provide a service |

- There is a **contract** between two parties, violated if:
 - The instructions are not followed. [Client’s fault]
 - Instructions followed, but service not satisfactory. [Supplier’s fault]

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Client, Supplier, Contract in OOP (1)



```
class Microwave {
    private boolean on;
    private boolean locked;
    void power() {on = true;}
    void lock() {locked = true;}
    void heat(Object stuff) {
        /* Assume: on && locked */
        /* stuff not explosive. */
    }
}
```

```
class MicrowaveUser {
    public static void main(...) {
        Microwave m = new Microwave();
        Object obj = ???;
        m.power(); m.lock();
        m.heat(obj);
    }
}
```

Method call `m.heat(obj)` indicates a client-supplier relation.

- **Client:** resident class of the method call [MicrowaveUser]
- **Supplier:** type of context object (or call target) `m` [Microwave]

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What is a Good Design?



- A “good” design should *explicitly* and *unambiguously* describe the **contract** between **clients** (e.g., users of Java classes) and **suppliers** (e.g., developers of Java classes). We call such a contractual relation a **specification**.
- When you conduct *software design*, you should be guided by the “appropriate” contracts between users and developers.
 - Instructions to **clients** should *not be unreasonable*.
e.g., asking them to assemble internal parts of a microwave
 - Working conditions for **suppliers** should *not be unconditional*.
e.g., expecting them to produce a microwave which can safely heat an explosive with its door open!
 - You as a designer should strike proper balance between **obligations** and **benefits** of clients and suppliers.
e.g., What is the obligation of a binary-search user (also benefit of a binary-search implementer)? [The input array is *sorted*.]
 - Upon contract violation, there should be the fault of **only one side**.
 - This design process is called **Design by Contract (DbC)**.

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Client, Supplier, Contract in OOP (2)



```
class Microwave {
    private boolean on;
    private boolean locked;
    void power() {on = true;}
    void lock() {locked = true;}
    void heat(Object stuff) {
        /* Assume: on && locked */
        /* stuff not explosive. */
    }
}
```

```
class MicrowaveUser {
    public static void main(...) {
        Microwave m = new Microwave();
        Object obj = ???;
        m.power(); m.lock();
        m.heat(obj);
    }
}
```

- The **contract** is *honoured* if:

Right **before** the method call:

- State of `m` is as assumed: `m.on==true` and `m.locked==ture`
- The input argument `obj` is valid (i.e., not explosive).

Right **after** the method call: `obj` is properly heated.

- If any of these fails, there is a **contract violation**.
 - `m.on` or `m.locked` is false ⇒ MicrowaveUser's fault.
 - `obj` is an explosive ⇒ MicrowaveUser's fault.
 - A fault from the client is identified ⇒ Method call will not start.
 - Method executed but `obj` not properly heated ⇒ Microwave's fault

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A Simple Problem: Bank Accounts



Provide an object-oriented solution to the following problem:

- REQ1**: Each account is associated with the *name* of its owner (e.g., "Jim") and an integer *balance* that is always positive.
- REQ2**: We may *withdraw* an integer amount from an account.
- REQ3**: Each bank stores a list of *accounts*.
- REQ4**: Given a bank, we may *add* a new account in it.
- REQ5**: Given a bank, we may *query* about the associated account of a owner (e.g., the account of "Jim").
- REQ6**: Given a bank, we may *withdraw* from a specific account, identified by its name, for an integer amount.

Let's first try to work on **REQ1** and **REQ2** in Java.
This may not be as easy as you might think!

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Playing the Various Versions in Java



- **Download** the project archive (a zip file) here:
<http://www.eecs.yorku.ca/~jackie/teaching/lectures/2020/W/EECS3311/codes/DbCIntro.zip>
- Follow this tutorial to learn how to **import** an project archive into your workspace in Eclipse:
<https://youtu.be/h-rqdOZa2aY>
- Follow this tutorial to learn how to **enable** assertions in Eclipse:
<https://youtu.be/OEgRV4a5Dza>

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Version 1: Why Not a Good Design? (1)



```
public class BankAppV1 {
    public static void main(String[] args) {
        System.out.println("Create an account for Alan with balance -10:");
        AccountV1 alan = new AccountV1("Alan", -10);
        System.out.println(alan);
    }
}
```

Console Output:

```
Create an account for Alan with balance -10:
Alan's current balance is: -10
```

- Executing AccountV1's constructor results in an account object whose **state** (i.e., values of attributes) is **invalid** (i.e., Alan's balance is negative). ⇒ Violation of **REQ1**
- Unfortunately, both client and supplier are to be blamed: BankAppV1 passed an invalid balance, but the API of AccountV1 does not require that! ⇒ A lack of defined contract

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Version 1: An Account Class



```
1 public class AccountV1 {
2     private String owner;
3     private int balance;
4     public String getOwner() { return owner; }
5     public int getBalance() { return balance; }
6     public AccountV1(String owner, int balance) {
7         this.owner = owner; this.balance = balance;
8     }
9     public void withdraw(int amount) {
10        this.balance = this.balance - amount;
11    }
12    public String toString() {
13        return owner + "'s current balance is: " + balance;
14    }
15 }
```

- Is this a good design? Recall **REQ1**: Each account is associated with ... an integer balance that is **always positive**.
- This requirement is **not** reflected in the above Java code.

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Version 1: Why Not a Good Design? (2)



```
public class BankAppV1 {
    public static void main(String[] args) {
        System.out.println("Create an account for Mark with balance 100:");
        AccountV1 mark = new AccountV1("Mark", 100);
        System.out.println(mark);
        System.out.println("Withdraw -1000000 from Mark's account:");
        mark.withdraw(-1000000);
        System.out.println(mark);
    }
}
```

```
Create an account for Mark with balance 100:
Mark's current balance is: 100
Withdraw -1000000 from Mark's account:
Mark's current balance is: 1000100
```

- Mark's account state is always valid (i.e., 100 and 1000100).
- Withdraw amount is never negative! ⇒ Violation of **REQ2**
- Again a lack of contract between BankAppV1 and AccountV1.

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Version 1: Why Not a Good Design? (3)



```
public class BankAppV1 {
    public static void main(String[] args) {
        System.out.println("Create an account for Tom with balance 100:");
        AccountV1 tom = new AccountV1("Tom", 100);
        System.out.println(tom);
        System.out.println("Withdraw 150 from Tom's account:");
        tom.withdraw(150);
        System.out.println(tom);
    }
}
```

```
Create an account for Tom with balance 100:
Tom's current balance is: 100
Withdraw 150 from Tom's account:
Tom's current balance is: -50
```

- Withdrawal was done via an “appropriate” reduction, but the resulting balance of Tom is *invalid*. ⇒ Violation of **REQ1**
- Again a lack of contract between BankAppV1 and AccountV1.

REQ1

Version 1: How Should We Improve it? (1)



Preconditions of a method specify the precise circumstances under which that method can be executed.

- Precond. of `divide(int x, int y)`? [$y \neq 0$]
- Precond. of `binSearch(int x, int[] xs)`? [`xs` is sorted]
- Precond. of `topoSort(Graph g)`? [`g` is a DAG]

REQ1

Version 1: How Should We Improve it? (2)



- The best we can do in Java is to encode the **logical negations** of preconditions as **exceptions**:
 - `divide(int x, int y)`
throws `DivisionByZeroException` when $y == 0$.
 - `binSearch(int x, int[] xs)`
throws `ArrayNotSortedException` when `xs` is *not* sorted.
 - `topoSort(Graph g)`
throws `NotDAGException` when `g` is *not* directed and acyclic.
- Design your method by specifying the **preconditions** (i.e., **service** conditions for **valid** inputs) it requires, not the **exceptions** (i.e., **error** conditions for **invalid** inputs) for it to fail.
- Create **Version 2** by adding **exceptional conditions** (an **approximation** of **preconditions**) to the constructor and `withdraw` method of the `Account` class.

REQ1

Version 2: Added Exceptions to Approximate Method Preconditions



```
1 public class AccountV2 {
2     public AccountV2(String owner, int balance) throws
3         BalanceNegativeException
4     {
5         if (balance < 0) { /* negated precondition */
6             throw new BalanceNegativeException(); }
7         else { this.owner = owner; this.balance = balance; }
8     }
9     public void withdraw(int amount) throws
10        WithdrawAmountNegativeException, WithdrawAmountTooLargeException {
11         if (amount < 0) { /* negated precondition */
12             throw new WithdrawAmountNegativeException(); }
13         else if (balance < amount) { /* negated precondition */
14             throw new WithdrawAmountTooLargeException(); }
15         else { this.balance = this.balance - amount; }
16     }
}
```

REQ1

Version 2: Why Better than Version 1? (1)



```
1 public class BankAppV2 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Alan with balance -10:");
4         try {
5             AccountV2 alan = new AccountV2("Alan", -10);
6             System.out.println(alan);
7         }
8         catch (BalanceNegativeException bne) {
9             System.out.println("Illegal negative account balance.");
10        }
11    }
12 }
```

```
Create an account for Alan with balance -10:
Illegal negative account balance.
```

L6: When attempting to call the constructor AccountV2 with a negative balance -10, a BalanceNegativeException (i.e., precondition violation) occurs, preventing further operations upon this invalid object.

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Version 2: Why Better than Version 1? (2.2)



Console Output:

```
Create an account for Mark with balance 100:
Mark's current balance is: 100
Withdraw -1000000 from Mark's account:
Illegal negative withdraw amount.
```

- L8: When attempting to call method withdraw with a negative amount -1000000, a WithdrawAmountNegativeException (i.e., precondition violation) occurs, preventing the withdrawal from proceeding.
- We should observe that adding preconditions to the supplier BankV2's code forces the client BankAppV2's code to get complicated by the try-catch statements.
- Adding clear contract (preconditions in this case) to the design should not be at the cost of complicating the client's code!!

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Version 2: Why Better than Version 1? (2.1)



```
1 public class BankAppV2 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Mark with balance 100:");
4         try {
5             AccountV2 mark = new AccountV2("Mark", 100);
6             System.out.println(mark);
7             System.out.println("Withdraw -1000000 from Mark's account:");
8             mark.withdraw(-1000000);
9             System.out.println(mark);
10        }
11        catch (BalanceNegativeException bne) {
12            System.out.println("Illegal negative account balance.");
13        }
14        catch (WithdrawAmountNegativeException wane) {
15            System.out.println("Illegal negative withdraw amount.");
16        }
17        catch (WithdrawAmountTooLargeException wane) {
18            System.out.println("Illegal too large withdraw amount.");
19        }
20    }
21 }
```

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Version 2: Why Better than Version 1? (3.1)



```
1 public class BankAppV2 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Tom with balance 100:");
4         try {
5             AccountV2 tom = new AccountV2("Tom", 100);
6             System.out.println(tom);
7             System.out.println("Withdraw 150 from Tom's account:");
8             tom.withdraw(150);
9             System.out.println(tom);
10        }
11        catch (BalanceNegativeException bne) {
12            System.out.println("Illegal negative account balance.");
13        }
14        catch (WithdrawAmountNegativeException wane) {
15            System.out.println("Illegal negative withdraw amount.");
16        }
17        catch (WithdrawAmountTooLargeException wane) {
18            System.out.println("Illegal too large withdraw amount.");
19        }
20    }
21 }
```

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Version 2: Why Better than Version 1? (3.2)



Console Output:

```
Create an account for Tom with balance 100:
Tom's current balance is: 100
Withdraw 150 from Tom's account:
Illegal too large withdraw amount.
```

- **L8:** When attempting to call method `withdraw` with a positive but too large amount 150, a `WithdrawAmountTooLargeException` (i.e., **precondition violation**) occurs, *preventing the withdrawal from proceeding*.
- We should observe that due to the **added preconditions** to the supplier `BankV2`'s code, the client `BankAppV2`'s code is forced to *repeat the long list of the try-catch statements*.
- Indeed, adding clear contract (**preconditions** in this case) **should not** be at the cost of complicating the client's code!!

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Version 2: Why Still Not a Good Design? (2.1)



```
1 public class BankAppV2 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Jim with balance 100:");
4         try {
5             AccountV2 jim = new AccountV2("Jim", 100);
6             System.out.println(jim);
7             System.out.println("Withdraw 100 from Jim's account:");
8             jim.withdraw(100);
9             System.out.println(jim);
10        }
11        catch (BalanceNegativeException bne) {
12            System.out.println("Illegal negative account balance.");
13        }
14        catch (WithdrawAmountNegativeException wane) {
15            System.out.println("Illegal negative withdraw amount.");
16        }
17        catch (WithdrawAmountTooLargeException wane) {
18            System.out.println("Illegal too large withdraw amount.");
19        }
20    }
21 }
```

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Version 2: Why Still Not a Good Design? (1)



```
1 public class AccountV2 {
2     public AccountV2(String owner, int balance) throws
3         BalanceNegativeException
4     {
5         if (balance < 0) { /* negated precondition */
6             throw new BalanceNegativeException(); }
7         else { this.owner = owner; this.balance = balance; }
8     }
9     public void withdraw(int amount) throws
10        WithdrawAmountNegativeException, WithdrawAmountTooLargeException {
11        if (amount < 0) { /* negated precondition */
12            throw new WithdrawAmountNegativeException(); }
13        else if (balance < amount) { /* negated precondition */
14            throw new WithdrawAmountTooLargeException(); }
15        else { this.balance = this.balance - amount; }
16    }
17 }
```

- Are all the **exception** conditions (\neg **preconditions**) appropriate?
- What if `amount == balance` when calling `withdraw`?

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Version 2: Why Still Not a Good Design? (2.2)



```
Create an account for Jim with balance 100:
Jim's current balance is: 100
Withdraw 100 from Jim's account:
Jim's current balance is: 0
```

- **L9:** When attempting to call method `withdraw` with an amount 100 (i.e., equal to Jim's current balance) that would result in a **zero balance** (clearly a violation of **REQ1**), there should have been a **precondition violation**.

Supplier `AccountV2`'s **exception** condition `balance < amount` has a **missing case**:

- Calling `withdraw` with `amount == balance` will also result in an invalid account state (i.e., the resulting account balance is **zero**).
- \therefore **L13** of `AccountV2` should be `balance <= amount`.

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Version 2: How Should We Improve it?



- Even without fixing this insufficient *precondition*, we could have avoided the above scenario by *checking at the end of each method that the resulting account is valid*.
 - ⇒ We consider the condition `this.balance > 0` as **invariant** throughout the lifetime of all instances of `Account`.
- **Invariants** of a class specify the precise conditions which *all instances/objects* of that class must satisfy.
 - Inv. of `CSMajorStudent`? [`gpa >= 4.5`]
 - Inv. of `BinarySearchTree`? [in-order trav. → sorted key seq.]
- The best we can do in Java is encode invariants as **assertions**:
 - `CSMajorStudent`: **assert** `this.gpa >= 4.5`
 - `BinarySearchTree`: **assert** `this.inOrder()` is sorted
 - Unlike exceptions, assertions are not in the class/method API.
- Create **Version 3** by adding **assertions** to the end of constructor and `withdraw` method of the `Account` class.

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Version 3: Added Assertions to Approximate Class Invariants



```
1 public class AccountV3 {
2     public AccountV3(String owner, int balance) throws
3         BalanceNegativeException
4     {
5         if(balance < 0) { /* negated precondition */
6             throw new BalanceNegativeException(); }
7         else { this.owner = owner; this.balance = balance; }
8         assert this.getBalance() > 0 : "Invariant: positive balance";
9     }
10    public void withdraw(int amount) throws
11        WithdrawAmountNegativeException, WithdrawAmountTooLargeException {
12        if(amount < 0) { /* negated precondition */
13            throw new WithdrawAmountNegativeException(); }
14        else if (balance < amount) { /* negated precondition */
15            throw new WithdrawAmountTooLargeException(); }
16        else { this.balance = this.balance - amount; }
17        assert this.getBalance() > 0 : "Invariant: positive balance";
18    }
```

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Version 3: Why Better than Version 2?



```
1 public class BankAppV3 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Jim with balance 100:");
4         try { AccountV3 jim = new AccountV3("Jim", 100);
5             System.out.println(jim);
6             System.out.println("Withdraw 100 from Jim's account:");
7             jim.withdraw(100);
8             System.out.println(jim); }
9         /* catch statements same as this previous slide:
10        * Version 2: Why Still Not a Good Design? (2.1) */
```

```
Create an account for Jim with balance 100:
Jim's current balance is: 100
Withdraw 100 from Jim's account:
Exception in thread "main"
```

java.lang.AssertionError: Invariant: positive balance

L8: Upon completion of `jim.withdraw(100)`, Jim has a zero balance, an assertion failure (i.e., **invariant** violation) occurs, *preventing further operations on this invalid account object.*

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Version 3: Why Still Not a Good Design?



Let's recall what we have added to the method `withdraw`:

- From **Version 2**: **exceptions** encoding **negated preconditions**
- From **Version 3**: **assertions** encoding the **class invariants**

```
1 public class AccountV3 {
2     public void withdraw(int amount) throws
3         WithdrawAmountNegativeException, WithdrawAmountTooLargeException {
4         if(amount < 0) { /* negated precondition */
5             throw new WithdrawAmountNegativeException(); }
6         else if (balance < amount) { /* negated precondition */
7             throw new WithdrawAmountTooLargeException(); }
8         else { this.balance = this.balance - amount; }
9         assert this.getBalance() > 0 : "Invariant: positive balance";
}
```

However, there is **no contract** in `withdraw` which specifies:

- Obligations of supplier (`AccountV3`) if preconditions are met.
 - Benefits of client (`BankAppV3`) after meeting preconditions.
- ⇒ We illustrate how problematic this can be by creating

Version 4, where deliberately mistakenly implement `withdraw`.

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Version 4: What If the Implementation of withdraw is Wrong? (1)

```

1 public class AccountV4 {
2     public void withdraw(int amount) throws
3         WithdrawAmountNegativeException, WithdrawAmountTooLargeException
4     { if(amount < 0) { /* negated precondition */
5         throw new WithdrawAmountNegativeException(); }
6         else if (balance < amount) { /* negated precondition */
7             throw new WithdrawAmountTooLargeException(); }
8         else { /* WRONG IMPLEMENTATION */
9             this.balance = this.balance + amount; }
10        assert this.getBalance() > 0 :
11            owner + "Invariant: positive balance"; }

```

- Apparently the implementation at L11 is **wrong**.
- Adding a positive amount to a valid (positive) account balance would not result in an invalid (negative) one.
⇒ The **class invariant** will **not** catch this flaw.
- When something goes wrong, a good **design** (with an appropriate **contract**) should report it via a **contract violation**.

REQ1

Version 4: What If the Implementation of withdraw is Wrong? (2)

```

1 public class BankAppV4 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Jeremy with balance 100:");
4         try { AccountV4 jeremy = new AccountV4("Jeremy", 100);
5             System.out.println(jeremy);
6             System.out.println("Withdraw 50 from Jeremy's account:");
7             jeremy.withdraw(50);
8             System.out.println(jeremy); }
9         /* catch statements same as this previous slide:
10          * Version 2: Why Still Not a Good Design? (2.1) */

```

```

Create an account for Jeremy with balance 100:
Jeremy's current balance is: 100
Withdraw 50 from Jeremy's account:
Jeremy's current balance is: 150

```

L7: Resulting balance of Jeremy is valid (150 > 0), but withdrawal was done via an **mistaken** increase. ⇒ Violation of **REQ2**

REQ1

Version 4: How Should We Improve it?

- **Postconditions** of a method specify the precise conditions which it will satisfy upon its completion.
 - This relies on the assumption that right before the method starts, its preconditions are satisfied (i.e., inputs valid) and invariants are satisfied (i.e., object state valid).
 - Postcondition of double divide(int x, int y)?
[**Result** × y == x]
 - Postcondition of boolean binSearch(int x, int[] xs)?
[x ∈ xs ⇔ **Result**]
- The best we can do in Java is, similar to the case of invariants, encode postconditions as **assertions**.
But again, unlike exceptions, these assertions will not be part of the class/method API.
- Create **Version 5** by adding **assertions** to the end of withdraw method of the Account class.

REQ1

Version 5: Added Assertions to Approximate Method Postconditions

```

1 public class AccountV5 {
2     public void withdraw(int amount) throws
3         WithdrawAmountNegativeException, WithdrawAmountTooLargeException {
4         int oldBalance = this.balance;
5         if(amount < 0) { /* negated precondition */
6             throw new WithdrawAmountNegativeException(); }
7         else if (balance < amount) { /* negated precondition */
8             throw new WithdrawAmountTooLargeException(); }
9         else { this.balance = this.balance - amount; }
10        assert this.getBalance() > 0 : "Invariant: positive balance";
11        assert this.getBalance() == oldBalance - amount :
12            "Postcondition: balance deducted"; }

```

A postcondition typically **relates** the **pre-execution value** and the **post-execution value** of each relevant attribute

(e.g., balance in the case of withdraw).
⇒ Extra code (L4) to capture the pre-execution value of balance for the comparison at L11.

REQ1

Version 5: Why Better than Version 4?



```

1 public class BankAppV5 {
2     public static void main(String[] args) {
3         System.out.println("Create an account for Jeremy with balance 100:");
4         try { AccountV5 jeremy = new AccountV5("Jeremy", 100);
5             System.out.println(jeremy);
6             System.out.println("Withdraw 50 from Jeremy's account:");
7             jeremy.withdraw(50);
8             System.out.println(jeremy); }
9         /* catch statements same as this previous slide:
10        * Version 2: Why Still Not a Good Design? (2.1) */
    
```

```

Create an account for Jeremy with balance 100:
Jeremy's current balance is: 100
Withdraw 50 from Jeremy's account:
Exception in thread "main"
    java.lang.AssertionError: Postcondition: balance deducted
    
```

L8: Upon completion of `jeremy.withdraw(50)`, Jeremy has a wrong balance 150, an assertion failure (i.e., **postcondition** violation) occurs, *preventing further operations on this invalid account object*.

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Version 5: Contract between Client and Supplier



| | benefits | obligations |
|-------------------------------|---|---|
| BankAppV5.main (CLIENT) | balance deduction positive balance | amount non-negative amount not too large |
| BankV5.withdraw (SUPPLIER) | amount non-negative amount not too large | balance deduction positive balance |

| | benefits | obligations |
|----------|---------------------------|---------------------------|
| CLIENT | postcondition & invariant | precondition |
| SUPPLIER | precondition | postcondition & invariant |

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Evolving from Version 1 to Version 5



| | Improvements Made | Design Flaws |
|----|---|--|
| V1 | – | Complete lack of Contract |
| V2 | Added exceptions as <i>method preconditions</i> | Preconditions not strong enough (i.e., with missing cases) may result in an invalid account state. |
| V3 | Added assertions as <i>class invariants</i> | – |
| V4 | Deliberately changed <code>withdraw</code> 's implementation to be incorrect . | Incorrect implementations do not necessarily result in a state that violates the class invariants. |
| V5 | Added assertions as <i>method postconditions</i> | – |

- In Versions 2, 3, 4, 5, **preconditions** approximated as *exceptions*.
 - These are **not preconditions**, but their **logical negation**.
 - Client `BankApp`'s code **complicated** by repeating the list of `try-catch` statements.
- In Versions 3, 4, 5, **class invariants** and **postconditions** approximated as *assertions*.
 - Unlike exceptions, these assertions will **not appear in the API** of `withdraw`. Potential clients of this method **cannot know**: 1) what their benefits are; and 2) what their suppliers' obligations are.
 - For postconditions, **extra code** needed to capture pre-execution values of attributes.

8/27/61

DbC in Java



DbC is possible in Java, but not appropriate for your learning:

- Preconditions** of a method:
 - Supplier**
 - Encode their logical negations as exceptions.
 - In the **beginning** of that method, a list of `if`-statements for throwing the appropriate exceptions.
 - Client**
 - A list of `try-catch`-statements for handling exceptions.
- Postconditions** of a method:
 - Supplier**
 - Encoded as a list of assertions, placed at the **end** of that method.
 - Client**
 - All such assertions do not appear in the API of that method.
- Invariants** of a class:
 - Supplier**
 - Encoded as a list of assertions, placed at the **end** of every method.
 - Client**
 - All such assertions do not appear in the API of that class.

8/27/61

Why Java Interfaces Unacceptable ADTs (1)



Interface List<E>

Type Parameters:

E - the type of elements in this list

All Superinterfaces:

Collection<E>, Iterable<E>

All Known Implementing Classes:

AbstractList, AbstractSequentialList, ArrayList, AttributeList, CopyOnWriteArrayList, LinkedList, RoleList, RoleUnresolvedList, Stack, Vector

```
public interface List<E>
    extends Collection<E>
```

An ordered collection (also known as a *sequence*). The user of this interface has precise control over where in the list each element is inserted. The user can access elements by their integer index (position in the list), and search for elements in the list.

It is useful to have:

- A **generic collection class** where the **homogeneous type** of elements are parameterized as E.
- A reasonably **intuitive overview** of the ADT.

Java 8 List API

3/3/2014

DbC in Eiffel: Supplier



DbC is supported natively in Eiffel for **supplier**:

```
class ACCOUNT
create
    make
feature -- Attributes
    owner : STRING
    balance : INTEGER
feature -- Constructors
    make(nn: STRING; nb: INTEGER)
        require -- precondition
            positive_balance: nb > 0
        do
            owner := nn
            balance := nb
        end
feature -- Commands
    withdraw(amount: INTEGER)
        require -- precondition
            non_negative_amount: amount > 0
            affordable_amount: amount <= balance -- problematic, why?
        do
            balance := balance - amount
        ensure -- postcondition
            balance_deducted: balance = old balance - amount
        end
invariant -- class invariant
    positive_balance: balance > 0
end
```

3/3/2014

Why Java Interfaces Unacceptable ADTs (2)



Methods described in a **natural language** can be **ambiguous**:

```
E set(int index, E element)
    Replaces the element at the specified position in this list with the specified element (optional operation).
```

```
set
E set(int index,
    E element)
```

Replaces the element at the specified position in this list with the specified element (optional operation).

Parameters:

index - index of the element to replace

element - element to be stored at the specified position

Returns:

the element previously at the specified position

Throws:

UnsupportedOperationException - if the set operation is not supported by this list

ClassCastException - if the class of the specified element prevents it from being added to this list

NullPointerException - if the specified element is null and this list does not permit null elements

IllegalArgumentException - if some property of the specified element prevents it from being added to this list

IndexOutOfBoundsException - if the index is out of range (index < 0 || index >= size())

3/3/2014

DbC in Eiffel: Contract View of Supplier



Any potential **client** who is interested in learning about the kind of services provided by a **supplier** can look through the **contract view** (without showing any implementation details):

```
class ACCOUNT
create
    make
feature -- Attributes
    owner : STRING
    balance : INTEGER
feature -- Constructors
    make(nn: STRING; nb: INTEGER)
        require -- precondition
            positive_balance: nb > 0
        end
feature -- Commands
    withdraw(amount: INTEGER)
        require -- precondition
            non_negative_amount: amount > 0
            affordable_amount: amount <= balance -- problematic, why?
        ensure -- postcondition
            balance_deducted: balance = old balance - amount
        end
invariant -- class invariant
    positive_balance: balance > 0
end
```

3/3/2014

DbC in Eiffel: Anatomy of a Class

```

class SOME_CLASS
create
  -- Explicitly list here commands used as constructors
feature -- Attributes
  -- Declare attribute here
feature -- Commands
  -- Declare commands (mutators) here
feature -- Queries
  -- Declare queries (accessors) here
invariant
  -- List of tagged boolean expressions for class invariants
end
  
```

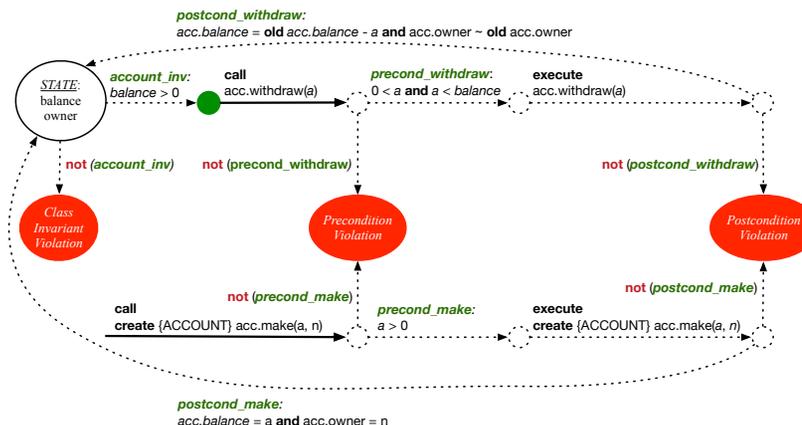
- Use **feature** clauses to group attributes, commands, queries.
- Explicitly declare list of commands under **create** clause, so that they can be used as class constructors.

[See the groups panel in Eiffel Studio.]

- The **class invariant invariant** clause may be omitted:
 - There's no class invariant: any resulting object state is acceptable.
 - The class invariant is equivalent to writing **invariant true**

Runtime Monitoring of Contracts (1)

In the specific case of **ACCOUNT** class with creation procedure **make** and command **withdraw**:



DbC in Eiffel: Anatomy of a Feature

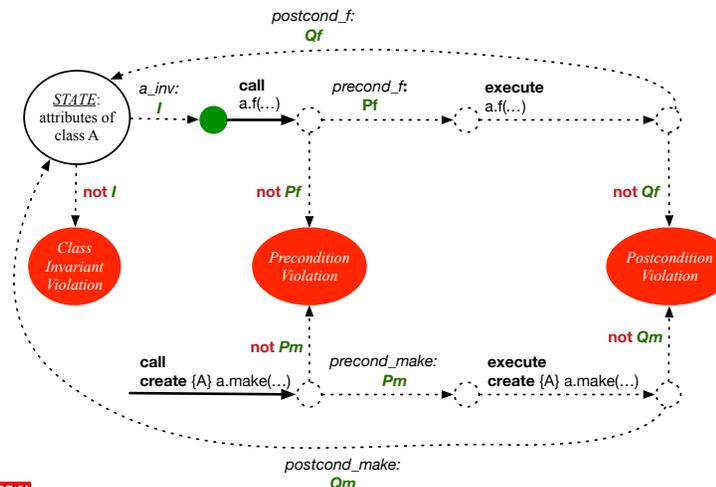
```

some_command
  -- Description of the command.
  require
    -- List of tagged boolean expressions for preconditions
  local
    -- List of local variable declarations
  do
    -- List of instructions as implementation
  ensure
    -- List of tagged boolean expressions for postconditions
  end
  
```

- The **precondition require** clause may be omitted:
 - There's no precondition: any starting state is acceptable.
 - The precondition is equivalent to writing **require true**
- The **postcondition ensure** clause may be omitted:
 - There's no postcondition: any resulting state is acceptable.
 - The postcondition is equivalent to writing **ensure true**

Runtime Monitoring of Contracts (2)

In general, class **C** with creation procedure **cp** and any feature **f**:



Runtime Monitoring of Contracts (3)



- All **contracts** are specified as **Boolean expressions**.
- Right **before** a feature call (e.g., `acc.withdraw(10)`):
 - The current state of `acc` is called the **pre-state**.
 - Evaluate feature `withdraw`'s **pre-condition** using current values of attributes and queries.
 - **Cache** values (**implicitly**) of all expressions involving the **old** keyword in the **post-condition**.
e.g., cache the value of `old balance` via `old_balance := balance`
- Right **after** the feature call:
 - The current state of `acc` is called the **post-state**.
 - Evaluate class `ACCOUNT`'s **invariant** using current values of attributes and queries.
 - Evaluate feature `withdraw`'s **post-condition** using both current and **"cached"** values of attributes and queries.

3/5/2011

DbC in Eiffel: Precondition Violation (1.2)



3/5/2011

DbC in Eiffel: Precondition Violation (1.1)



The **client** need not handle all possible contract violations:

```
class BANK_APP
inherit
  ARGUMENTS
create
  make
feature -- Initialization
  make
  -- Run application.
  local
    alan: ACCOUNT
  do
    -- A precondition violation with tag "positive_balance"
    create {ACCOUNT} alan.make ("Alan", -10)
  end
end
```

By executing the above code, the runtime monitor of Eiffel Studio will report a **contract violation** (precondition violation with tag "positive_balance").

3/5/2011

DbC in Eiffel: Precondition Violation (2.1)



The **client** need not handle all possible contract violations:

```
class BANK_APP
inherit
  ARGUMENTS
create
  make
feature -- Initialization
  make
  -- Run application.
  local
    mark: ACCOUNT
  do
    create {ACCOUNT} mark.make ("Mark", 100)
    -- A precondition violation with tag "non_negative_amount"
    mark.withdraw(-1000000)
  end
end
```

By executing the above code, the runtime monitor of Eiffel Studio will report a **contract violation** (precondition violation with tag "non_negative_amount").

3/5/2011

DbC in Eiffel: Precondition Violation (2.2)



```
APPLICATION: ACCOUNT
Feature
bank ACCOUNT withdraw
Flat view of feature 'withdraw' of class ACCOUNT
withdraw (amount: INTEGER_32)
  require
    non_negative_amount: amount >= 0
    affordable_amount: amount <= balance
  do
    balance := balance - amount
  ensure
    balance = old balance - amount
  end
```

Contract

DbC in Eiffel: Precondition Violation (3.2)



```
APPLICATION: ACCOUNT
Feature
bank ACCOUNT withdraw
Flat view of feature 'withdraw' of class ACCOUNT
withdraw (amount: INTEGER_32)
  require
    non_negative_amount: amount >= 0
    affordable_amount: amount <= balance
  do
    balance := balance - amount
  ensure
    balance = old balance - amount
  end
```

Contract

DbC in Eiffel: Precondition Violation (3.1)



The **client** need not handle all possible contract violations:

```
class BANK_APP
inherit
  ARGUMENTS
create
  make
feature -- Initialization
  make
  -- Run application.
local
  tom: ACCOUNT
do
  create {ACCOUNT} tom.make ("Tom", 100)
  -- A precondition violation with tag "affordable_amount"
  tom.withdraw(150)
end
end
```

By executing the above code, the runtime monitor of Eiffel Studio will report a **contract violation** (precondition violation with tag "affordable_amount").

Contract

DbC in Eiffel: Class Invariant Violation (4.1)



The **client** need not handle all possible contract violations:

```
class BANK_APP
inherit
  ARGUMENTS
create
  make
feature -- Initialization
  make
  -- Run application.
local
  jim: ACCOUNT
do
  create {ACCOUNT} tom.make ("Jim", 100)
  jim.withdraw(100)
  -- A class invariant violation with tag "positive_balance"
end
end
```

By executing the above code, the runtime monitor of Eiffel Studio will report a **contract violation** (class invariant violation with tag "positive_balance").

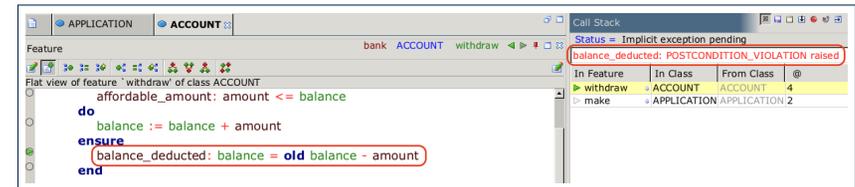
Contract

DbC in Eiffel: Class Invariant Violation (4.2)



5/20/2011

DbC in Eiffel: Postcondition Violation (5.2)



5/20/2011

DbC in Eiffel: Postcondition Violation (5.1)



The **client** need not handle all possible contract violations:

```
class BANK_APP
inherit ARGUMENTS
create make
feature -- Initialization
  make
  -- Run application.
  local
    jeremy: ACCOUNT
  do
  -- Faulty implementation of withdraw in ACCOUNT:
  -- balance := balance + amount
  create {ACCOUNT} jeremy.make ("Jeremy", 100)
  jeremy.withdraw(150)
  -- A postcondition violation with tag "balance_deducted"
  end
end
```

By executing the above code, the runtime monitor of Eiffel Studio will report a **contract violation** (postcondition violation with tag "balance_deducted").

5/20/2011

Beyond this lecture...



- Study this tutorial series on DbC and TDD:

<https://www.youtube.com/playlist?list=PL5dxAmCm1v6r5VfzCQ5bTznoDDqg> KS

5/20/2011



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[Beyond this lecture...](#)

20/04/20

Syntax of Eiffel: a Brief Overview



EECS3311 A: Software Design
Winter 2020

[CHEN-WEI WANG](#)

Escape Sequences



Escape sequences are special characters to be placed in your program text.

- In Java, an escape sequence starts with a backward slash \
e.g., \n for a new line character.
- In Eiffel, an escape sequence starts with a percentage sign %
e.g., %N for a new line character.

See here for more escape sequences in Eiffel: https://www.eiffel.org/doc/eiffel/Eiffel%20programming%20language%20syntax#Special_characters

20/04/20

Commands, and Queries, and Features



- In a Java class:
 - **Attributes:** Data
 - **Mutators:** Methods that change attributes without returning
 - **Accessors:** Methods that access attribute values and returning
- In an Eiffel class:
 - Everything can be called a *feature*.
 - But if you want to be specific:
 - Use *attributes* for data
 - Use *commands* for mutators
 - Use *queries* for accessors

20/04/20

Naming Conventions



- Cluster names: all lower-cases separated by underscores
e.g., root, model, tests, cluster_number_one
- Classes/Type names: all upper-cases separated by underscores
e.g., ACCOUNT, BANK_ACCOUNT_APPLICATION
- Feature names (attributes, commands, and queries): all lower-cases separated by underscores
e.g., account_balance, deposit_into, withdraw_from

make

Class Declarations



- In Java:

```
class BankAccount {  
    /* attributes and methods */  
}
```

- In Eiffel:

```
class BANK_ACCOUNT  
    /* attributes, commands, and queries */  
end
```

make

Class Constructor Declarations (1)



- In Eiffel, constructors are just commands that have been *explicitly* declared as **creation features**:

```
class BANK_ACCOUNT  
    -- List names commands that can be used as constructors  
    create  
        make  
    feature -- Commands  
        make (b: INTEGER)  
            do balance := b end  
        make2  
            do balance := 10 end  
    end
```

- Only the command `make` can be used as a constructor.
- Command `make2` is not declared explicitly, so it cannot be used as a constructor.

make

Creations of Objects (1)



- In Java, we use a constructor `Account (int b)` by:
 - Writing `Account acc = new Account (10)` to create a named object `acc`
 - Writing `new Account (10)` to create an anonymous object
- In Eiffel, we use a creation feature (i.e., a command explicitly declared under `create`) `make (int b)` in class `ACCOUNT` by:

- Writing `create {ACCOUNT} acc.make (10)` to create a named object `acc`
- Writing `create {ACCOUNT}.make (10)` to create an anonymous object

- Writing `create {ACCOUNT} acc.make (10)`

is really equivalent to writing

```
acc := create {ACCOUNT}.make (10)
```

make

Attribute Declarations



- In Java, you write: `int i, Account acc`
- In Eiffel, you write: `i: INTEGER, acc: ACCOUNT`
Think of `:` as the set membership operator \in :
e.g., The declaration `acc: ACCOUNT` means object `acc` is a member of all possible instances of `ACCOUNT`.

10/14/18

Operators: Assignment vs. Equality



- In Java:
 - Equal sign `=` is for assigning a value expression to some variable.
e.g., `x = 5 * y` changes `x`'s value to `5 * y`
This is actually controversial, since when we first learned about `=`, it means the mathematical equality between numbers.
 - Equal-equal `==` and bang-equal `!=` are used to denote the equality and inequality.
e.g., `x == 5 * y` evaluates to `true` if `x`'s value is equal to the value of `5 * y`, or otherwise it evaluates to `false`.
- In Eiffel:
 - Equal `=` and slash equal `/=` denote equality and inequality.
e.g., `x = 5 * y` evaluates to `true` if `x`'s value is equal to the value of `5 * y`, or otherwise it evaluates to `false`.
 - We use `:=` to denote variable assignment.
e.g., `x := 5 * y` changes `x`'s value to `5 * y`
 - Also, you are not allowed to write shorthands like `x++`, just write `x := x + 1`.

10/14/18

Method Declaration



• Command

```
deposit (amount: INTEGER)
do
  balance := balance + amount
end
```

Notice that you don't use the return type `void`

• Query

```
sum_of (x: INTEGER; y: INTEGER): INTEGER
do
  Result := x + y
end
```

- Input parameters are separated by semicolons `;`
- Notice that you don't use `return`; instead assign the return value to the pre-defined variable **Result**.

10/14/18

Operators: Division and Modulo



| | Division | Modulo (Remainder) |
|--------|---------------------------|---------------------------|
| Java | <code>20 / 3</code> is 6 | <code>20 % 3</code> is 2 |
| Eiffel | <code>20 // 3</code> is 6 | <code>20 \% 3</code> is 2 |

10/14/18

Operators: Logical Operators (1)



- Logical operators (what you learned from EECS1090) are for combining Boolean expressions.
- In Eiffel, we have operators that **EXACTLY** correspond to these logical operators:

| | LOGIC | EIFFEL |
|-------------|---------------|----------------|
| Conjunction | \wedge | and |
| Disjunction | \vee | or |
| Implication | \Rightarrow | implies |
| Equivalence | \equiv | = |

▶

Selections (1)



```

if B1 then
  -- B1
  -- do something
elseif B2 then
  -- B2  $\wedge$  ( $\neg$ B1)
  -- do something else
else
  -- ( $\neg$ B1)  $\wedge$  ( $\neg$ B2)
  -- default action
end
    
```

▶

Operators: Logical Operators (2)



- How about Java?
 - Java does not have an operator for logical implication.
 - The == operator can be used for logical equivalence.
 - The && and || operators only **approximate** conjunction and disjunction, due to the **short-circuit effect (SCE)**:
 - When evaluating e1 && e2, if e1 already evaluates to **false**, then e2 will **not** be evaluated.

e.g., In (y != 0) && (x / y > 10), the SCE guards the division against division-by-zero error.
 - When evaluating e1 || e2, if e1 already evaluates to **true**, then e2 will **not** be evaluated.

e.g., In (y == 0) || (x / y > 10), the SCE guards the division against division-by-zero error.
 - However, in math, the order of the two sides should not matter.
- In Eiffel, we also have the version of operators with SCE:

| | short-circuit conjunction | short-circuit disjunction |
|--------|---------------------------|---------------------------|
| Java | && | |
| Eiffel | and then | or else |

▶

Selections (2)



- An **if-statement** is considered as:
- An **instruction** if its branches contain **instructions**.
 - An **expression** if its branches contain Boolean **expressions**.

```

class
  FOO
  feature --Attributes
    x, y: INTEGER
  feature -- Commands
    command
      -- A command with if-statements in implementation and contracts.
      require
        if x \ 2 /= 0 then True else False end -- Or: x \ 2 /= 0
      do
        if x > 0 then y := 1 elseif x < 0 then y := -1 else y := 0 end
      ensure
        y = if old x > 0 then 1 elseif old x < 0 then -1 else 0 end
        -- Or: (old x > 0 implies y = 1)
        -- and (old x < 0 implies y = -1) and (old x = 0 implies y = 0)
      end
  end
end
    
```

▶

Loops (1)

- In Java, the Boolean conditions in `for` and `while` loops are **stay** conditions.

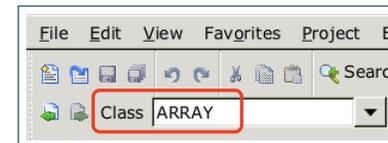
```
void printStuffs() {
    int i = 0;
    while(i < 10 /* stay condition */) {
        System.out.println(i);
        i = i + 1;
    }
}
```

- In the above Java loop, we **stay** in the loop as long as `i < 10` is true.
- In Eiffel, we think the opposite: we **exit** the loop as soon as `i >= 10` is true.

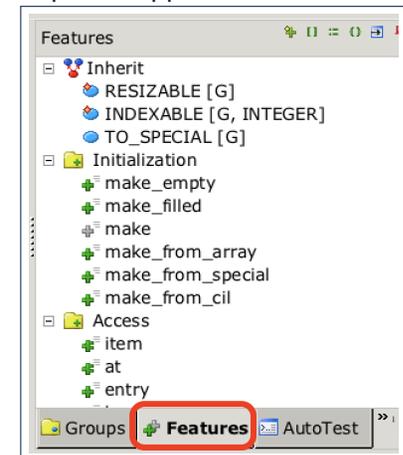
1/2/2008

Library Data Structures

Enter a DS name.



Explore supported features.



1/2/2008

Loops (2)

In Eiffel, the Boolean conditions you need to specify for loops are **exit** conditions (logical negations of the stay conditions).

```
print_stuffs
local
  i: INTEGER
do
  from
    i := 0
  until
    i >= 10 -- exit condition
  loop
    print (i)
    i := i + 1
  end -- end loop
end -- end command
```

- Don't put `()` after a command or query with no input parameters.
- Local variables must all be declared in the beginning.

1/2/2008

Data Structures: Arrays

- Creating an empty array:

```
local a: ARRAY[INTEGER]
do create {ARRAY[INTEGER]} a.make_empty
```

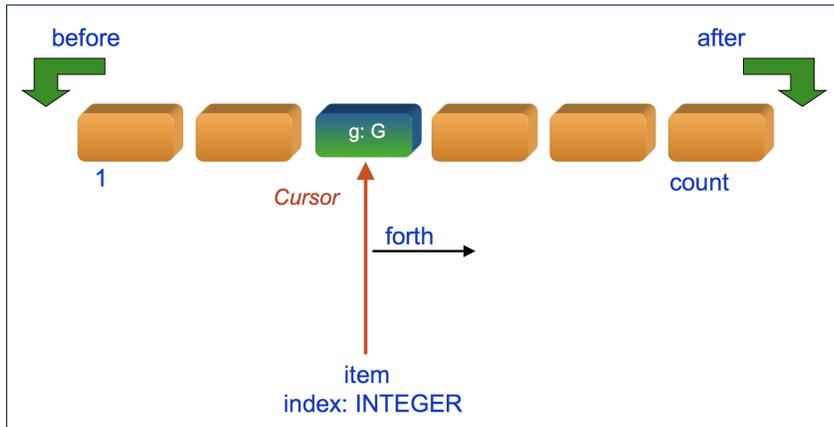
- This creates an array of lower and upper indices 1 and 0.
- Size of array a: `a.upper - a.lower + 1`.

- Typical loop structure to iterate through an array:

```
local
  a: ARRAY[INTEGER]
  i, j: INTEGER
do
  ...
  from
    j := a.lower
  until
    j > a.upper
  do
    i := a [j]
    j := j + 1
  end
```

1/2/2008

Data Structures: Linked Lists (1)



28/04/20

Iterable Structures



- Eiffel collection types (like in Java) are **iterable**.
- If indices are irrelevant for your application, use:

```
across ... as ... loop ... end
```

e.g.,

```
...
local
  a: ARRAY[INTEGER]
  l: LINKED_LIST[INTEGER]
  sum1, sum2: INTEGER
do
  ...
  across a as cursor loop sum1 := sum1 + cursor.item end
  across l as cursor loop sum2 := sum2 + cursor.item end
  ...
end
```

28/04/20

Data Structures: Linked Lists (2)



- Creating an empty linked list:

```
local
  list: LINKED_LIST[INTEGER]
do
  create {LINKED_LIST[INTEGER]} list.make
```

- Typical loop structure to iterate through a linked list:

```
local
  list: LINKED_LIST[INTEGER]
  i: INTEGER
do
  ...
  from
    list.start
  until
    list.after
  do
    i := list.item
    list.forth
  end
```

28/04/20

Using across for Quantifications (1.1)



- **across** ... **as** ... **all** ... **end**

A Boolean expression acting as a universal quantification (\forall)

```
1 local
2   allPositive: BOOLEAN
3   a: ARRAY[INTEGER]
4 do
5   ...
6   Result :=
7     across
8       a.lower |..| a.upper as i
9     all
10      a [i.item] > 0
11   end
```

- **L8**: a.lower |..| a.upper denotes a list of integers.
- **L8**: as i declares a list cursor for this list.
- **L10**: i.item denotes the value pointed to by cursor i.
- **L9**: Changing the keyword **all** to **some** makes it act like an existential quantification \exists .

28/04/20

Using across for Quantifications (1.2)



- Alternatively: **across** ... **is** ... **all** ... **end**
A Boolean expression acting as a universal quantification (\forall)

```
1 local
2   allPositive: BOOLEAN
3   a: ARRAY[INTEGER]
4 do
5   ...
6   Result :=
7     across
8     a.lower |..| a.upper is i
9     all
10    a [i] > 0
11  end
```

- L8**: `a.lower |..| a.upper` denotes a list of integers.
- L8**: `is i` declares a variable for storing a member of the list.
- L10**: `i` denotes the value itself.
- L9**: Changing the keyword **all** to **some** makes it act like an existential quantification \exists .

25/01/88

Using across for Quantifications (3)



```
class BANK
...
accounts: LIST [ACCOUNT]
binary_search (acc_id: INTEGER): ACCOUNT
  -- Search on accounts sorted in non-descending order.
  require
    --  $\forall i: \text{INTEGER} \mid 1 \leq i < \text{accounts.count} \bullet \text{accounts}[i].\text{id} \leq \text{accounts}[i+1].\text{id}$ 
    across
      1 |..| (accounts.count - 1) as cursor
      all
        accounts [cursor.item].id <= accounts [cursor.item + 1].id
      end
    do
      ...
    ensure
      Result.id = acc_id
    end
```

25/01/88

Using across for Quantifications (2)



```
class
  CHECKER
  feature -- Attributes
    collection: ITERABLE [INTEGER] -- ARRAY, LIST, HASH_TABLE
  feature -- Queries
    is_all_positive: BOOLEAN
    -- Are all items in collection positive?
    do
      ...
    ensure
      across
        collection as cursor
        all
          cursor.item > 0
        end
      end
```

- Using **all** corresponds to a universal quantification (i.e., \forall).
- Using **some** corresponds to an existential quantification (i.e., \exists).

25/01/88

Using across for Quantifications (4)



```
class BANK
...
accounts: LIST [ACCOUNT]
contains_duplicate: BOOLEAN
  -- Does the account list contain duplicate?
  do
    ...
  ensure
     $\forall i, j: \text{INTEGER} \mid 1 \leq i \leq \text{accounts.count} \wedge 1 \leq j \leq \text{accounts.count} \bullet \text{accounts}[i] \sim \text{accounts}[j] \Rightarrow i = j$ 
  end
```

- Exercise**: Convert this mathematical predicate for postcondition into Eiffel.
- Hint**: Each **across** construct can only introduce one dummy variable, but you may nest as many **across** constructs as necessary.

25/01/88

Equality



- To compare references between two objects, use `=`.
- To compare “contents” between two objects *of the same type*, use the *redefined* version of `is_equal` feature.
- You may also use the binary operator `~`
 - `o1 ~ o2` evaluates to:
 - true* if both `o1` and `o2` are void
 - false* if one is void but not the other
 - `o1.is_equal(o2)` if both are not void

28/04/20

Use of `~`: Caution



```
1 class
2   BANK
3   feature -- Attribute
4     accounts: ARRAY[ACCOUNT]
5   feature -- Queries
6     get_account (id: STRING): detachable ACCOUNT
7       -- Account object with 'id'.
8       do
9         across
10          accounts as cursor
11          loop
12            if cursor.item ~ id then
13              Result := cursor.item
14            end
15          end
16        end
17      end
```

L15 should be: `cursor.item.id ~ id`

28/04/20

Review of Propositional Logic (1)



- A *proposition* is a statement of claim that must be of either *true* or *false*, but not both.
- Basic logical operands are of type Boolean: *true* and *false*.
- We use logical operators to construct compound statements.
 - Binary logical operators: conjunction (\wedge), disjunction (\vee), implication (\Rightarrow), and equivalence (a.k.a if-and-only-if \Leftrightarrow)

| p | q | $p \wedge q$ | $p \vee q$ | $p \Rightarrow q$ | $p \Leftrightarrow q$ |
|--------------|--------------|--------------|--------------|-------------------|-----------------------|
| <i>true</i> | <i>true</i> | <i>true</i> | <i>true</i> | <i>true</i> | <i>true</i> |
| <i>true</i> | <i>false</i> | <i>false</i> | <i>true</i> | <i>false</i> | <i>false</i> |
| <i>false</i> | <i>true</i> | <i>false</i> | <i>true</i> | <i>true</i> | <i>false</i> |
| <i>false</i> | <i>false</i> | <i>false</i> | <i>false</i> | <i>true</i> | <i>true</i> |

- Unary logical operator: negation (\neg)

| p | $\neg p$ |
|--------------|--------------|
| <i>true</i> | <i>false</i> |
| <i>false</i> | <i>true</i> |

28/04/20

Review of Propositional Logic: Implication



- Written as $p \Rightarrow q$
- Pronounced as “ p implies q ”
- We call p the antecedent, assumption, or premise.
- We call q the consequence or conclusion.
- Compare the *truth* of $p \Rightarrow q$ to whether a contract is *honoured*: $p \approx$ promised terms; and $q \approx$ obligations.
- When the promised terms are met, then:
 - The contract is *honoured* if the obligations are fulfilled.
 - The contract is *breached* if the obligations are not fulfilled.
- When the promised terms are not met, then:
 - Fulfilling the obligation (q) or not ($\neg q$) does *not breach* the contract.

| p | q | $p \Rightarrow q$ |
|--------------|--------------|-------------------|
| <i>true</i> | <i>true</i> | <i>true</i> |
| <i>true</i> | <i>false</i> | <i>false</i> |
| <i>false</i> | <i>true</i> | <i>true</i> |
| <i>false</i> | <i>false</i> | <i>true</i> |

28/04/20



Review of Propositional Logic (2)

- **Axiom:** Definition of \Rightarrow

$$p \Rightarrow q \equiv \neg p \vee q$$

- **Theorem:** Identity of \Rightarrow

$$\text{true} \Rightarrow p \equiv p$$

- **Theorem:** Zero of \Rightarrow

$$\text{false} \Rightarrow p \equiv \text{true}$$

- **Axiom:** De Morgan

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

- **Axiom:** Double Negation

$$p \equiv \neg(\neg p)$$

- **Theorem:** Contrapositive

$$p \Rightarrow q \equiv \neg q \Rightarrow \neg p$$

EXAMPLE



Review of Predicate Logic (2.1)

- A **universal quantification** has the form $(\forall X | R \bullet P)$
 - X is a list of variable *declarations*
 - R is a *constraint on ranges* of declared variables
 - P is a *property*
 - $(\forall X | R \bullet P) \equiv (\forall X \bullet R \Rightarrow P)$
 - e.g., $(\forall X | \text{True} \bullet P) \equiv (\forall X \bullet \text{True} \Rightarrow P) \equiv (\forall X \bullet P)$
 - e.g., $(\forall X | \text{False} \bullet P) \equiv (\forall X \bullet \text{False} \Rightarrow P) \equiv (\forall X \bullet \text{True}) \equiv \text{True}$
- **For all** (combinations of) values of variables declared in X that satisfies R , it is the case that P is satisfied.
 - $\forall i | i \in \mathbb{N} \bullet i \geq 0$ [true]
 - $\forall i | i \in \mathbb{Z} \bullet i \geq 0$ [false]
 - $\forall i, j | i \in \mathbb{Z} \wedge j \in \mathbb{Z} \bullet i < j \vee i > j$ [false]
- The range constraint of a variable may be moved to where the variable is declared.
 - $\forall i : \mathbb{N} \bullet i \geq 0$
 - $\forall i : \mathbb{Z} \bullet i \geq 0$
 - $\forall i, j : \mathbb{Z} \bullet i < j \vee i > j$

EXAMPLE



Review of Predicate Logic (1)

- A **predicate** is a *universal* or *existential* statement about objects in some universe of disclosure.
- Unlike propositions, predicates are typically specified using *variables*, each of which declared with some *range* of values.
- We use the following symbols for common numerical ranges:
 - \mathbb{Z} : the set of integers
 - \mathbb{N} : the set of natural numbers
- Variable(s) in a predicate may be *quantified*:
 - **Universal quantification**:
All values that a variable may take satisfy certain property.
e.g., Given that i is a natural number, i is *always* non-negative.
 - **Existential quantification**:
Some value that a variable may take satisfies certain property.
e.g., Given that i is an integer, i *can be* negative.

EXAMPLE



Review of Predicate Logic (2.2)

- An **existential quantification** has the form $(\exists X | R \bullet P)$
 - X is a list of variable *declarations*
 - R is a *constraint on ranges* of declared variables
 - P is a *property*
 - $(\exists X | R \bullet P) \equiv (\exists X \bullet R \wedge P)$
 - e.g., $(\exists X | \text{True} \bullet P) \equiv (\exists X \bullet \text{True} \wedge P) \equiv (\exists X \bullet P)$
 - e.g., $(\exists X | \text{False} \bullet P) \equiv (\exists X \bullet \text{False} \wedge P) \equiv (\exists X \bullet \text{False}) \equiv \text{False}$
- **There exists** a combination of values of variables declared in X that satisfies R and P .
 - $\exists i | i \in \mathbb{N} \bullet i \geq 0$ [true]
 - $\exists i | i \in \mathbb{Z} \bullet i \geq 0$ [true]
 - $\exists i, j | i \in \mathbb{Z} \wedge j \in \mathbb{Z} \bullet i < j \vee i > j$ [true]
- The range constraint of a variable may be moved to where the variable is declared.
 - $\exists i : \mathbb{N} \bullet i \geq 0$
 - $\exists i : \mathbb{Z} \bullet i \geq 0$
 - $\exists i, j : \mathbb{Z} \bullet i < j \vee i > j$

EXAMPLE

Predicate Logic (3)



- Conversion between \forall and \exists

$$\begin{aligned}(\forall X \mid R \bullet P) &\iff \neg(\exists X \bullet R \Rightarrow \neg P) \\(\exists X \mid R \bullet P) &\iff \neg(\forall X \bullet R \Rightarrow \neg P)\end{aligned}$$

- Range Elimination

$$\begin{aligned}(\forall X \mid R \bullet P) &\iff (\forall X \bullet R \Rightarrow P) \\(\exists X \mid R \bullet P) &\iff (\exists X \bullet R \wedge P)\end{aligned}$$

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Common Eiffel Errors: Contracts vs. Implementations



EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG

Contracts vs. Implementations: Where?



- Instructions for *Implementations*: $inst_1, inst_2$
- Boolean expressions for Contracts: $exp_1, exp_2, exp_3, exp_4, exp_5$

```
class
  ACCOUNT
  feature -- Queries
    balance: INTEGER
    require
      exp1
    do
      inst1
    ensure
      exp2
    end
```

```
feature -- Commands
  withdraw
    require
      exp3
    do
      inst2
    ensure
      exp4
    end
  invariant
    exp5
end -- end of class ACCOUNT
```

2020

Contracts vs. Implementations: Definitions



In Eiffel, there are two categories of constructs:

- **Implementations**
 - are step-by-step **instructions** that have *side-effects*

e.g., `... := ...`, `across ... as ... loop ... end`

- change attribute values
- do not return values
- \approx commands
- **Contracts**
 - are Boolean **expressions** that have *no side-effects*

e.g., `... = ...`, `across ... as ... all ... end`

- use attribute and parameter values to specify a condition
- return a Boolean value (i.e., *True* or *False*)
- \approx queries

2020

Implementations: Instructions with No Return Values



- Assignments

```
balance := balance + a
```

- Selections with branching instructions:

```
if a > 0 then acc.deposit (a) else acc.withdraw (-a) end
```

- Loops

```
from
  i := a.lower
until
  i > a.upper
loop
  Result :=
    Result + a[i]
  i := i + 1
end
```

```
from
  list.start
until
  list.after
loop
  list.item.wdw(10)
  list.forth
end
```

```
across
  list as cursor
loop
  sum :=
    sum + cursor.item
end
```

2020



Contracts: Expressions with Boolean Return Values

- Relational Expressions (using =, /=, ~, /~, >, <, >=, <=)

```
a > 0
```

- Binary Logical Expressions (using **and**, **and then**, **or**, **or else**, **implies**)

```
(a.lower <= index) and (index <= a.upper)
```

- Logical Quantification Expressions (using **all**, **some**)

```
across
  a.lower |..| a.upper as cursor
all
  a [cursor.item] >= 0
end
```

- old** keyword can only appear in postconditions (i.e., **ensure**).

```
balance = old balance + a
```

not



Contracts: Common Mistake (1) Fixed

```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
  do
    ...
  ensure
    balance = old balance - a
  end
...

```

not



Contracts: Common Mistake (1)

```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
  do
    ...
  ensure
    balance := old balance - a
  end
...

```

Colon-Equal sign (: =) is used to write assignment instructions.

not



Contracts: Common Mistake (2)

```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
  do
    ...
  ensure
    across
      a as cursor
    loop
      ...
    end
  end
...

```

across ... loop ... end is used to create loop instructions.

not

Contracts: Common Mistake (2) Fixed



```
class
  ACCOUNT
  feature
    withdraw (a: INTEGER)
    do
      ...
    ensure
      across
        a as cursor
      all -- if you meant  $\forall$ , or use some if you meant  $\exists$ 
        ... -- A Boolean expression is expected here!
      end
    end
  ...
end
```

10/28

Contracts: Common Mistake (3) Fixed



```
class
  ACCOUNT
  feature
    withdraw (a: INTEGER)
    do
      ...
    ensure
      postcond_1: balance = old balance - a
      postcond_2: old balance > 0
    end
  ...
end
```

10/28

Contracts: Common Mistake (3)



```
class
  ACCOUNT
  feature
    withdraw (a: INTEGER)
    do
      ...
    ensure
      old balance - a
    end
  ...
end
```

Contracts can only be specified as Boolean expressions.

10/28

Contracts: Common Mistake (4)



```
class
  ACCOUNT
  feature
    withdraw (a: INTEGER)
    require
      old balance > 0
    do
      ...
    ensure
      ...
    end
  ...
end
```

- Only **postconditions** may use the **old** keyword to specify *the relationship between pre-state values* (before the execution of *withdraw*) *and post-state values* (after the execution of *withdraw*).
- **Pre-state values** (right before the feature is executed) are **considered the old values**, so there's no need to qualify them!

10/28

Contracts: Common Mistake (4) Fixed



```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
    require
      balance > 0
    do
      ...
    ensure
      ...
    end
  ...
end
...
```

ERROR

Contracts: Common Mistake (5) Fixed



```
class LINEAR_CONTAINER
create make
feature -- Attributes
  a: ARRAY[STRING]
feature -- Queries
  count: INTEGER do Result := a.count end
  get (i: INTEGER): STRING do Result := a[i] end
feature -- Commands
  make do create a.make_empty end
  update (i: INTEGER; v: STRING)
  do ...
  ensure -- Others Unchanged
    across
      1 |..| count as j
    all
      j.item /= i implies (old Current).get(j.item) ~ get(j.item)
    end
  end
end
end
```

- o The idea is that the **old** expression should not involve the local cursor variable *j* that is introduced in the postcondition.
- o Whether to put (old *Current.twin*) or (old *Current.deep_twin*) is up to your need.

ERROR

Contracts: Common Mistake (5)



```
class LINEAR_CONTAINER
create make
feature -- Attributes
  a: ARRAY[STRING]
feature -- Queries
  count: INTEGER do Result := a.count end
  get (i: INTEGER): STRING do Result := a[i] end
feature -- Commands
  make do create a.make_empty end
  update (i: INTEGER; v: STRING)
  do ...
  ensure -- Others Unchanged
    across
      1 |..| count as j
    all
      j.item /= i implies old get(j.item) ~ get(j.item)
    end
  end
end
end
```

Compilation Error:

- o Expression value to be cached before executing update?
[*Current.get(j.item)*]
- o But, in the **pre-state**, integer cursor *j* does not exist!

ERROR

Implementations: Common Mistake (1)



```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
    do
      balance = balance + 1
    end
  ...
end
```

- Equal sign (=) is used to write Boolean expressions.
- In the context of implementations, Boolean expression values must appear:
 - o on the RHS of an **assignment**;
 - o as one of the **branching conditions** of an if-then-else statement; or
 - o as the **exit condition** of a loop instruction.

ERROR

Implementations: Common Mistake (1) Fixed



```
class
  ACCOUNT
feature
  withdraw (a: INTEGER)
  do
    balance := balance + 1
  end
  ...
```

2012

Implementations: Common Mistake (2) Fixed



```
1 class
2   BANK
3 feature
4   min_credit: REAL
5   accounts: LIST[ACCOUNT]
6
7   no_warning_accounts: BOOLEAN
8   do
9     Result :=
10      across
11        accounts as cursor
12      all
13        cursor.item.balance > min_credit
14      end
15    end
16  ...
```

Rewrite L10 – L14 using `across ... as ... some ... end`.

Hint: $\forall x \bullet P(x) \equiv \neg(\exists x \bullet \neg P(x))$

2012

Implementations: Common Mistake (2)



```
class
  BANK
feature
  min_credit: REAL
  accounts: LIST[ACCOUNT]

  no_warning_accounts: BOOLEAN
  do
    across
      accounts as cursor
    all
      cursor.item.balance > min_credit
    end
  end
  ...
```

Again, in implementations, Boolean expressions cannot appear alone without their values being “captured”.

2012

Implementations: Common Mistake (3)



```
class
  BANK
feature
  accounts: LIST[ACCOUNT]

  total_balance: REAL
  do
    Result :=
      across
        accounts as cursor
      loop
        Result := Result + cursor.item.balance
      end
    ...
  end
  ...
```

In implementations, since instructions do not return values, they cannot be used on the RHS of assignments.

2012

Implementations: Common Mistake (3) Fixed



```
class
  BANK
feature
  accounts: LIST[ACCOUNT]

  total_balance: REAL
  do
    across
      accounts as cursor
    loop
      Result := Result + cursor.item.balance
    end
  end
end
```

2/20/20

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2/20/20

Drawing a Design Diagram using the Business Object Notation (BON)



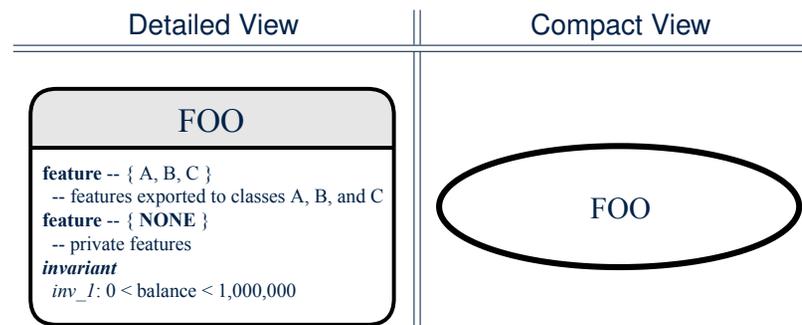
Why a Design Diagram?

- **SOURCE CODE** is **not** an appropriate form for communication.
- Use a **DESIGN DIAGRAM** showing **selective** sets of important:
 - clusters (i.e., packages)
 - classes [deferred vs. effective] [generic vs. non-generic]
 - architectural relations [client-supplier vs. inheritance]
 - features (queries and commands) [deferred vs. effective vs. redefined]
 - **contracts** [precondition vs. postcondition vs. class invariant]
- Your design diagram is called an **abstraction** of your system:
 - Being **selective** on what to show, filtering out **irrelevant details**
 - Presenting **contractual specification** in a **mathematical form** (e.g., \forall instead of **across ... all ... end**).

2012



Classes: Detailed View vs. Compact View (2)



2012



Classes: Detailed View vs. Compact View (1)

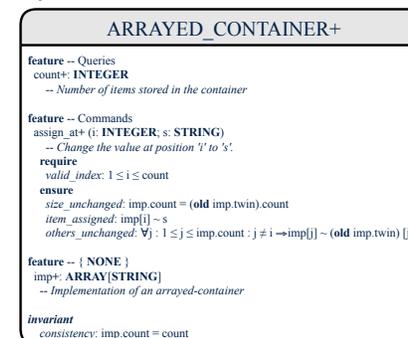
- **Detailed view** shows a selection of:
 - **features** (queries and/or commands)
 - **contracts** (class invariant and feature pre-post-conditions)
 - Use the detailed view if readers of your design diagram **should know** such details of a class. e.g., Classes critical to your design or implementation
- **Compact view** shows only the class name.
 - Use the compact view if readers **should not be bothered with** such details of a class. e.g., Minor "helper" classes of your design or implementation e.g., Library classes (e.g., ARRAY, LINKED_LIST, HASH_TABLE)

2012



Contracts: Mathematical vs. Programming

- When presenting the detailed view of a class, you should include **contracts** of features which you judge as **important**.
- Consider an array-based linear container:



- A **tag** should be included for each contract.
- Use **mathematical** symbols (e.g., \forall , \exists , \leq) instead of **programming** symbols (e.g., **across ... all ...**, **across ... some ...**, **<=**).

2012

Classes: Generic vs. Non-Generic



- A class is **generic** if it declares **at least one** type parameters.
 - Collection classes are generic: `ARRAY [G]`, `HASH_TABLE [G, H]`, *etc.*
 - Type parameter(s) of a class may or may not be **instantiated**:



- If necessary, present a generic class in the detailed form:



- A class is **non-generic** if it declares **no** type parameters.



Classes: Deferred vs. Effective



- A **deferred class** has **at least one** feature **unimplemented**.
 - A **deferred class** may only be used as a **static** type (for declaration), but cannot be used as a **dynamic** type.
 - e.g., By declaring `list: LIST [INTEGER]` (where `LIST` is a **deferred class**), it is invalid to write:
 - `create list.make`
 - `create {LIST [INTEGER]} list.make`
- An **effective class** has **all** features **implemented**.
 - An **effective class** may be used as both **static** and **dynamic** types.
 - e.g., By declaring `list: LIST [INTEGER]`, it is valid to write:
 - `create {LINKED_LIST [INTEGER]} list.make`
 - `create {ARRAYED_LIST [INTEGER]} list.make`where `LINKED_LIST` and `ARRAYED_LIST` are both **effective** descendants of `LIST`.



Deferred vs. Effective



Deferred means **unimplemented** (\approx **abstract** in Java)

Effective means **implemented**



Features: Deferred, Effective, Redefined (1)



- A **deferred feature** is declared with its **header** only (i.e., name, parameters, return type).
 - The word "**deferred**" means a descendant class would later implement this feature.
 - The resident class of the **deferred** feature must also be **deferred**.

```
deferred class
  DATABASE[G]
  feature -- Queries
  search (g: G): BOOLEAN
    -- Does item 'g' exist in database?
  deferred end
end
```



Features: Deferred, Effective, Redefined (2)



- An **effective feature** *implements* some inherited deferred feature.

```
class
  DATABASE_V1[G]
inherit
  DATABASE
feature -- Queries
  search (g: G): BOOLEAN
    -- Perform a linear search on the database.
  deferred end
end
```

- A descendant class may still later *re-implement* this feature.

13/02/25

Features: Deferred, Effective, Redefined (3)



- A **redefined feature** *re-implements* some inherited effective feature.

```
class
  DATABASE_V2[G]
inherit
  DATABASE_V1[G]
  redefine search end
feature -- Queries
  search (g: G): BOOLEAN
    -- Perform a binary search on the database.
  deferred end
end
```

- A descendant class may still later *re-implement* this feature.

13/02/25

Classes: Deferred vs. Effective (2.1)



- Append a star * to the name of a **deferred** class or feature.
- Append a plus + to the name of an **effective** class or feature.
- Append two pluses ++ to the name of a **redefined** feature.
- Deferred or effective classes may be in the compact form:

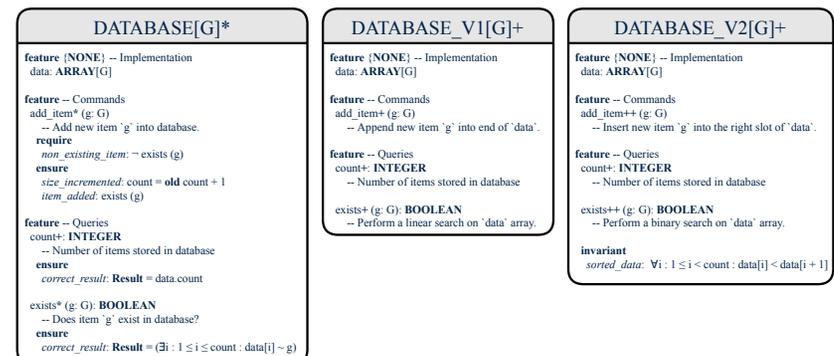


13/02/25

Classes: Deferred vs. Effective (2.2)



- Append a star * to the name of a **deferred** class or feature.
- Append a plus + to the name of an **effective** class or feature.
- Append two pluses ++ to the name of a **redefined** feature.
- Deferred or effective classes may be in the detailed form:



13/02/25

Class Relations: Inheritance (1)

- An **inheritance hierarchy** is formed using **red arrows**.
 - Arrow's **origin** indicates the **child/descendant** class.
 - Arrow's **destination** indicates the **parent/ancestor** class.
- You may choose to present each class in an inheritance hierarchy in either the **detailed** form or the **compact** form:

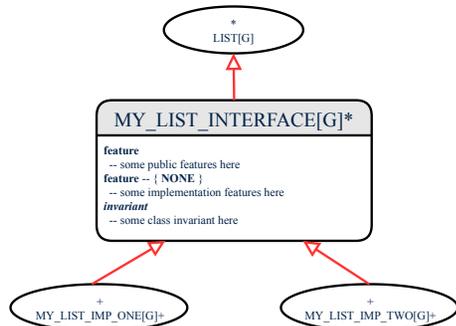


FIGURE 2.5

Class Relations: Client-Supplier (1)

- A **client-supplier (CS) relation** exists between two classes: one (the **client**) uses the service of another (the **supplier**).
- Programmatically, there is CS relation if in class CLIENT there is a **variable declaration** `s1: SUPPLIER`.
 - A variable may be an **attribute**, a **parameter**, or a **local variable**.
- A **green arrow** is drawn between the two classes.
 - Arrow's **origin** indicates the **client** class.
 - Arrow's **destination** indicates the **supplier** class.
 - Above the label there should be a **label** indicating the **supplier name** (i.e., variable name).
 - In the case where supplier is an **attribute**, indicate after the label name if it is deferred (*****), effective (**+**), or redefined (**++**).

FIGURE 2.5

Class Relations: Inheritance (2)

More examples (emphasizing different aspects of DATABASE):

Inheritance Hierarchy | Features being (Re-)Implemented

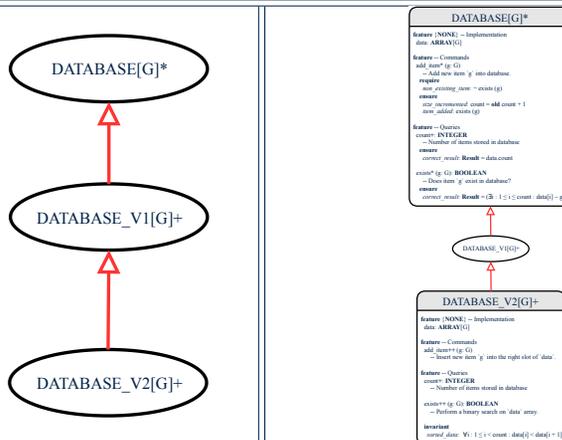


FIGURE 2.5

Class Relations: Client-Supplier (2.1)

```

class DATABASE
feature {NONE} -- implementation
data: ARRAY[STRING]
feature -- Commands
add_name (nn: STRING)
    -- Add name 'nn' to database.
    require ... do ... ensure ... end

name_exists (n: STRING): BOOLEAN
    -- Does name 'n' exist in database?
    require ...
    local
    u: UTILITIES
    do ... ensure ... end
invariant
...
end
    
```

```

class UTILITIES
feature -- Queries
search (a: ARRAY[STRING]; n: STRING): BOOLEAN
    -- Does name 'n' exist in array 'a'?
    require ... do ... ensure ... end
end
    
```

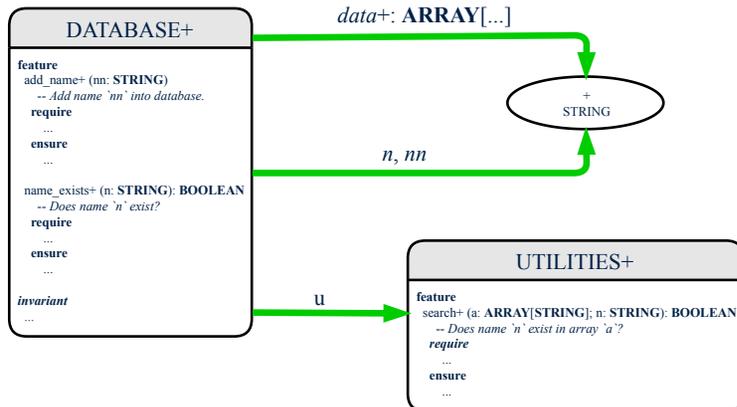
- Attribute `data: ARRAY[STRING]` indicates two suppliers: **STRING** and **ARRAY**.
- Parameters `nn` and `n` may have an arrow with label `nn, n`, pointing to the **STRING** class.
- Local variable `u` may have an arrow with label `u`, pointing to the **UTILITIES** class.

FIGURE 2.5

Class Relations: Client-Supplier (2.2.1)



If STRING is to be emphasized, label is `data: ARRAY[...]`, where ... denotes the supplier class STRING being pointed to.



2.2.2.1

Class Relations: Client-Supplier (3.1)



Known: The *deferred* class LIST has two *effective* descendants ARRAY_LIST and LINKED_LIST).

• DESIGN ONE:

```
class DATABASE_V1
feature {NONE} -- implementation
imp: ARRAYED_LIST[PERSON]
... -- more features and contracts
end
```

• DESIGN TWO:

```
class DATABASE_V2
feature {NONE} -- implementation
imp: LIST[PERSON]
... -- more features and contracts
end
```

Question: Which design is better? [DESIGN TWO]

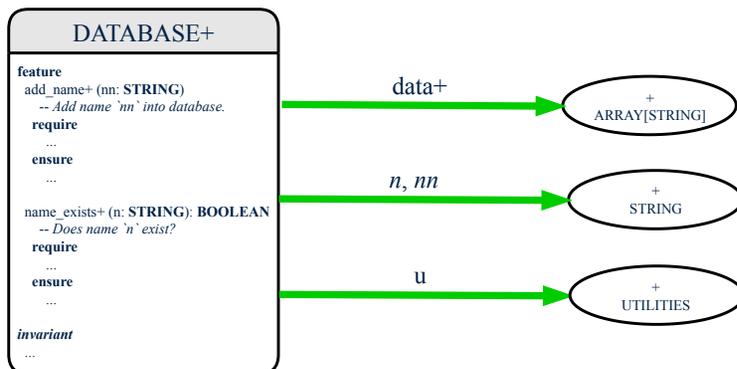
Rationale: Program to the *interface*, not the *implementation*.

3.1.1

Class Relations: Client-Supplier (2.2.2)



If ARRAY is to be emphasized, label is `data`.
The supplier's name should be complete: `ARRAY [STRING]`

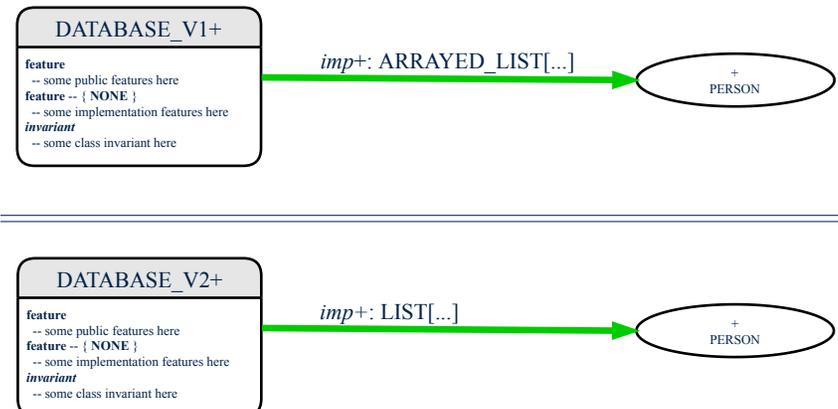


2.2.2.2

Class Relations: Client-Supplier (3.2.1)



We may focus on the PERSON supplier class, which may not help judge which design is better.

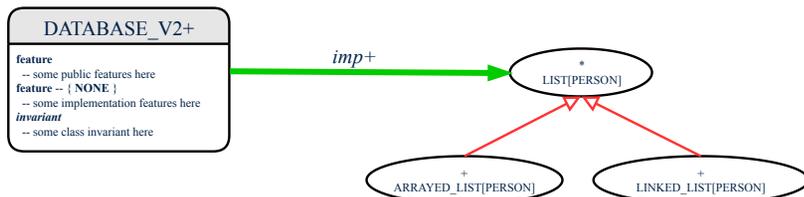
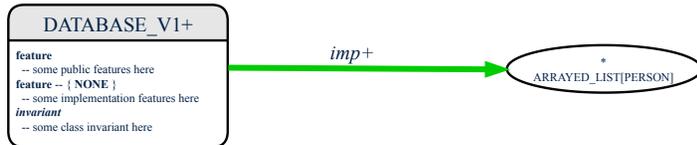


3.2.1

Class Relations: Client-Supplier (3.2.2)



Alternatively, we may focus on the LIST supplier class, which in this case helps us judge which design is better.

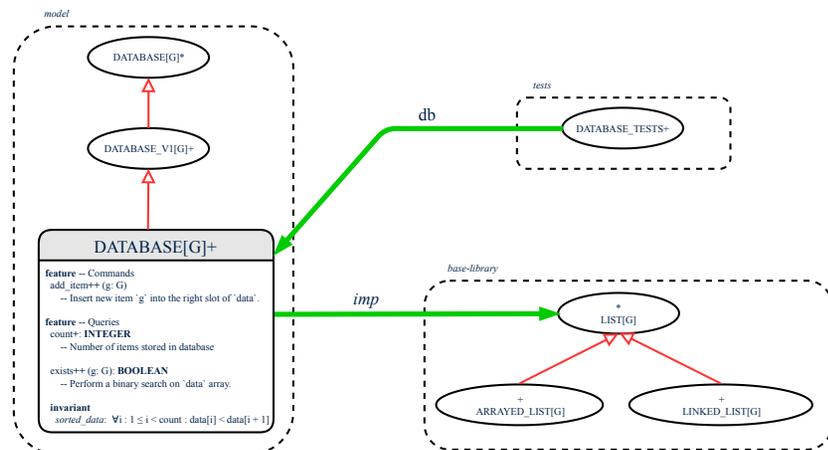


25/01/25

Clusters: Grouping Classes



Use **clusters** to group classes into logical units.



25/01/25

Index (1)



Why a Design Diagram?

Classes:

Detailed View vs. Compact View (1)

Classes:

Detailed View vs. Compact View (2)

Contracts: Mathematical vs. Programming

Classes: Generic vs. Non-Generic

Deferred vs. Effective

Classes: Deferred vs. Effective

Features: Deferred, Effective, Redefined (1)

Features: Deferred, Effective, Redefined (2)

Features: Deferred, Effective, Redefined (3)

Classes: Deferred vs. Effective (2.1)

Classes: Deferred vs. Effective (2.2)

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Index (2)



Class Relations: Inheritance (1)

Class Relations: Inheritance (2)

Class Relations: Client-Supplier (1)

Class Relations: Client-Supplier (2.1)

Class Relations: Client-Supplier (2.2.1)

Class Relations: Client-Supplier (2.2.2)

Class Relations: Client-Supplier (3.1)

Class Relations: Client-Supplier (3.2.1)

Class Relations: Client-Supplier (3.2.2)

Clusters: Grouping Classes

25/01/25

Copies: Reference vs. Shallow vs. Deep Writing Complete Postconditions



EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG



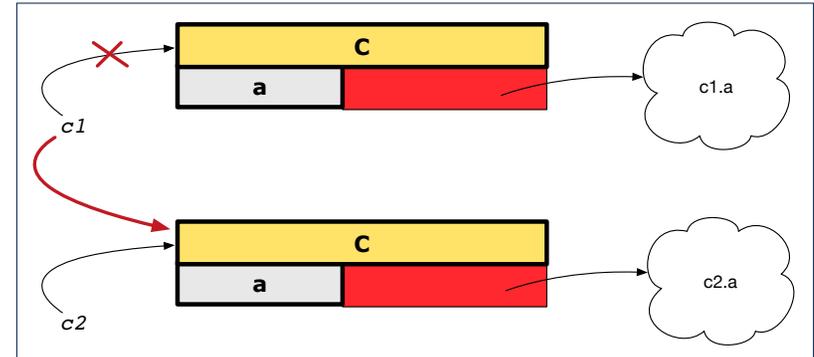
Copying Objects: Reference Copy

Reference Copy

`c1 := c2`

- Copy the address stored in variable `c2` and store it in `c1`.
 - ⇒ Both `c1` and `c2` point to the same object.
 - ⇒ Updates performed via `c1` also visible to `c2`.

[*aliasing*]



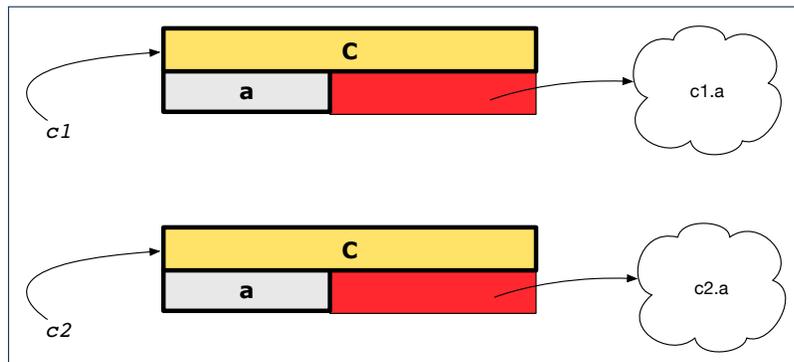
© 2014

Copying Objects



Say variables `c1` and `c2` are both declared of type `C`. [`c1, c2: C`]

- There is only one attribute `a` declared in class `C`.
- `c1.a` and `c2.a` are references to objects.



© 2014

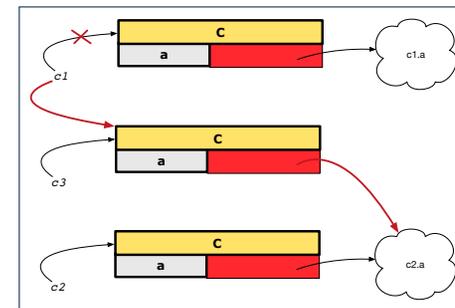
Copying Objects: Shallow Copy



Shallow Copy

`c1 := c2.twin`

- Create a temporary, behind-the-scene object `c3` of type `C`.
- Initialize each attribute `a` of `c3` via *reference copy*: `c3.a := c2.a`
- Make a *reference copy* of `c3`: `c1 := c3`
 - ⇒ `c1` and `c2` **are not** pointing to the same object. [`c1 != c2`]
 - ⇒ `c1.a` and `c2.a` **are** pointing to the same object.
 - ⇒ *Aliasing* still occurs: at 1st level (i.e., attributes of `c1` and `c2`)

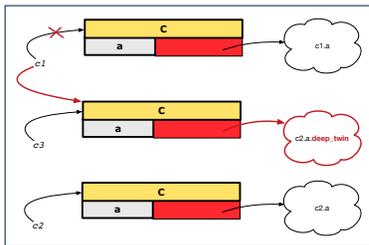


© 2014

Copying Objects: Deep Copy

Deep Copy

- Create a temporary, behind-the-scenes object `c3` of type `C`.
- **Recursively** initialize each attribute `a` of `c3` as follows:
 - Base Case:** `a` is primitive (e.g., `INTEGER`). $\Rightarrow c3.a := c2.a$.
 - Recursive Case:** `a` is referenced. $\Rightarrow c3.a := c2.a.deep_twin$
- Make a **reference copy** of `c3`: $c1 := c3$
 - $\Rightarrow c1$ and `c2` **are not** pointing to the same object.
 - $\Rightarrow c1.a$ and `c2.a` **are not** pointing to the same object.
 - \Rightarrow **No aliasing** occurs at any levels.



Example: Collection Objects (1)

- In any OOP, when a variable is declared of a **type** that corresponds to a **known class** (e.g., `STRING`, `ARRAY`, `LINKED_LIST`, etc.):
 - At **runtime**, that variable stores the **address** of an object of that type (as opposed to storing the object in its entirety).
- Assume the following variables of the same type:

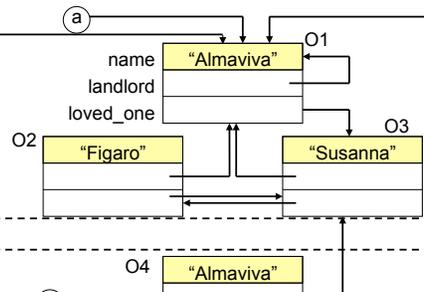
```

local
  imp : ARRAY[STRING]
  old_imp: ARRAY[STRING]
do
  create {ARRAY[STRING]} imp.make_empty
  imp.force("Alan", 1)
  imp.force("Mark", 2)
  imp.force("Tom", 3)
    
```

- **Before** we undergo a change on `imp`, we **copy** it to `old_imp`.
- **After** the change is completed, we compare `imp` vs. `old_imp`.
- Can a change always be **visible** between **"old"** and **"new"** `imp`?

Copying Objects

- Initial situation:

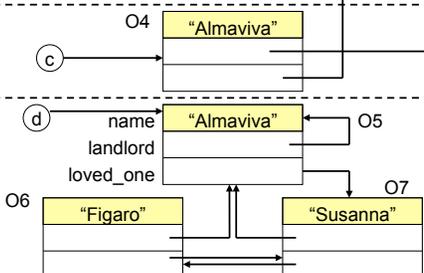


- Result of:

$b := a$

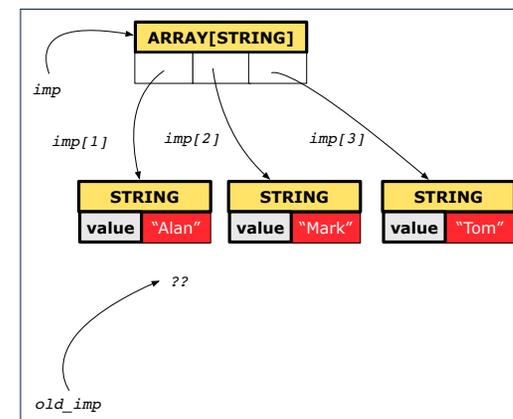
$c := a.twin$

$d := a.deep_twin$



Example: Collection Objects (2)

- Variables `imp` and `old_imp` store address(es) of some array(s).
- Each "slot" of these arrays stores a `STRING` object's address.



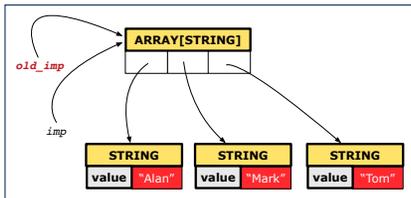
Reference Copy of Collection Object



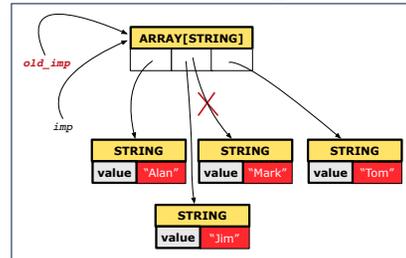
```

1 old_imp := imp
2 Result := old_imp = imp -- Result = true
3 imp[2] := "Jim"
4 Result :=
5   across 1 |..| imp.count is j
6   all imp [j] ~ old_imp [j]
7 end -- Result = true
    
```

Before Executing L3



After Executing L3



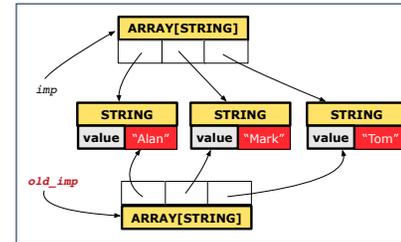
Shallow Copy of Collection Object (2)



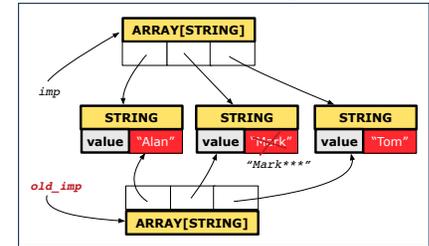
```

1 old_imp := imp.twin
2 Result := old_imp = imp -- Result = false
3 imp[2].append ("**")
4 Result :=
5   across 1 |..| imp.count is j
6   all imp [j] ~ old_imp [j]
7 end -- Result = true
    
```

Before Executing L3



After Executing L3



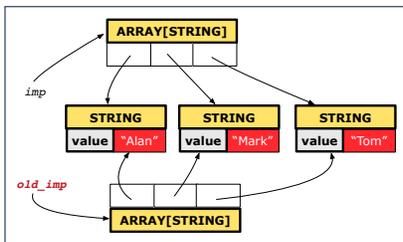
Shallow Copy of Collection Object (1)



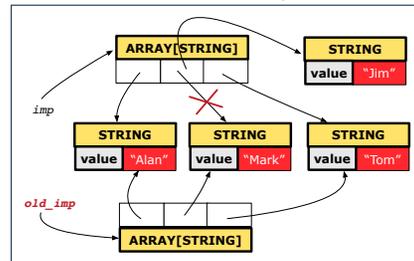
```

1 old_imp := imp.twin
2 Result := old_imp = imp -- Result = false
3 imp[2] := "Jim"
4 Result :=
5   across 1 |..| imp.count is j
6   all imp [j] ~ old_imp [j]
7 end -- Result = false
    
```

Before Executing L3



After Executing L3



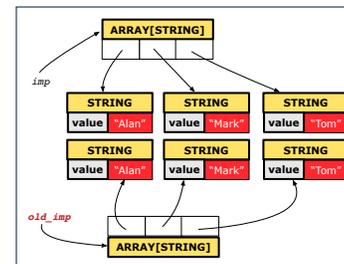
Deep Copy of Collection Object (1)



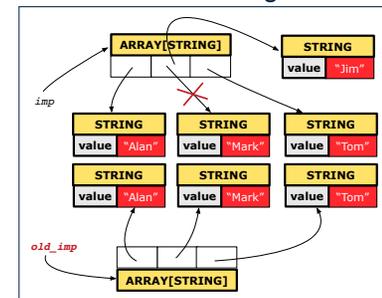
```

1 old_imp := imp.deep_twin
2 Result := old_imp = imp -- Result = false
3 imp[2] := "Jim"
4 Result :=
5   across 1 |..| imp.count is j
6   all imp [j] ~ old_imp [j] end -- Result = false
    
```

Before Executing L3



After Executing L3



Deep Copy of Collection Object (2)

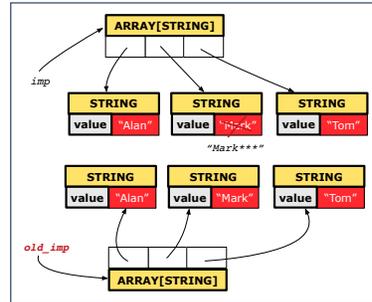
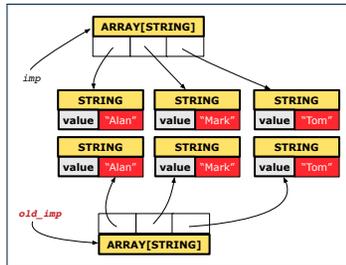
```

1  old_imp := imp.deep.twin
2  Result := old_imp = imp -- Result = false
3  imp[2].append ("**")
4  Result :=
5  across 1 |..| imp.count is j
6  all imp [j] ~ old_imp [j] end -- Result = false

```

Before Executing L3

After Executing L3



1/2/2018

When are contracts complete?

- In **post-condition**, for **each attribute**, specify the relationship between its **pre-state** value and its **post-state** value.
 - Eiffel supports this purpose using the **old** keyword.
- This is tricky for attributes whose structures are **composite** rather than **simple**:
 - e.g., *ARRAY*, *LINKED_LIST* are composite-structured.
 - e.g., *INTEGER*, *BOOLEAN* are simple-structured.
- Rule of thumb:** For an attribute whose structure is composite, we should specify that after the update:
 - The intended change is present; **and**
 - The rest of the structure is unchanged**.
- The second contract is much harder to specify:
 - Reference aliasing [ref copy vs. shallow copy vs. deep copy]
 - Iterable structure [use across]

1/2/2018

How are contracts checked at runtime?

- All contracts are specified as Boolean expressions.
- Right **before** a feature call (e.g., `acc.withdraw(10)`):
 - The current state of `acc` is called its **pre-state**.
 - Evaluate **pre-condition** using **current values** of attributes/queries.
 - Cache values, via `:=`, of **old expressions** in the **post-condition**.
 - e.g., `old accounts[i].id` [`old_accounts.i.id := accounts[i].id`]
 - e.g., `(old accounts[i]).id` [`old_accounts.i := accounts[i]`]
 - e.g., `(old accounts[i].twin).id` [`old_accounts.i.twin := accounts[i].twin`]
 - e.g., `(old accounts)[i].id` [`old_accounts := accounts`]
 - e.g., `(old accounts.twin)[i].id` [`old_accounts.twin := accounts.twin`]
 - e.g., `(old Current).accounts[i].id` [`old_current := Current`]
 - e.g., `(old Current.twin).accounts[i].id` [`old_current.twin := Current.twin`]
- Right **after** the feature call:
 - The current state of `acc` is called its **post-state**.
 - Evaluate **invariant** using **current values** of attributes and queries.
 - Evaluate **post-condition** using both **current values** and **"cached" values** of attributes and queries.

1/2/2018

Account

```

class
  ACCOUNT

inherit
  ANY
  redefine is_equal end

create
  make

feature -- Attributes
  owner: STRING
  balance: INTEGER

feature -- Commands
  make (n: STRING)
  do
    owner := n
    balance := 0
  end

```

```

deposit(a: INTEGER)
do
  balance := balance + a
ensure
  balance = old balance + a
end

is_equal(other: ACCOUNT): BOOLEAN
do
  Result :=
    owner ~ other.owner
  and balance = other.balance
end
end

```

1/2/2018



Bank

```

class BANK
create make
feature
  accounts: ARRAY[ACCOUNT]
  make do create accounts.make_empty end
  account_of (n: STRING): ACCOUNT
  require -- the input name exists
    existing: across accounts is acc some acc.owner ~ n end
    -- not (across accounts is acc all acc.owner /~ n end)
  do ... ensure Result.owner ~ n end
  add (n: STRING)
  require -- the input name does not exist
    non_existing: across accounts is acc all acc.owner /~ n end
    -- not (across accounts is acc some acc.owner ~ n end)
  local new_account: ACCOUNT
  do
    create new_account.make (n)
    accounts.force (new_account, accounts.upper + 1)
  end
end

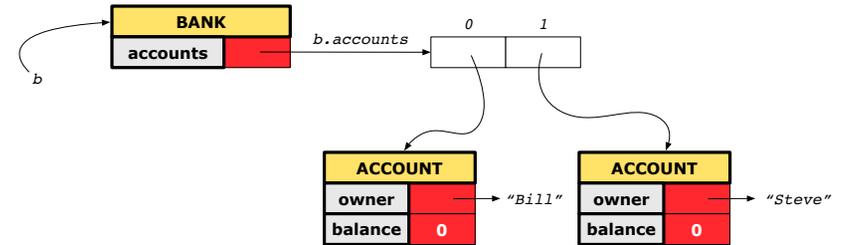
```

201808



Object Structure for Illustration

We will test each version by starting with the same runtime object structure:



201808



Roadmap of Illustrations

We examine 5 different versions of a command

deposit_on (n: STRING; a: INTEGER)

| VERSION | IMPLEMENTATION | CONTRACTS | SATISFACTORY? |
|---------|----------------|---------------------------|---------------|
| 1 | Correct | Incomplete | No |
| 2 | Wrong | Incomplete | No |
| 3 | Wrong | Complete (reference copy) | No |
| 4 | Wrong | Complete (shallow copy) | No |
| 5 | Wrong | Complete (deep copy) | Yes |

201808



Version 1: Incomplete Contracts, Correct Implementation

```

class BANK
  deposit_on_v1 (n: STRING; a: INTEGER)
  require across accounts is acc some acc.owner ~ n end
  local i: INTEGER
  do
    from i := accounts.lower
    until i > accounts.upper
    loop
      if accounts[i].owner ~ n then accounts[i].deposit(a) end
      i := i + 1
    end
  end
  ensure
    num_of_accounts_unchanged:
      accounts.count = old accounts.count
    balance_of_n_increased:
      Current.account_of(n).balance =
        old Current.account_of(n).balance + a
  end
end

```

201808

Test of Version 1



```

class TEST_BANK
  test_bank_deposit_correct_imp_incomplete_contract: BOOLEAN
  local
    b: BANK
  do
    comment("t1: correct imp and incomplete contract")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v1 ("Steve", 100)
    Result :=
      b.account_of("Bill").balance = 0
      and b.account_of("Steve").balance = 100
    check Result end
  end
end
  
```

2/2/2020

Version 2: Incomplete Contracts, Wrong Implementation



```

class BANK
  deposit_on_v2 (n: STRING; a: INTEGER)
  require across accounts is acc some acc.owner ~ n end
  local i: INTEGER
  do ...
    -- imp. of version 1, followed by a deposit into 1st account
    accounts[accounts.lower].deposit(a)
  ensure
    num_of_accounts_unchanged:
      accounts.count = old accounts.count
    balance_of_n_increased:
      Current.account_of(n).balance =
        old Current.account_of(n).balance + a
  end
end
  
```

Current postconditions lack a check that accounts other than n are unchanged.

2/2/2020

Test of Version 1: Result



APPLICATION

Note: * indicates a violation test case

| PASSED (1 out of 1) | | |
|---------------------|--------------------|--|
| Case Type | Passed | Total |
| Violation | 0 | 0 |
| Boolean | 1 | 1 |
| All Cases | 1 | 1 |
| State | Contract Violation | Test Name |
| Test1 | | TEST_BANK |
| PASSED | NONE | t1: test deposit_on with correct imp and incomplete contract |

2/2/2020

Test of Version 2



```

class TEST_BANK
  test_bank_deposit_wrong_imp_incomplete_contract: BOOLEAN
  local
    b: BANK
  do
    comment("t2: wrong imp and incomplete contract")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v2 ("Steve", 100)
    Result :=
      b.account_of("Bill").balance = 0
      and b.account_of("Steve").balance = 100
    check Result end
  end
end
  
```

2/2/2020

Test of Version 2: Result



APPLICATION

Note: * indicates a violation test case

| FAILED (1 failed & 1 passed out of 2) | | |
|---------------------------------------|---------------------------|--|
| Case Type | Passed | Total |
| Violation | 0 | 0 |
| Boolean | 1 | 2 |
| All Cases | 1 | 2 |
| State | Contract Violation | Test Name |
| Test1 | TEST_BANK | |
| PASSED | NONE | t1: test deposit_on with correct imp and incomplete contract |
| FAILED | Check assertion violated. | t2: test deposit_on with wrong imp but incomplete contract |

25/07/20

Test of Version 3



```
class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_ref_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t3: wrong imp and complete contract with ref copy")
    create b.make
    b.add("Bill")
    b.add("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on.v3("Steve", 100)
    Result :=
      b.account_of("Bill").balance = 0
      and b.account_of("Steve").balance = 100
    check Result end
  end
end
```

25/07/20

Version 3: Complete Contracts with Reference Copy



```
class BANK
  deposit_on_v3 (n: STRING; a: INTEGER)
  require across accounts is acc some acc.owner ~ n end
  local i: INTEGER
  do ...
    -- imp. of version 1, followed by a deposit into 1st account
    accounts[accounts.lower].deposit(a)
  ensure
    num_of_accounts_unchanged: accounts.count = old accounts.count
    balance_of_n_increased:
      Current.account_of(n).balance =
        old Current.account_of(n).balance + a
    others_unchanged:
      across old accounts is acc
      all
        acc.owner /~ n implies acc ~ Current.account_of(acc.owner)
      end
    end
  end
end
```

25/07/20

Test of Version 3: Result



APPLICATION

Note: * indicates a violation test case

| FAILED (2 failed & 1 passed out of 3) | | |
|---------------------------------------|---------------------------|---|
| Case Type | Passed | Total |
| Violation | 0 | 0 |
| Boolean | 1 | 3 |
| All Cases | 1 | 3 |
| State | Contract Violation | Test Name |
| Test1 | TEST_BANK | |
| PASSED | NONE | t1: test deposit_on with correct imp and incomplete contract |
| FAILED | Check assertion violated. | t2: test deposit_on with wrong imp but incomplete contract |
| FAILED | Check assertion violated. | t3: test deposit_on with wrong imp, complete contract with reference copy |

25/07/20



Version 4: Complete Contracts with Shallow Object Copy

```
class BANK
  deposit_on_v4 (n: STRING; a: INTEGER)
    require across accounts is acc some acc.owner ~ n end
    local i: INTEGER
    do ...
      -- imp. of version 1, followed by a deposit into 1st account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged: accounts.count = old accounts.count
      balance_of_n_increased:
        Current.account_of(n).balance =
          old Current.account_of(n).balance + a
      others_unchanged:
        across old accounts.twin is acc
        all
          acc.owner /~ n implies acc ~ Current.account_of(acc.owner)
        end
      end
    end
end
```

2/17/2018



Test of Version 4: Result

APPLICATION

Note: * indicates a violation test case

| FAILED (3 failed & 1 passed out of 4) | | |
|---------------------------------------|---------------------------|--|
| Case Type | Passed | Total |
| Violation | 0 | 0 |
| Boolean | 1 | 4 |
| All Cases | 1 | 4 |
| State | Contract Violation | Test Name |
| Test1 | | TEST_BANK |
| PASSED | NONE | t1: test deposit_on with correct imp and incomplete contract |
| FAILED | Check assertion violated. | t2: test deposit_on with wrong imp but incomplete contract |
| FAILED | Check assertion violated. | t3: test deposit_on with wrong imp, complete contract with reference copy |
| FAILED | Check assertion violated. | t4: test deposit_on with wrong imp, complete contract with shallow object copy |

2/17/2018



Test of Version 4

```
class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_shallow_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t4: wrong imp and complete contract with shallow copy")
    create b.make
    b.add("Bill")
    b.add("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v4("Steve", 100)
  Result :=
    b.account_of("Bill").balance = 0
    and b.account_of("Steve").balance = 100
  check Result end
end
```

2/17/2018



Version 5: Complete Contracts with Deep Object Copy

```
class BANK
  deposit_on_v5 (n: STRING; a: INTEGER)
    require across accounts is acc some acc.owner ~ n end
    local i: INTEGER
    do ...
      -- imp. of version 1, followed by a deposit into 1st account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged: accounts.count = old accounts.count
      balance_of_n_increased:
        Current.account_of(n).balance =
          old Current.account_of(n).balance + a
      others_unchanged:
        across old accounts.deep.twin is acc
        all
          acc.owner /~ n implies acc ~ Current.account_of(acc.owner)
        end
      end
    end
end
```

2/17/2018

Test of Version 5



```

class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_deep_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t5: wrong imp and complete contract with deep copy")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v5 ("Steve", 100)
  Result :=
    b.account_of("Bill").balance = 0
    and b.account_of("Steve").balance = 100
  check Result end
end
end
  
```

START

Test of Version 5: Result



APPLICATION

Note: * indicates a violation test case

| FAILED (4 failed & 1 passed out of 5) | | |
|---------------------------------------|---------------------------|--|
| Case Type | Passed | Total |
| Violation | 0 | 0 |
| Boolean | 1 | 5 |
| All Cases | 1 | 5 |
| State | Contract Violation | Test Name |
| Test1 | | TEST_BANK |
| PASSED | NONE | t1: test deposit_on with correct imp and incomplete contract |
| FAILED | Check assertion violated. | t2: test deposit_on with wrong imp but incomplete contract |
| FAILED | Check assertion violated. | t3: test deposit_on with wrong imp, complete contract with reference copy |
| FAILED | Check assertion violated. | t4: test deposit_on with wrong imp, complete contract with shallow object copy |
| FAILED | Postcondition violated. | t5: test deposit_on with wrong imp, complete contract with deep object copy |

START

Exercise



- Consider the query *account_of* (*n*: *STRING*) of *BANK*.
- How do we specify (part of) its postcondition to assert that the state of the bank remains unchanged:
 - `accounts = old accounts` [×]
 - `accounts = old accounts.twin` [×]
 - `accounts = old accounts.deep_twin` [×]
 - `accounts ~ old accounts` [×]
 - `accounts ~ old accounts.twin` [×]
 - `accounts ~ old accounts.deep_twin` [✓]
- Which equality of the above is appropriate for the postcondition?
- Why is each one of the other equalities not appropriate?

START

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Abstractions via Mathematical Models

Motivating Problem: Complete Contracts



- Recall what we learned in the *Complete Contracts* lecture:
 - In **post-condition**, for **each attribute**, specify the relationship between its **pre-state** value and its **post-state** value.
 - Use the **old** keyword to refer to **post-state** values of expressions.
 - For a **composite**-structured attribute (e.g., arrays, linked-lists, hash-tables, etc.), we should specify that after the update:
 - The intended change is present; **and**
 - The rest of the structure is unchanged**.
- Let's now revisit this technique by specifying a *LIFO stack*.

not 41

Motivating Problem: LIFO Stack (2.1)



```
class LIFO_STACK[G] create make
feature {NONE} -- Strategy 1: array
  imp: ARRAY[G]
feature -- Initialization
  make do create imp.make_empty ensure imp.count = 0 end
feature -- Commands
  push(g: G)
  do imp.force(g, imp.count + 1)
  ensure
    changed: imp[count] ~ g
    unchanged: across 1 |..| count - 1 as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item] end
  end
  pop
  do imp.remove_tail(1)
  ensure
    changed: count = old count - 1
    unchanged: across 1 |..| count as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item] end
  end
end
```

not 41

Motivating Problem: LIFO Stack (1)



- Let's consider three different implementation strategies:

| Stack Feature | Array | Linked List | |
|----------------|-----------------------------|---------------------------|--------------------------|
| | Strategy 1 | Strategy 2 | Strategy 3 |
| <i>count</i> | imp.count | | |
| <i>top</i> | imp[imp.count] | imp.first | imp.last |
| <i>push(g)</i> | imp.force(g, imp.count + 1) | imp.put_front(g) | imp.extend(g) |
| <i>pop</i> | imp.list.remove_tail(1) | list.start list.remove | imp.finish imp.remove |

- Given that all strategies are meant for implementing the **same ADT**, will they have **identical** contracts?

not 41

Motivating Problem: LIFO Stack (2.2)



```
class LIFO_STACK[G] create make
feature {NONE} -- Strategy 2: linked-list first item as top
  imp: LINKED_LIST[G]
feature -- Initialization
  make do create imp.make ensure imp.count = 0 end
feature -- Commands
  push(g: G)
  do imp.put_front(g)
  ensure
    changed: imp.first ~ g
    unchanged: across 2 |..| count as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item - 1] end
  end
  pop
  do imp.start ; imp.remove
  ensure
    changed: count = old count - 1
    unchanged: across 1 |..| count as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item + 1] end
  end
end
```

not 41

Motivating Problem: LIFO Stack (2.3)



```
class LIFO_STACK[G] create make
feature {NONE} -- Strategy 3: linked-list last item as top
  imp: LINKED_LIST[G]
feature -- Initialization
  make do create imp.make ensure imp.count = 0 end
feature -- Commands
  push(g: G)
  do imp.extend(g)
  ensure
    changed: imp.last ~ g
    unchanged: across 1 |..| count - 1 as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item] end
  end
  pop
  do imp.finish ; imp.remove
  ensure
    changed: count = old count - 1
    unchanged: across 1 |..| count as i all
      imp[i.item] ~ (old imp.deep_twin)[i.item] end
  end
end
```

not 41

Motivating Problem: LIFO Stack (3)



- **Postconditions** of all 3 versions of stack are **complete**. i.e., Not only the new item is **pushed/popped**, but also the remaining part of the stack is **unchanged**.
- But they violate the principle of **information hiding**: Changing the **secret**, internal workings of data structures should not affect any existing clients.
- How so?
 - The private attribute `imp` is referenced in the **postconditions**, exposing the implementation strategy not relevant to clients:
 - Top of stack may be `imp[count]`, `imp.first`, or `imp.last`.
 - Remaining part of stack may be `across 1 |..| count - 1` or `across 2 |..| count`.
 - ⇒ **Changing the implementation strategy** from one to another will also **change the contracts for all features**.
 - ⇒ This also violates the **Single Choice Principle**.

not 41

Design Principles: Information Hiding & Single Choice



- **Information Hiding** (IH):
 - Hide supplier's **design decisions** that are **likely to change**.
 - Violation of IH means that your design's public API is **unstable**.
 - **Change of supplier's secrets** should not affect clients relying upon the existing API.
- **Single Choice Principle** (SCP):
 - When a **change** is needed, there should be **a single place** (or **a minimal number of places**) where you need to make that change.
 - Violation of SCP means that your design contains **redundancies**.

not 41

Math Models: Command vs Query



- Use MATHMODELS library to create math objects (SET, REL, SEQ).
- State-changing **commands**: Implement an **Abstraction Function**

```
class LIFO_STACK[G -> attached ANY] create make
feature {NONE} -- Implementation
  imp: LINKED_LIST[G]
feature -- Abstraction function of the stack ADT
  model: SEQ[G]
  do create Result.make_empty
  across imp as cursor loop Result.append(cursor.item) end
end
```

- Side-effect-free **queries**: Write Complete Contracts

```
class LIFO_STACK[G -> attached ANY] create make
feature -- Abstraction function of the stack ADT
  model: SEQ[G]
feature -- Commands
  push(g: G)
  ensure model ~ (old model.deep_twin).appended(g) end
```

not 41

Implementing an Abstraction Function (1)



```

class LIFO_STACK[G -> attached ANY] create make
feature {NONE} -- Implementation Strategy 1
  imp: ARRAY[G]
feature -- Abstraction function of the stack ADT
  model: SEQ[G]
  do create Result.make_from_array (imp)
  ensure
    counts: imp.count = Result.count
    contents: across 1 |..| Result.count as i all
      Result[i.item] ~ imp[i.item]
  end
feature -- Commands
  make do create imp.make_empty ensure model.count = 0 end
  push (g: G) do imp.force(g, imp.count + 1)
    ensure pushed: model ~ (old model.deep_twin).appended(g) end
  pop do imp.remove_tail(1)
    ensure popped: model ~ (old model.deep_twin).front end
end
end
    
```

13/07/11

Implementing an Abstraction Function (2)



```

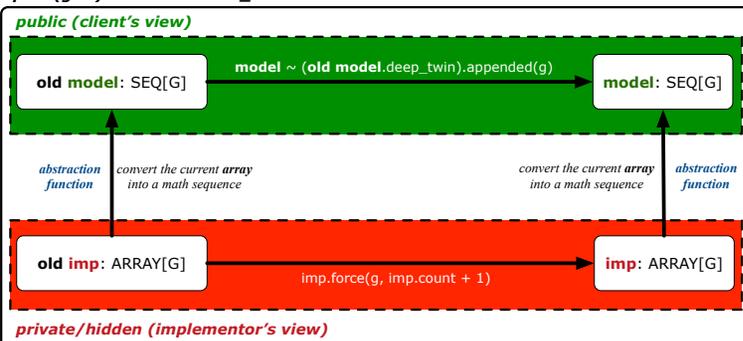
class LIFO_STACK[G -> attached ANY] create make
feature {NONE} -- Implementation Strategy 2 (first as top)
  imp: LINKED_LIST[G]
feature -- Abstraction function of the stack ADT
  model: SEQ[G]
  do create Result.make_empty
  across imp as cursor loop Result.prepend(cursor.item) end
  ensure
    counts: imp.count = Result.count
    contents: across 1 |..| Result.count as i all
      Result[i.item] ~ imp[count - i.item + 1]
  end
feature -- Commands
  make do create imp.make ensure model.count = 0 end
  push (g: G) do imp.put_front(g)
    ensure pushed: model ~ (old model.deep_twin).appended(g) end
  pop do imp.start ; imp.remove
    ensure popped: model ~ (old model.deep_twin).front end
end
end
    
```

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Abstracting ADTs as Math Models (1)



'push(g: G)' feature of LIFO_STACK ADT



- **Strategy 1** **Abstraction function**: Convert the *implementation array* to its corresponding *model sequence*.
- **Contract** for the `put (g: G)` feature remains the **same**:

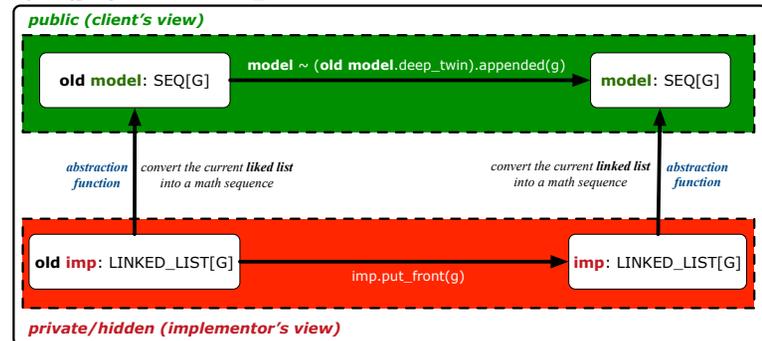
$model \sim (old\ model.deep_twin).appended(g)$

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Abstracting ADTs as Math Models (2)



'push(g: G)' feature of LIFO_STACK ADT



- **Strategy 2** **Abstraction function**: Convert the *implementation list* (first item is top) to its corresponding *model sequence*.
- **Contract** for the `put (g: G)` feature remains the **same**:

$model \sim (old\ model.deep_twin).appended(g)$

13/07/11

Implementing an Abstraction Function (3)



```

class LIFO_STACK[G -> attached ANY] create make
feature {NONE} -- Implementation Strategy 3 (last as top)
  imp: LINKED_LIST[G]
feature -- Abstraction function of the stack ADT
  model: SEQ[G]
  do create Result.make_empty
    across imp as cursor loop Result.append(cursor.item) end
  ensure
    counts: imp.count = Result.count
    contents: across 1 |..| Result.count as i all
      Result[i.item] ~ imp[i.item]
  end
feature -- Commands
  make do create imp.make ensure model.count = 0 end
  push (g: G) do imp.extend(g)
    ensure pushed: model ~ (old model.deep_twin).appended(g) end
  pop do imp.finish ; imp.remove
    ensure popped: model ~ (old model.deep_twin).front end
end
  
```

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Solution: Abstracting ADTs as Math Models



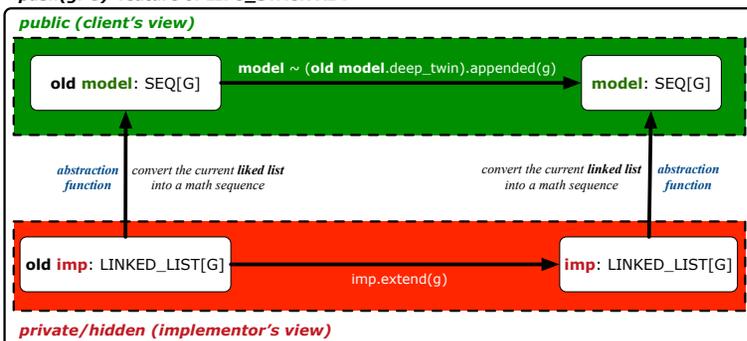
- Writing contracts in terms of *implementation attributes* (arrays, LL's, hash tables, etc.) violates **information hiding** principle.
- Instead:
 - For each ADT, create an **abstraction** via a **mathematical model**.
e.g., Abstract a LIFO_STACK as a mathematical **sequence**.
 - For each ADT, define an **abstraction function** (i.e., a query) whose return type is a kind of **mathematical model**.
e.g., Convert *implementation array* to *mathematical sequence*
 - Write contracts in terms of the **abstract math model**.
e.g., When pushing an item g onto the stack, specify it as appending g into its model sequence.
 - Upon **changing the implementation**:
 - **No change on what** the abstraction is, hence *no change on contracts*.
 - **Only change how** the abstraction is constructed, hence *changes on the body of the abstraction function*.
e.g., Convert *implementation linked-list* to *mathematical sequence*
⇒ The **Single Choice Principle** is obeyed.

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Abstracting ADTs as Math Models (3)



'push(g: G)' feature of LIFO_STACK ADT



- **Strategy 3** **Abstraction function**: Convert the *implementation list* (last item is top) to its corresponding *model sequence*.
- **Contract** for the `put (g: G)` feature remains the **same**:

```
model ~ (old model.deep_twin).appended(g)
```

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Math Review: Set Definitions and Membership



- A **set** is a collection of objects.
 - Objects in a set are called its *elements* or *members*.
 - **Order** in which elements are arranged does not matter.
 - An element can appear **at most once** in the set.
- We may define a set using:
 - **Set Enumeration**: Explicitly list all members in a set.
e.g., $\{1, 3, 5, 7, 9\}$
 - **Set Comprehension**: Implicitly specify the condition that all members satisfy.
e.g., $\{x \mid 1 \leq x \leq 10 \wedge x \text{ is an odd number}\}$
- An empty set (denoted as $\{\}$ or \emptyset) has no members.
- We may check if an element is a *member* of a set:
 - e.g., $5 \in \{1, 3, 5, 7, 9\}$ [true]
 - e.g., $4 \notin \{x \mid x \leq 1 \leq 10, x \text{ is an odd number}\}$ [true]
- The number of elements in a set is called its *cardinality*.
e.g., $|\emptyset| = 0$, $|\{x \mid x \leq 1 \leq 10, x \text{ is an odd number}\}| = 5$

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Math Review: Set Relations



Given two sets S_1 and S_2 :

- S_1 is a *subset* of S_2 if every member of S_1 is a member of S_2 .

$$S_1 \subseteq S_2 \iff (\forall x \bullet x \in S_1 \Rightarrow x \in S_2)$$

- S_1 and S_2 are *equal* iff they are the subset of each other.

$$S_1 = S_2 \iff S_1 \subseteq S_2 \wedge S_2 \subseteq S_1$$

- S_1 is a *proper subset* of S_2 if it is a strictly smaller subset.

$$S_1 \subset S_2 \iff S_1 \subseteq S_2 \wedge |S_1| < |S_2|$$

210741

Math Review: Set Operations



Given two sets S_1 and S_2 :

- *Union* of S_1 and S_2 is a set whose members are in either.

$$S_1 \cup S_2 = \{x \mid x \in S_1 \vee x \in S_2\}$$

- *Intersection* of S_1 and S_2 is a set whose members are in both.

$$S_1 \cap S_2 = \{x \mid x \in S_1 \wedge x \in S_2\}$$

- *Difference* of S_1 and S_2 is a set whose members are in S_1 but not S_2 .

$$S_1 \setminus S_2 = \{x \mid x \in S_1 \wedge x \notin S_2\}$$

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Math Review: Power Sets



The *power set* of a set S is a *set* of all S ' *subsets*.

$$\mathbb{P}(S) = \{s \mid s \subseteq S\}$$

The power set contains subsets of *cardinalities* $0, 1, 2, \dots, |S|$.
e.g., $\mathbb{P}(\{1, 2, 3\})$ is a set of sets, where each member set s has cardinality $0, 1, 2$, or 3 :

$$\left\{ \begin{array}{l} \emptyset, \\ \{1\}, \{2\}, \{3\}, \\ \{1, 2\}, \{2, 3\}, \{3, 1\}, \\ \{1, 2, 3\} \end{array} \right\}$$

210741

Math Review: Set of Tuples



Given n sets S_1, S_2, \dots, S_n , a *cross product* of these sets is a set of n -tuples.

Each *n -tuple* (e_1, e_2, \dots, e_n) contains n elements, each of which a member of the corresponding set.

$$S_1 \times S_2 \times \dots \times S_n = \{(e_1, e_2, \dots, e_n) \mid e_i \in S_i \wedge 1 \leq i \leq n\}$$

e.g., $\{a, b\} \times \{2, 4\} \times \{\$, \&\}$ is a set of triples:

$$\begin{aligned} & \{a, b\} \times \{2, 4\} \times \{\$, \&\} \\ = & \{(e_1, e_2, e_3) \mid e_1 \in \{a, b\} \wedge e_2 \in \{2, 4\} \wedge e_3 \in \{\$, \&\}\} \\ = & \{(a, 2, \$), (a, 2, \&), (a, 4, \$), (a, 4, \&), \\ & (b, 2, \$), (b, 2, \&), (b, 4, \$), (b, 4, \&)\} \end{aligned}$$

210741

Math Models: Relations (1)



- A **relation** is a collection of mappings, each being an *ordered pair* that maps a member of set S to a member of set T .
e.g., Say $S = \{1, 2, 3\}$ and $T = \{a, b\}$
 - \emptyset is an empty relation.
 - $S \times T$ is a relation (say r_1) that maps from each member of S to each member in T : $\{(1, a), (1, b), (2, a), (2, b), (3, a), (3, b)\}$
 - $\{(x, y) : S \times T \mid x \neq 1\}$ is a relation (say r_2) that maps only some members in S to every member in T : $\{(2, a), (2, b), (3, a), (3, b)\}$.
- Given a relation r :
 - Domain** of r is the set of S members that r maps from.

$$\text{dom}(r) = \{s : S \mid (\exists t \bullet (s, t) \in r)\}$$

e.g., $\text{dom}(r_1) = \{1, 2, 3\}$, $\text{dom}(r_2) = \{2, 3\}$

- Range** of r is the set of T members that r maps to.

$$\text{ran}(r) = \{t : T \mid (\exists s \bullet (s, t) \in r)\}$$

e.g., $\text{ran}(r_1) = \{a, b\} = \text{ran}(r_2)$

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Math Models: Relations (3.1)



Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- r.domain**: set of first-elements from r
 - $\text{r.domain} = \{d \mid (d, r) \in r\}$
 - e.g., $\text{r.domain} = \{a, b, c, d, e, f\}$
- r.range**: set of second-elements from r
 - $\text{r.range} = \{r \mid (d, r) \in r\}$
 - e.g., $\text{r.range} = \{1, 2, 3, 4, 5, 6\}$
- r.inverse**: a relation like r except elements are in reverse order
 - $\text{r.inverse} = \{(r, d) \mid (d, r) \in r\}$
 - e.g., $\text{r.inverse} = \{(1, a), (2, b), (3, c), (4, a), (5, b), (6, c), (1, d), (2, e), (3, f)\}$

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Math Models: Relations (2)



- We use the power set operator to express the set of *all possible relations* on S and T :

$$\mathbb{P}(S \times T)$$

- To declare a relation variable r , we use the colon ($:$) symbol to mean *set membership*:

$$r : \mathbb{P}(S \times T)$$

- Or alternatively, we write:

$$r : S \leftrightarrow T$$

where the set $S \leftrightarrow T$ is synonymous to the set $\mathbb{P}(S \times T)$

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Math Models: Relations (3.2)



Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- r.domain_restricted(ds)**: sub-relation of r with domain ds .
 - $\text{r.domain_restricted}(ds) = \{(d, r) \mid (d, r) \in r \wedge d \in ds\}$
 - e.g., $\text{r.domain_restricted}(\{a, b\}) = \{(a, 1), (b, 2), (a, 4), (b, 5)\}$
- r.domain_subtracted(ds)**: sub-relation of r with domain not ds .
 - $\text{r.domain_subtracted}(ds) = \{(d, r) \mid (d, r) \in r \wedge d \notin ds\}$
 - e.g., $\text{r.domain_subtracted}(\{a, b\}) = \{(c, 6), (d, 1), (e, 2), (f, 3)\}$
- r.range_restricted(rs)**: sub-relation of r with range rs .
 - $\text{r.range_restricted}(rs) = \{(d, r) \mid (d, r) \in r \wedge r \in rs\}$
 - e.g., $\text{r.range_restricted}(\{1, 2\}) = \{(a, 1), (b, 2), (d, 1), (e, 2)\}$
- r.range_subtracted(ds)**: sub-relation of r with range not ds .
 - $\text{r.range_subtracted}(rs) = \{(d, r) \mid (d, r) \in r \wedge r \notin rs\}$
 - e.g., $\text{r.range_subtracted}(\{1, 2\}) = \{(c, 3), (a, 4), (b, 5), (c, 6)\}$

25/07/21

Math Models: Relations (3.3)



Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- $r.\text{overridden}(t)$: a relation which agrees on r outside domain of $t.\text{domain}$, and agrees on t within domain of $t.\text{domain}$
 - $r.\text{overridden}(t) = t \cup r.\text{domain_subtracted}(t.\text{domain})$
 -

$$\begin{aligned} & r.\text{overridden}(\{(a, 3), (c, 4)\}) \\ &= \underbrace{\{(a, 3), (c, 4)\}}_t \cup \underbrace{\{(b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\}}_{r.\text{domain_subtracted}(t.\text{domain})} \\ &= \{(a, 3), (c, 4), (b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\} \end{aligned}$$

25/07/11

Math Review: Functions (1)



A **function** f on sets S and T is a *specialized form* of relation: it is forbidden for a member of S to map to more than one members of T .

$$\forall s : S; t_1 : T; t_2 : T \bullet (s, t_1) \in f \wedge (s, t_2) \in f \Rightarrow t_1 = t_2$$

e.g., Say $S = \{1, 2, 3\}$ and $T = \{a, b\}$, which of the following relations are also functions?

- $S \times T$ [No]
- $(S \times T) - \{(x, y) \mid (x, y) \in S \times T \wedge x = 1\}$ [No]
- $\{(1, a), (2, b), (3, a)\}$ [Yes]
- $\{(1, a), (2, b)\}$ [Yes]

25/07/11

Math Review: Functions (2)



- We use *set comprehension* to express the set of all possible functions on S and T as those relations that satisfy the **functional property**:

$$\{r : S \leftrightarrow T \mid (\forall s : S; t_1 : T; t_2 : T \bullet (s, t_1) \in r \wedge (s, t_2) \in r \Rightarrow t_1 = t_2)\}$$

- This set (of possible functions) is a subset of the set (of possible relations): $\mathbb{P}(S \times T)$ and $S \leftrightarrow T$.
- We abbreviate this set of possible functions as $S \rightarrow T$ and use it to declare a function variable f :

$$f : S \rightarrow T$$

25/07/11

Math Review: Functions (3.1)



Given a function $f : S \rightarrow T$:

- f is *injective* (or an injection) if f does not map a member of S to more than one members of T .

$$f \text{ is injective} \iff (\forall s_1 : S; s_2 : S; t : T \bullet (s_1, t) \in f \wedge (s_2, t) \in f \Rightarrow s_1 = s_2)$$

e.g., Considering an array as a function from integers to objects, being injective means that the array does not contain any duplicates.

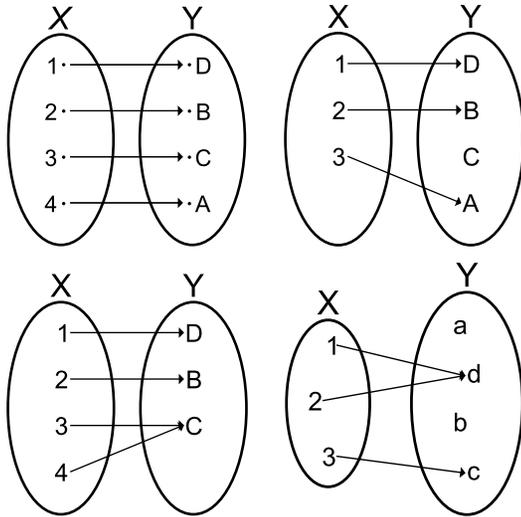
- f is *surjective* (or a surjection) if f maps to all members of T .

$$f \text{ is surjective} \iff \text{ran}(f) = T$$

- f is *bijective* (or a bijection) if f is both injective and surjective.

25/07/11

Math Review: Functions (3.2)



3/20/21

Math Models: Example Test



```
test_rel: BOOLEAN
local
  r, t: REL[STRING, INTEGER]
  ds: SET[STRING]
do
  create r.make_from_tuple_array (
    <<["a", 1], ["b", 2], ["c", 3],
      ["a", 4], ["b", 5], ["c", 6],
      ["d", 1], ["e", 2], ["f", 3]>>)
  create ds.make_from_array (<<"a">>)
  -- r is not changed by the query 'domain_subtracted'
  t := r.domain_subtracted (ds)
  Result :=
    t /~ r and not t.domain.has ("a") and r.domain.has ("a")
  check Result end
  -- r is changed by the command 'domain_subtract'
  r.domain_subtract (ds)
  Result :=
    t ~ r and not t.domain.has ("a") and not r.domain.has ("a")
end
```

3/20/21

Math Models: Command-Query Separation



| Command | Query |
|--------------------|----------------------|
| domain_restrict | domain_restricted |
| domain_restrict_by | domain_restricted_by |
| domain_subtract | domain_subtracted |
| domain_subtract_by | domain_subtracted_by |
| range_restrict | range_restricted |
| range_restrict_by | range_restricted_by |
| range_subtract | range_subtracted |
| range_subtract_by | range_subtracted_by |
| override | overridden |
| override_by | overridden_by |

Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- Commands** modify the context relation objects.

`r.domain_restrict({a})` changes r to $\{(a, 1), (a, 4)\}$

- Queries** return new relations without modifying context objects.

`r.domain_restricted({a})` returns $\{(a, 1), (a, 4)\}$ with r untouched

3/20/21

Case Study: A Birthday Book



- A birthday book stores a collection of entries, where each entry is a pair of a person's name and their birthday.
- No two entries stored in the book are allowed to have the same name.
- Each birthday is characterized by a month and a day.
- A birthday book is first created to contain an empty collection of entries.
- Given a birthday book, we may:
 - Inquire about the number of entries currently stored in the book
 - Add a new entry by supplying its name and the associated birthday
 - Remove the entry associated with a particular person
 - Find the birthday of a particular person
 - Get a reminder list of names of people who share a given birthday

3/20/21

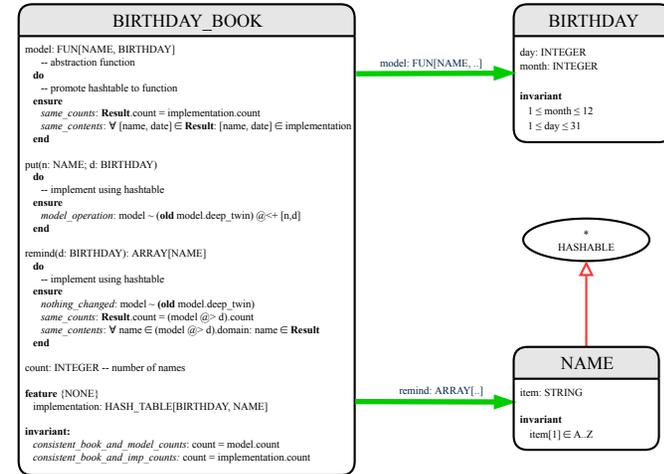
Birthday Book: Decisions



- **Design** Decision
 - Classes
 - Client Supplier vs. Inheritance
 - Mathematical Model? [e.g., REL or FUN]
 - Contracts
- **Implementation** Decision
 - Two linear structures (e.g., arrays, lists) [$O(n)$]
 - A balanced search tree (e.g., AVL tree) [$O(\log \cdot n)$]
 - A hash table [$O(1)$]
- Implement an **abstract function** that maps implementation to the math model.

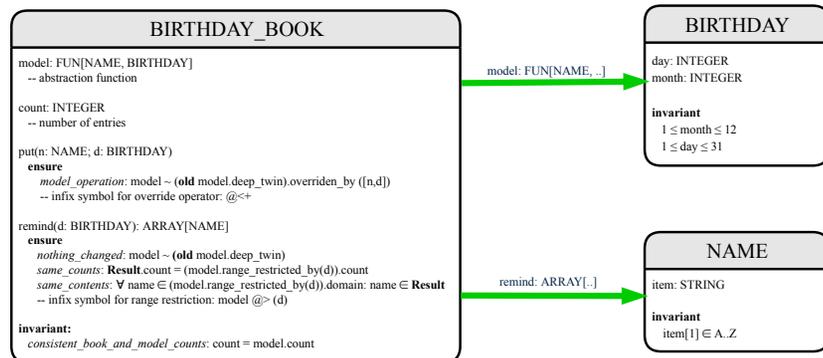
3.5.07.41

Birthday Book: Implementation



3.5.07.41

Birthday Book: Design



3.5.07.41

Beyond this lecture ...



- Familiarize yourself with the features of classes SEQ, REL, FUN, and SET for the lab test.
- Play with the source code of the Birthday Book example: <https://www.eecs.vorku.ca/~jackie/teaching/lectures/2020/W/EECS3311/codes/birthday-book.zip>.
- **Exercise:**
 - Consider an alternative implementation using two linear structures (e.g., [here in Java](#)).
 - Implement the design of birthday book covered in lectures.
 - Create another LINEAR_BIRTHDAY_BOOK class and modify the implementation of abstraction function accordingly. Do all contracts still pass? What should change? What remain unchanged?

3.5.07.41

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Use of Generic Parameters Iterator and Singleton Patterns



EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG

Generic Collection Class: Motivation (2)



```
class ACCOUNT_STACK
feature {NONE} -- Implementation
  imp: ARRAY[ACCOUNT] ; i: INTEGER
feature -- Queries
  count: INTEGER do Result := i end
  -- Number of items on stack.
  top: ACCOUNT do Result := imp [i] end
  -- Return top of stack.
feature -- Commands
  push (v: ACCOUNT) do imp[i] := v; i := i + 1 end
  -- Add 'v' to top of stack.
  pop do i := i - 1 end
  -- Remove top of stack.
end
```

- Does how we implement integer stack operations (e.g., top, push, pop) depends on features specific to element type ACCOUNT (e.g., deposit, withdraw)? [NO!]
- A **collection** (e.g., table, tree, graph) is meant for the **storage** and **retrieval** of elements, not how those elements are manipulated.

not 4c

Generic Collection Class: Motivation (1)



```
class STRING_STACK
feature {NONE} -- Implementation
  imp: ARRAY[STRING] ; i: INTEGER
feature -- Queries
  count: INTEGER do Result := i end
  -- Number of items on stack.
  top: STRING do Result := imp [i] end
  -- Return top of stack.
feature -- Commands
  push (v: STRING) do imp[i] := v; i := i + 1 end
  -- Add 'v' to top of stack.
  pop do i := i - 1 end
  -- Remove top of stack.
end
```

- Does how we implement integer stack operations (e.g., top, push, pop) depends on features specific to element type STRING (e.g., at, append)? [NO!]
- How would you implement another class ACCOUNT_STACK?

not 4c

Generic Collection Class: Supplier



- Your design **"smells"** if you have to create an **almost identical** new class (hence **code duplicates**) for every stack element type you need (e.g., INTEGER, CHARACTER, PERSON, etc.).
- Instead, as **supplier**, use **G** to **parameterize** element type:

```
class STACK[G]
feature {NONE} -- Implementation
  imp: ARRAY[G] ; i: INTEGER
feature -- Queries
  count: INTEGER do Result := i end
  -- Number of items on stack.
  top: G do Result := imp [i] end
  -- Return top of stack.
feature -- Commands
  push (v: G) do imp[i] := v; i := i + 1 end
  -- Add 'v' to top of stack.
  pop do i := i - 1 end
  -- Remove top of stack.
end
```

not 4c

Generic Collection Class: Client (1.1)



As **client**, declaring `ss: STACK[STRING]` instantiates every occurrence of `G` as `STRING`.

```
class STACK [G STRING]
feature {NONE} -- Implementation
  imp: ARRAY[G STRING] ; i: INTEGER
feature -- Queries
  count: INTEGER do Result := i end
  -- Number of items on stack.
  top: G STRING do Result := imp [i] end
  -- Return top of stack.
feature -- Commands
  push (v: G STRING) do imp[i] := v; i := i + 1 end
  -- Add 'v' to top of stack.
  pop do i := i - 1 end
  -- Remove top of stack.
end
```

not 4c

Generic Collection Class: Client (2)



As **client**, instantiate the type of `G` to be the one needed.

```
1 test_stacks: BOOLEAN
2   local
3     ss: STACK[STRING] ; sa: STACK[ACCOUNT]
4     s: STRING ; a: ACCOUNT
5   do
6     ss.push("A")
7     ss.push(create {ACCOUNT}.make ("Mark", 200))
8     s := ss.top
9     a := ss.top
10    sa.push(create {ACCOUNT}.make ("Alan", 100))
11    sa.push("B")
12    a := sa.top
13    s := sa.top
14  end
```

- **L3** commits that `ss` stores `STRING` objects only.
 - **L8** and **L10** *valid*; **L9** and **L11** *invalid*.
- **L4** commits that `sa` stores `ACCOUNT` objects only.
 - **L12** and **L14** *valid*; **L13** and **L15** *invalid*.

not 4c

Generic Collection Class: Client (1.2)



As **client**, declaring `ss: STACK[ACCOUNT]` instantiates every occurrence of `G` as `ACCOUNT`.

```
class STACK [G ACCOUNT]
feature {NONE} -- Implementation
  imp: ARRAY[G ACCOUNT] ; i: INTEGER
feature -- Queries
  count: INTEGER do Result := i end
  -- Number of items on stack.
  top: G ACCOUNT do Result := imp [i] end
  -- Return top of stack.
feature -- Commands
  push (v: G ACCOUNT) do imp[i] := v; i := i + 1 end
  -- Add 'v' to top of stack.
  pop do i := i - 1 end
  -- Remove top of stack.
end
```

not 4c

What are design patterns?



- Solutions to *recurring problems* that arise when software is being developed within a particular *context*.
 - Heuristics for structuring your code so that it can be systematically maintained and extended.
 - **Caveat**: A pattern is only suitable for a particular problem.
 - Therefore, always understand *problems* before *solutions*!

not 4c

Iterator Pattern: Motivation (1)



Supplier:

```
class
  CART
  feature
    orders: ARRAY[ORDER]
  end

class
  ORDER
  feature
    price: INTEGER
    quantity: INTEGER
  end
```

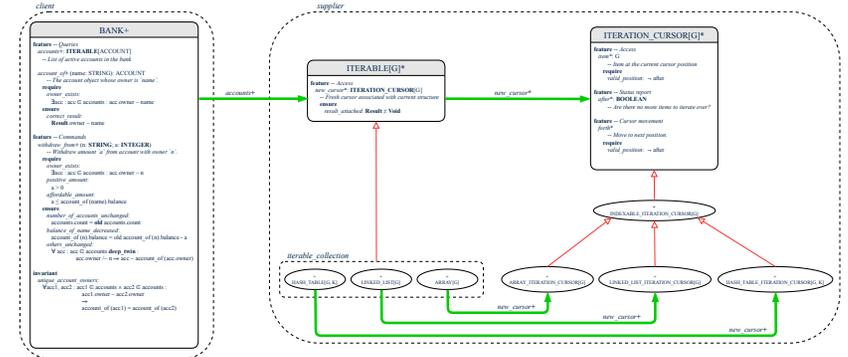
Client:

```
class
  SHOP
  feature
    cart: CART
    checkout: INTEGER
  do
    from
      i := cart.orders.lower
    until
      i > cart.orders.upper
    do
      Result := Result +
        cart.orders[i].price
        *
        cart.orders[i].quantity
      i := i + 1
    end
  end
end
```

Problems?

10/14

Iterator Pattern: Architecture



10/14

Iterator Pattern: Motivation (2)



Supplier:

```
class
  CART
  feature
    orders: LINKED_LIST[ORDER]
  end

class
  ORDER
  feature
    price: INTEGER
    quantity: INTEGER
  end
```

Client:

```
class
  SHOP
  feature
    cart: CART
    checkout: INTEGER
  do
    from
      cart.orders.start
    until
      cart.orders.after
    do
      Result := Result +
        cart.orders.item.price
        *
        cart.orders.item.quantity
    end
  end
end
```

Client's code must be modified to adapt to the supplier's change on implementation.

10/14

Iterator Pattern: Supplier's Side



- **Information Hiding Principle**:
 - Hide design decisions that are *likely to change* (i.e., *stable API*).
 - *Change of secrets* does not affect clients using the existing API. e.g., changing from *ARRAY* to *LINKED_LIST* in the *CART* class
- Steps:
 1. Let the supplier class inherit from the deferred class *ITERABLE[G]*.
 2. This forces the supplier class to implement the inherited feature: *new_cursor: ITERATION_CURSOR [G]*, where the type parameter *G* may be instantiated (e.g., *ITERATION_CURSOR[ORDER]*).
 - 2.1 If the internal, library data structure is already *iterable* e.g., *imp: ARRAY[ORDER]*, then simply return *imp.new_cursor*.
 - 2.2 Otherwise, say *imp: MY_TREE[ORDER]*, then create a new class *MY_TREE.ITERATION_CURSOR* that inherits from *ITERATION_CURSOR[ORDER]*, then implement the 3 inherited features *after*, *item*, and *forth* accordingly.

10/14

Iterator Pattern: Supplier's Implementation (1)



```
class
  CART
inherit
  ITERABLE[ORDER]
...

feature {NONE} -- Information Hiding
  orders: ARRAY[ORDER]

feature -- Iteration
  new_cursor: ITERATION_CURSOR[ORDER]
  do
    Result := orders.new_cursor
  end
```

When the secrete implementation is already *iterable*, reuse it!

13/07/09

Iterator Pattern: Supplier's Imp. (2.2)



```
class
  MY_ITERATION_CURSOR[G]
inherit
  ITERATION_CURSOR[ TUPLE[STRING, G] ]
feature -- Constructor
  make (ns: ARRAY[STRING]; rs: ARRAY[G])
  do ... end
feature {NONE} -- Information Hiding
  cursor_position: INTEGER
  names: ARRAY[STRING]
  records: ARRAY[G]
feature -- Cursor Operations
  item: TUPLE[STRING, G]
  do ... end
  after: Boolean
  do ... end
  forth
  do ... end
```

You need to implement the three inherited features:
item, *after*, and *forth*.

13/07/09

Iterator Pattern: Supplier's Imp. (2.1)



```
class
  GENERIC_BOOK[G]
inherit
  ITERABLE[ TUPLE[STRING, G] ]
...
feature {NONE} -- Information Hiding
  names: ARRAY[STRING]
  records: ARRAY[G]
feature -- Iteration
  new_cursor: ITERATION_CURSOR[ TUPLE[STRING, G] ]
  local
    cursor: MY_ITERATION_CURSOR[G]
  do
    create cursor.make (names, records)
    Result := cursor
  end
```

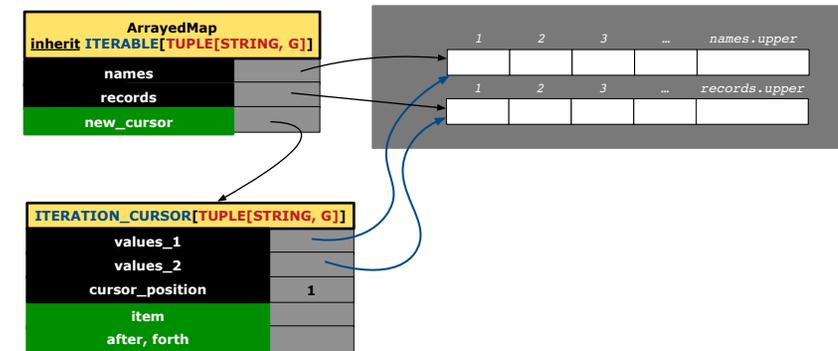
No Eiffel library support for iterable arrays ⇒ Implement it yourself!

13/07/09

Iterator Pattern: Supplier's Imp. (2.3)



Visualizing iterator pattern at runtime:



13/07/09

Exercises



1. Draw the BON diagram showing how the iterator pattern is applied to the *CART* (supplier) and *SHOP* (client) classes.
2. Draw the BON diagram showing how the iterator pattern is applied to the supplier classes:
 - *GENERIC_BOOK* (a descendant of *ITERABLE*) and
 - *MY_ITERATION_CURSOR* (a descendant of *ITERATION_CURSOR*).

13/07/19

Iterator Pattern: Client's Side



Information hiding: the clients do not at all depend on *how* the supplier implements the collection of data; they are only interested in iterating through the collection in a linear manner.

Steps:

1. Obey the **code to interface, not to implementation** principle.
2. Let the client declare an attribute of **interface** type *ITERABLE[G]* (rather than **implementation** type *ARRAY*, *LINKED_LIST*, or *MY_TREE*).
e.g., `cart: CART`, where *CART* inherits *ITERABLE[ORDER]*
3. Eiffel supports, in **both** implementation and **contracts**, the **across** syntax for iterating through anything that's *iterable*.

13/07/19

Resources



- [Tutorial Videos on Generic Parameters and the Iterator Pattern](#)
- [Tutorial Videos on Information Hiding and the Iterator Pattern](#)

13/07/19

Iterator Pattern: Clients using across for Contracts (1)



```
class
  CHECKER
  feature -- Attributes
    collection: ITERABLE [INTEGER]
  feature -- Queries
    is_all_positive: BOOLEAN
    -- Are all items in collection positive?
    do
      ...
    ensure
      across
        collection is item
      all
        item > 0
      end
    end
end
```

- Using **all** corresponds to a universal quantification (i.e., \forall).
- Using **some** corresponds to an existential quantification (i.e., \exists).

13/07/19



Iterator Pattern: Clients using across for Contracts (2)

```
class BANK
...
accounts: LIST [ACCOUNT]
binary_search (acc_id: INTEGER): ACCOUNT
  -- Search on accounts sorted in non-descending order.
  require
  across
    1 |..| (accounts.count - 1) is i
  all
    accounts [i].id <= accounts [i + 1].id
  end
do
...
ensure
  Result.id = acc_id
end
```

This precondition corresponds to:

$\forall i: \text{INTEGER} \mid 1 \leq i < \text{accounts.count} \bullet \text{accounts}[i].\text{id} \leq \text{accounts}[i+1].\text{id}$

22/07/19



Iterator Pattern: Clients using Iterable in Imp. (1)

```
class BANK
accounts: ITERABLE [ACCOUNT]
max_balance: ACCOUNT
  -- Account with the maximum balance value.
require ??
local
cursor: ITERATION_CURSOR [ACCOUNT]; max: ACCOUNT
do
  from cursor := accounts.new_cursor; max := cursor.item
  until cursor.after
  do
    if cursor.item.balance > max.balance then
      max := cursor.item
    end
    cursor.forth
  end
end
ensure ??
end
```

22/07/19



Iterator Pattern: Clients using across for Contracts (3)

```
class BANK
...
accounts: LIST [ACCOUNT]
contains_duplicate: BOOLEAN
  -- Does the account list contain duplicate?
do
...
ensure
   $\forall i, j: \text{INTEGER} \mid$ 
     $1 \leq i \leq \text{accounts.count} \wedge 1 \leq j \leq \text{accounts.count} \bullet$ 
     $\text{accounts}[i] \sim \text{accounts}[j] \Rightarrow i = j$ 
end
```

- **Exercise:** Convert this mathematical predicate for postcondition into Eiffel.
- **Hint:** Each **across** construct can only introduce one dummy variable, but you may nest as many **across** constructs as necessary.

22/07/19



Iterator Pattern: Clients using Iterable in Imp. (2)

```
1 class SHOP
2   cart: CART
3   checkout: INTEGER
4   -- Total price calculated based on orders in the cart.
5   require ??
6   do
7     across
8       cart is order
9     loop
10      Result := Result + order.price * order.quantity
11    end
12  ensure ??
13  end
```

- Class *CART* should inherit from *ITERABLE[ORDER]*.
- **L10** implicitly declares `cursor: ITERATION_CURSOR[ORDER]` and does `cursor := cart.new_cursor`

22/07/19

Iterator Pattern: Clients using Iterable in Imp. (3)

```

class BANK
  accounts: LIST[ACCOUNT] -- Q: Can ITERABLE[ACCOUNT] work?
  max_balance: ACCOUNT
  -- Account with the maximum balance value.
  require ??
  local
    max: ACCOUNT
  do
    max := accounts [1]
    across
      accounts is acc
    loop
      if acc.balance > max.balance then
        max := acc
      end
    end
  ensure ??
end
  
```

25/07/09

Expanded Class: Programming (2)

```

class KEYBOARD ... end class CPU ... end
class MONITOR ... end class NETWORK ... end
class WORKSTATION
  k: expanded KEYBOARD
  c: expanded CPU
  m: expanded MONITOR
  n: NETWORK
end
  
```

Alternatively:

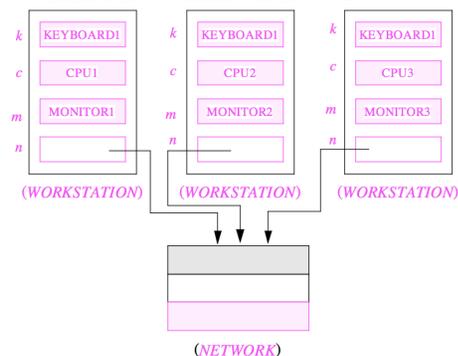
```

expanded class KEYBOARD ... end
expanded class CPU ... end
expanded class MONITOR ... end
class NETWORK ... end
class WORKSTATION
  k: KEYBOARD
  c: CPU
  m: MONITOR
  n: NETWORK
end
  
```

27/07/09

Expanded Class: Modelling

- We may want to have objects which are:
 - Integral parts of some other objects
 - Not** shared among objects
- e.g., Each workstation has its own CPU, monitor, and keyboard.
All workstations share the same network.



28/07/09

Expanded Class: Programming (3)

```

expanded class
  B
  feature
    change_i (ni: INTEGER)
    do
      i := ni
    end
  feature
    i: INTEGER
  end
end
  
```

```

1 test_expanded
2 local
3   eb1, eb2: B
4 do
5   check eb1.i = 0 and eb2.i = 0 end
6   check eb1 = eb2 end
7   eb2.change_i (15)
8   check eb1.i = 0 and eb2.i = 15 end
9   check eb1 /= eb2 end
10  eb1 := eb2
11  check eb1.i = 15 and eb2.i = 15 end
12  eb1.change_i (10)
13  check eb1.i = 10 and eb2.i = 15 end
14  check eb1 /= eb2 end
15 end
  
```

- L5**: object of expanded type is automatically initialized.
- L10,L12,L13**: no sharing among objects of expanded type.
- L6,L9,L14**: = compares contents between expanded objects.

28/07/09

Reference vs. Expanded (1)



- Every entity must be declared to be of a certain type (based on a class).
- Every type is either *referenced* or *expanded*.
- In *reference* types:
 - y denotes *a reference* to some object
 - $x := y$ attaches x to same object as does y
 - $x = y$ compares references
- In *expanded* types:
 - y denotes *some object* (of expanded type)
 - $x := y$ copies contents of y into x
 - $x = y$ compares contents

$[x \sim y]$

20/07/09

Singleton Pattern: Motivation



Consider two problems:

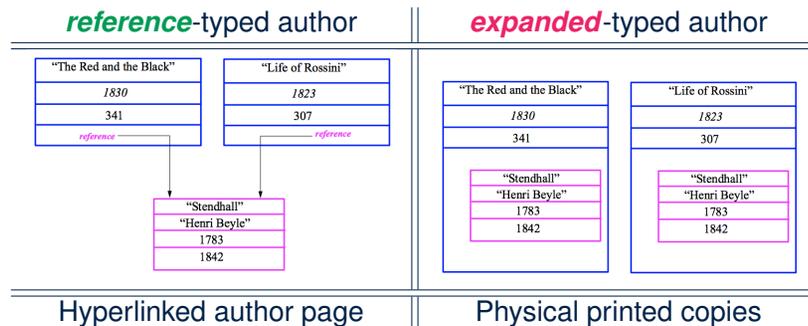
1. *Bank accounts* share a set of data.
e.g., interest and exchange rates, minimum and maximum balance, etc.
2. *Processes* are regulated to access some shared, limited resources.
e.g., printers

20/07/09

Reference vs. Expanded (2)



Problem: Every published book has an author. Every author may publish more than one books. Should the author field of a book *reference*-typed or *expanded*-typed?



20/07/09

Shared Data via Inheritance



Descendant:

```
class DEPOSIT inherit SHARED_DATA
  -- 'maximum_balance' relevant
end

class WITHDRAW inherit SHARED_DATA
  -- 'minimum_balance' relevant
end

class INT_TRANSFER inherit SHARED_DATA
  -- 'exchange_rate' relevant
end

class ACCOUNT inherit SHARED_DATA
feature
  -- 'interest_rate' relevant
  deposits: DEPOSIT_LIST
  withdraws: WITHDRAW_LIST
end
```

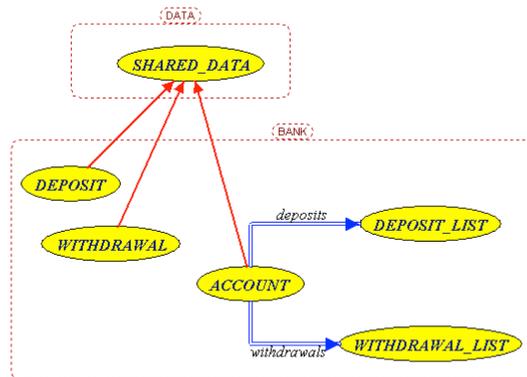
Ancestor:

```
class
  SHARED_DATA
feature
  interest_rate: REAL
  exchange_rate: REAL
  minimum_balance: INTEGER
  maximum_balance: INTEGER
  ...
end
```

Problems?

20/07/09

Sharing Data via Inheritance: Architecture



- *Irreverent* features are inherited.
⇒ Descendants' **cohesion** is broken.
- Same set of data is *duplicated* as instances are created.
⇒ Updates on these data may result in **inconsistency**.

3/5/2019

Introducing the Once Routine in Eiffel (1.1)



```
1 class A
2 create make
3 feature -- Constructor
4   make do end
5 feature -- Query
6   new_once_array (s: STRING): ARRAY[STRING]
7     -- A once query that returns an array.
8     once
9       create {ARRAY[STRING]} Result.make_empty
10      Result.force (s, Result.count + 1)
11    end
12   new_array (s: STRING): ARRAY[STRING]
13     -- An ordinary query that returns an array.
14     do
15       create {ARRAY[STRING]} Result.make_empty
16       Result.force (s, Result.count + 1)
17     end
18 end
```

L9 & L10 executed **only once** for initialization.

L15 & L16 executed **whenever** the feature is called.

3/5/2019

Sharing Data via Inheritance: Limitation



- Each descendant instance at runtime owns a separate copy of the shared data.
- This makes inheritance *not* an appropriate solution for both problems:
 - What if the interest rate changes? Apply the change to all instantiated account objects?
 - An update to the global lock must be observable by all regulated processes.

Solution:

- Separate notions of **data** and its **shared access** in two separate classes.
- **Encapsulate** the shared access itself in a separate class.

3/5/2019

Introducing the Once Routine in Eiffel (1.2)



```
1 test_query: BOOLEAN
2 local
3   a: A
4   arr1, arr2: ARRAY[STRING]
5 do
6   create a.make
7
8   arr1 := a.new_array ("Alan")
9   Result := arr1.count = 1 and arr1[1] ~ "Alan"
10  check Result end
11
12  arr2 := a.new_array ("Mark")
13  Result := arr2.count = 1 and arr2[1] ~ "Mark"
14  check Result end
15
16  Result := not (arr1 = arr2)
17  check Result end
18 end
```

3/5/2019

Introducing the Once Routine in Eiffel (1.3)



```
1 test_once_query: BOOLEAN
2   local
3     a: A
4     arr1, arr2: ARRAY[STRING]
5   do
6     create a.make
7
8     arr1 := a.new_once_array ("Alan")
9     Result := arr1.count = 1 and arr1[1] ~ "Alan"
10    check Result end
11
12    arr2 := a.new_once_array ("Mark")
13    Result := arr2.count = 1 and arr2[1] ~ "Alan"
14    check Result end
15
16    Result := arr1 = arr2
17    check Result end
18  end
```

3/3/2016

Introducing the Once Routine in Eiffel (2)



```
r (...): T
  once
    -- Some computations on Result
    ...
  end
```

- The ordinary **do ... end** is replaced by **once ... end**.
- The first time the **once** routine *r* is called by some client, it executes the body of computations and returns the computed result.
- From then on, the computed result is “*cached*”.
- In every subsequent call to *r*, possibly by different clients, the body of *r* is not executed at all; instead, it just returns the “*cached*” result, which was computed in the very first call.
- **How does this help us?**

Cache the reference to the same shared object !

3/3/2016

Approximating Once Routine in Java (1)



We may encode Eiffel once routines in Java:

```
class BankData {
  BankData() { }
  double interestRate;
  void setIR(double r);
  ...
}
```

```
class Account {
  BankData data;
  Account() {
    data = BankDataAccess.getData();
  }
}
```

```
class BankDataAccess {
  static boolean initOnce;
  static BankData data;
  static BankData getData() {
    if(!initOnce) {
      data = new BankData();
      initOnce = true;
    }
    return data;
  }
}
```

Problem?

Multiple *BankData* objects may be created in *Account*, breaking the singleton!

```
Account() {
  data = new BankData();
}
```

3/3/2016

Approximating Once Routine in Java (2)



We may encode Eiffel once routines in Java:

```
class BankData {
  private BankData() { }
  double interestRate;
  void setIR(double r);
  static boolean initOnce;
  static BankData data;
  static BankData getData() {
    if(!initOnce) {
      data = new BankData();
      initOnce = true;
    }
    return data;
  }
}
```

Problem?

Loss of Cohesion: *Data* and *Access to Data* are two separate concerns, so should be decoupled into two different classes!

3/3/2016

Singleton Pattern in Eiffel (1)

Supplier:

```
class DATA
create {DATA_ACCESS} make
feature {DATA_ACCESS}
  make do v := 10 end
feature -- Data Attributes
  v: INTEGER
  change_v (nv: INTEGER)
    do v := nv end
end
```

```
expanded class
  DATA_ACCESS
feature
  data: DATA
  -- The one and only access
  once create Result.make end
invariant data = data
```

Client:

```
test: BOOLEAN
local
  access: DATA_ACCESS
  d1, d2: DATA
do
  d1 := access.data
  d2 := access.data
  Result := d1 = d2
  and d1.v = 10 and d2.v = 10
  check Result end
  d1.change_v (15)
  Result := d1 = d2
  and d1.v = 15 and d2.v = 15
end
end
```

Writing `create d1.make` in test feature does not compile. Why?

12/14/09

Testing Singleton Pattern in Eiffel

```
test_bank_shared_data: BOOLEAN
-- Test that a single data object is manipulated
local acc1, acc2: ACCOUNT
do
  comment ("t1: test that a single data object is shared")
  create acc1.make ("Bill")
  create acc2.make ("Steve")
  Result := acc1.data = acc2.data
  check Result end
  Result := acc1.data ~ acc2.data
  check Result end
  acc1.data.set_interest_rate (3.11)
  Result :=
    acc1.data.interest_rate = acc2.data.interest_rate
  and acc1.data.interest_rate = 3.11
  check Result end
  acc2.data.set_interest_rate (2.98)
  Result :=
    acc1.data.interest_rate = acc2.data.interest_rate
  and acc1.data.interest_rate = 2.98
end
end
```

12/14/09

Singleton Pattern in Eiffel (2)

Supplier:

```
class BANK_DATA
create {BANK_DATA_ACCESS} make
feature {BANK_DATA_ACCESS}
  make do ... end
feature -- Data Attributes
  interest_rate: REAL
  set_interest_rate (r: REAL)
  ...
end
```

```
expanded class
  BANK_DATA_ACCESS
feature
  data: BANK_DATA
  -- The one and only access
  once create Result.make end
invariant data = data
```

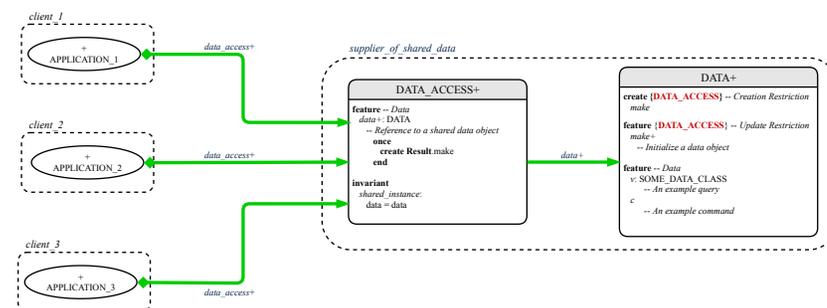
Client:

```
class
  ACCOUNT
feature
  data: BANK_DATA
  make (...)
  -- Init. access to bank data.
local
  data_access: BANK_DATA_ACCESS
do
  data := data_access.data
  ...
end
end
```

Writing `create data.make` in client's make feature does not compile. Why?

12/14/09

Singleton Pattern: Architecture



Important Exercises: Instantiate this architecture to both problems of shared bank data and shared lock. Draw them in draw.io.

12/14/09

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Bank ATM

The ATM application has a variety of *concrete* user interfaces.



Eiffel Testing Framework (ETF): Acceptance Tests via Abstract User Interface

Separation of Concerns

- The (Concrete) User Interface
 - The executable of your application *hides* the implementing classes and features.
 - Users typically interact with your application via some GUI. e.g., web app, mobile app, or desktop app
- The *Business Logic (Model)*
 - When you develop your application software, you implement classes and features. e.g., How the bank stores, processes, retrieves information about accounts and transactions

In practice:

- You need to test your software as if it were a real app *way before* dedicating to the design of an actual GUI.
- The model should be *independent* of the View, Input and Output.

Prototyping System with Abstract UI



- For you to quickly prototype a working system, you do not need to spend time on developing a fancy GUI.
- The **Eiffel Testing Framework (ETF)** allows you to:
 - Focus on developing the business model;
 - Test your business model as if it were a real app.
- In ETF, observable interactions with the application GUI (e.g., “button clicks”) are **abstracted** as monitored events.

| Events | Features |
|------------------|-------------------------|
| interactions | computations |
| external | internal |
| observable | hidden |
| acceptance tests | unit tests |
| users, customers | programmers, developers |



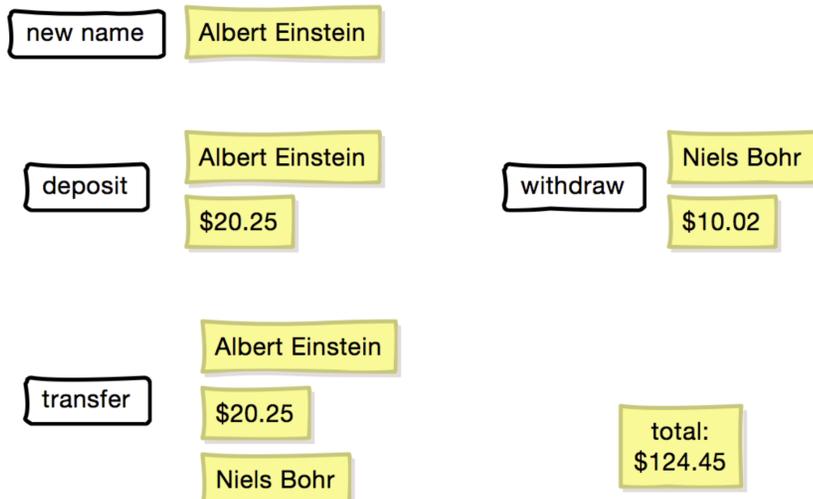
ETF in a Nutshell



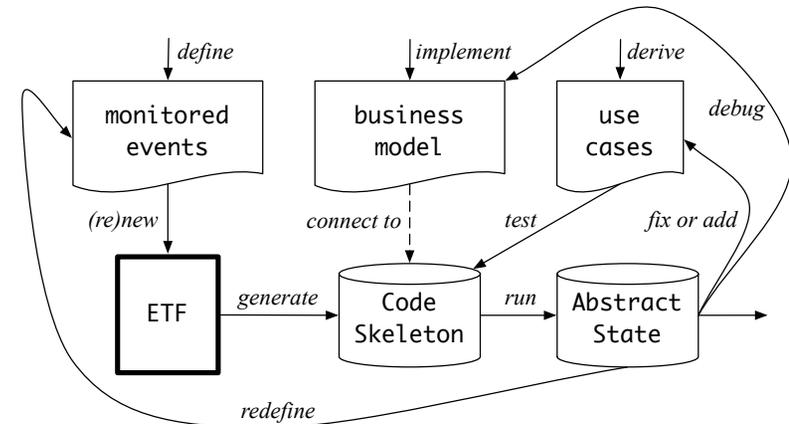
- **Eiffel Testing Framework (ETF)** facilitates engineers to write and execute **input-output-based acceptance tests**.
 - **Inputs** are specified as traces of events (or sequences).
 - The **boundary** of the system under development (SUD) is defined by declaring the list of input events that might occur.
 - **Outputs** (from executing events in the input trace) are by default logged onto the terminal, and their formats may be customized.
- An executable ETF that is tailored for the SUD can already be generated, using these event declarations (documented in a plain text file), with a default **business model**.
- Once the **business model** is implemented, there is only a small number of steps to follow for the developers to connect it to the generated ETF.
- Once connected, developers may re-run all use cases and observe if the expected state effects take place.



Abstract Events: Bank ATM



Workflow: Develop-Connect-Test



ETF: Abstract User Interface



Input Grammar

```

system bank
type NAME = STRING

new(name1: NAME)
-- create a new bank account for "id"

deposit(name1: NAME; amount: VALUE)
-- deposit "amount" into the account of "id"

withdraw(name1: NAME; amount: VALUE)
-- withdraw "amount" from the account of "id"

transfer(name1: NAME; name2: NAME; amount: VALUE)
-- transfer "amount" from "id1" to "id2"
                
```

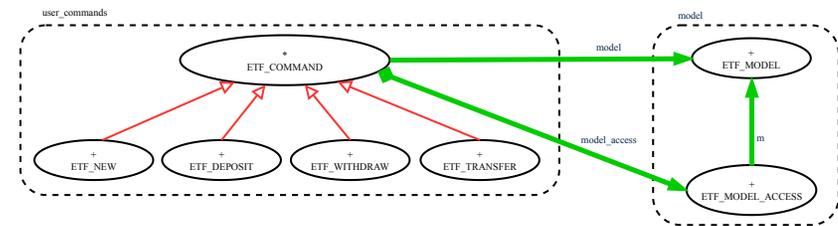
```

%bank -b at1.txt
init
->new("Steve")
name: Steve, balance: 0.00
->new("Bill")
name: Bill, balance: 0.00
name: Steve, balance: 0.00
->deposit("Steve",520)
name: Bill, balance: 0.00
name: Steve, balance: 520.00
->new("Pam")
name: Bill, balance: 0.00
name: Pam, balance: 0.00
name: Steve, balance: 520.00
->deposit("Bill",100)
name: Bill, balance: 100.00
name: Pam, balance: 0.00
name: Steve, balance: 520.00
->withdraw("Steve",20)
name: Bill, balance: 100.00
name: Pam, balance: 0.00
name: Steve, balance: 500.00
                
```

User Interface



ETF: Architecture



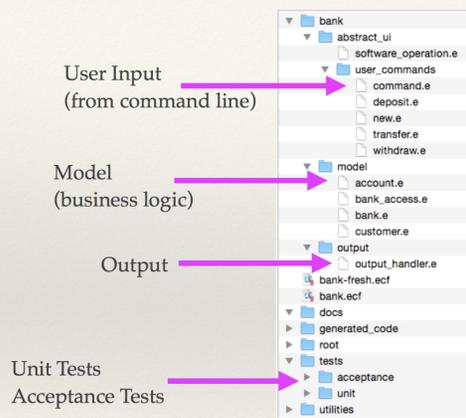
- Classes in the `model` cluster are hidden from the users.
- All commands reference to the same model (bank) instance.
- When a user's request is made:
 - A **command object** of the corresponding type is created, which invokes relevant feature(s) in the `model` cluster.
 - Updates to the model are published to the output handler.



ETF: Generating a New Project



`etf -new bank.input.txt <directory>`



ETF: Input Errors



```

class
  ETF_DEPOSIT
inherit
  ETF_DEPOSIT_INTERFACE
  redefine deposit end
create
  make
feature -- command
  deposit(id: STRING ; amount: REAL_64)
  do
    if not model.has_user (id) then
      -- Set some error message
    elseif not amount <= model.get_balance (id) then
      -- Set some other error message
    else
      -- perform some update on the model state
      model.deposit (id, amount)
    end
    -- Publish model update
    etf_cmd_container.on_change.notify ([Current])
  end
end
end
                
```



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Bank ATM

Separation of Concerns

Prototyping System with Abstract UI

Abstract Events: Bank ATM

ETF in a Nutshell

Workflow: Develop-Connect-Test

ETF: Abstract User Interface

ETF: Generating a New Project

ETF: Architecture

ETF: Input Errors

2019

Inheritance

Readings: OOSCS2 Chapters 14 – 16



EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG

Aspects of Inheritance



- **Code Reuse**
- Substitutability
 - **Polymorphism** and **Dynamic Binding** [compile-time type checks]
 - **Sub-contracting** [runtime behaviour checks]

2019

Why Inheritance: A Motivating Example



Problem: A *student management system* stores data about students. There are two kinds of university students: *resident* students and *non-resident* students. Both kinds of students have a *name* and a list of *registered courses*. Both kinds of students are restricted to *register* for no more than 30 courses. When *calculating the tuition* for a student, a base amount is first determined from the list of courses they are currently registered (each course has an associated fee). For a non-resident student, there is a *discount rate* applied to the base amount to waive the fee for on-campus accommodation. For a resident student, there is a *premium rate* applied to the base amount to account for the fee for on-campus accommodation and meals.

Tasks: Design classes that satisfy the above problem statement. At runtime, each type of student must be able to register a course and calculate their tuition fee.

2019

The COURSE Class



```
class
  COURSE

create -- Declare commands that can be used as constructors
  make

feature -- Attributes
  title: STRING
  fee: REAL

feature -- Commands
  make (t: STRING; f: REAL)
    -- Initialize a course with title 't' and fee 'f'.
    do
      title := t
      fee := f
    end
end
```

not fix

No Inheritance: NON_RESIDENT_STUDENT Class



```
class NON_RESIDENT_STUDENT
create make
feature -- Attributes
  name: STRING
  courses: LINKED_LIST[COURSE]
  discount_rate: REAL
feature -- Constructor
  make (n: STRING)
    do name := n ; create courses.make end
feature -- Commands
  set_dr (r: REAL) do discount_rate := r end
  register (c: COURSE) do courses.extend (c) end
feature -- Queries
  tuition: REAL
  local base: REAL
  do base := 0.0
    across courses as c loop base := base + c.item.fee end
  Result := base * discount_rate
end
end
```

not fix

No Inheritance: RESIDENT_STUDENT Class



```
class RESIDENT_STUDENT
create make
feature -- Attributes
  name: STRING
  courses: LINKED_LIST[COURSE]
  premium_rate: REAL
feature -- Constructor
  make (n: STRING)
    do name := n ; create courses.make end
feature -- Commands
  set_pr (r: REAL) do premium_rate := r end
  register (c: COURSE) do courses.extend (c) end
feature -- Queries
  tuition: REAL
  local base: REAL
  do base := 0.0
    across courses as c loop base := base + c.item.fee end
  Result := base * premium_rate
end
end
```

not fix

No Inheritance: Testing Student Classes



```
test_students: BOOLEAN
local
  c1, c2: COURSE
  jim: RESIDENT_STUDENT
  jeremy: NON_RESIDENT_STUDENT
do
  create c1.make ("EECS2030", 500.0)
  create c2.make ("EECS3311", 500.0)
  create jim.make ("J. Davis")
  jim.set_pr (1.25)
  jim.register (c1)
  jim.register (c2)
  Result := jim.tuition = 1250
  check Result end
  create jeremy.make ("J. Gibbons")
  jeremy.set_dr (0.75)
  jeremy.register (c1)
  jeremy.register (c2)
  Result := jeremy.tuition = 750
end
```

not fix

No Inheritance: Issues with the Student Classes



- Implementations for the two student classes seem to work. But can you see any potential problems with it?
- The code of the two student classes share a lot in common.
- **Duplicates of code make it hard to maintain your software!**
- This means that when there is a change of policy on the common part, we need modify *more than one places*.
⇒ This violates the **Single Choice Principle**:
when a **change** is needed, there should be **a single place** (or **a minimal number of places**) where you need to make that change.

10:16

No Inheritance: Maintainability of Code (2)



What if a **new** way for base tuition calculation is to be implemented?

e.g.,

```
tuition: REAL
  local base: REAL
  do base := 0.0
    across courses as c loop base := base + c.item.fee end
  Result := base * inflation_rate * ...
end
```

We need to change the `tuition` query in **both** student classes.

⇒ **Violation** of the **Single Choice Principle**

10:18

No Inheritance: Maintainability of Code (1)



What if a **new** way for course registration is to be implemented?

e.g.,

```
register(Course c)
do
  if courses.count >= MAX_CAPACITY then
    -- Error: maximum capacity reached.
  else
    courses.extend (c)
  end
end
```

We need to change the `register` commands in **both** student classes!

⇒ **Violation** of the **Single Choice Principle**

10:16

No Inheritance: A Collection of Various Kinds of Students



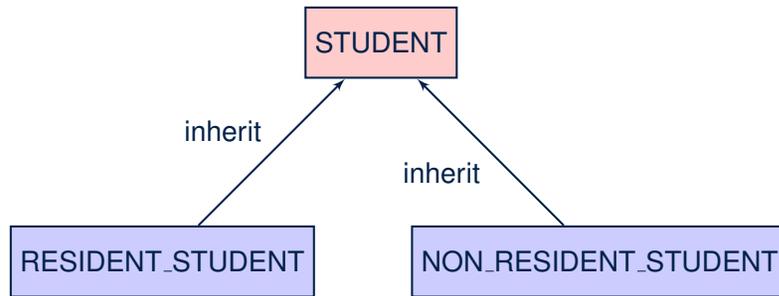
How do you define a class `StudentManagementSystem` that contains a list of **resident** and **non-resident** students?

```
class STUDENT_MANAGEMENT_SYSETM
  rs : LINKED_LIST[RESIDENT_STUDENT]
  nrs : LINKED_LIST[NON_RESIDENT_STUDENT]
  add_rs (rs: RESIDENT_STUDENT) do ... end
  add_nrs (nrs: NON_RESIDENT_STUDENT) do ... end
  register_all (Course c) -- Register a common course 'c'.
  do
    across rs as c loop c.item.register (c) end
    across nrs as c loop c.item.register (c) end
  end
end
```

But what if we later on introduce **more kinds of students**?
Inconvenient to handle each list of students, in pretty much the **same** manner, **separately**!

10:18

Inheritance Architecture



13/07/88

Inheritance: The RESIDENT_STUDENT Child Class



```
1 class
2   RESIDENT_STUDENT
3 inherit
4   STUDENT
5   redefine tuition end
6 create make
7 feature -- Attributes
8   premium_rate : REAL
9 feature -- Commands
10  set_pr (r: REAL) do premium_rate := r end
11 feature -- Queries
12  tuition: REAL
13    local base: REAL
14    do base := Precursor ; Result := base * premium_rate end
15 end
```

- L3: RESIDENT_STUDENT inherits all features from STUDENT.
- There is no need to repeat the register command
- L14: *Precursor* returns the value from query tuition in STUDENT.

13/07/88

Inheritance: The STUDENT Parent Class



```
1 class STUDENT
2 create make
3 feature -- Attributes
4   name: STRING
5   courses: LINKED_LIST[COURSE]
6 feature -- Commands that can be used as constructors.
7   make (n: STRING) do name := n ; create courses.make end
8 feature -- Commands
9   register (c: COURSE) do courses.extend (c) end
10 feature -- Queries
11  tuition: REAL
12    local base: REAL
13    do base := 0.0
14      across courses as c loop base := base + c.item.fee end
15    Result := base
16  end
17 end
```

13/07/88

Inheritance: The NON_RESIDENT_STUDENT Child Class

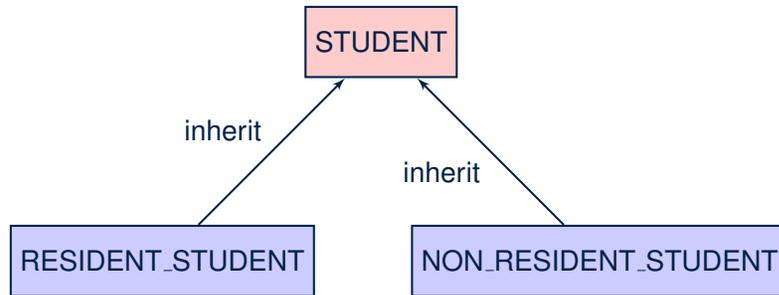


```
1 class
2   NON_RESIDENT_STUDENT
3 inherit
4   STUDENT
5   redefine tuition end
6 create make
7 feature -- Attributes
8   discount_rate : REAL
9 feature -- Commands
10  set_dr (r: REAL) do discount_rate := r end
11 feature -- Queries
12  tuition: REAL
13    local base: REAL
14    do base := Precursor ; Result := base * discount_rate end
15 end
```

- L3: NON_RESIDENT_STUDENT inherits all features from STUDENT.
- There is no need to repeat the register command
- L14: *Precursor* returns the value from query tuition in STUDENT.

13/07/88

Inheritance Architecture Revisited



- The class that defines the common features (attributes, commands, queries) is called the **parent**, **super**, or **ancestor** class.
- Each “specialized” class is called a **child**, **sub**, or **descendent** class.

19/07/88

Testing the Two Student Sub-Classes



```
test_students: BOOLEAN
local
  c1, c2: COURSE
  jim: RESIDENT_STUDENT ; jeremy: NON_RESIDENT_STUDENT
do
  create c1.make ("EECS2030", 500.0); create c2.make ("EECS3311", 500.0)
  create jim.make ("J. Davis")
  jim.set_pr (1.25) ; jim.register (c1); jim.register (c2)
  Result := jim.tuition = 1250
check Result end
create jeremy.make ("J. Gibbons")
jeremy.set_dr (0.75); jeremy.register (c1); jeremy.register (c2)
Result := jeremy.tuition = 750
end
```

- The software can be used in exactly the same way as before (because we did not modify **feature signatures**).
- But now the internal structure of code has been made **maintainable** using **inheritance**.

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Using Inheritance for Code Reuse



Inheritance in Eiffel (or any OOP language) allows you to:

- Factor out **common features** (attributes, commands, queries) in a separate class.
e.g., the `STUDENT` class
- Define an “specialized” version of the class which:
 - **inherits** definitions of all attributes, commands, and queries
e.g., attributes `name`, `courses`
e.g., command `register`
e.g., query on base amount in `tuition`
This means code reuse and elimination of code duplicates!
 - **defines new** features if necessary
e.g., `set_pr` for `RESIDENT_STUDENT`
e.g., `set_dr` for `NON_RESIDENT_STUDENT`
 - **redefines** features if necessary
e.g., compounded tuition for `RESIDENT_STUDENT`
e.g., discounted tuition for `NON_RESIDENT_STUDENT`

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Static Type vs. Dynamic Type



- In **object orientation**, an entity has two kinds of types:
 - **static type** is declared at compile time [**unchangeable**]
An entity's **ST** determines what features may be called upon it.
 - **dynamic type** is changeable at runtime

- In Java:

```
Student s = new Student("Alan");
Student rs = new ResidentStudent("Mark");
```

- In Eiffel:

```
local s: STUDENT
      rs: STUDENT
do create {STUDENT} s.make ("Alan")
   create {RESIDENT_STUDENT} rs.make ("Mark")
```

- In Eiffel, the **dynamic type** can be omitted if it is meant to be the same as the **static type**:

```
local s: STUDENT
do create s.make ("Alan")
```

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Inheritance Architecture Revisited

```

register (Course c)
tuition: REAL
class STUDENT
name: STRING
courses: LINKED_LIST[COURSE]

/* new features */
premium_rate: REAL
set_pr (r: REAL)
/* redefined features */
tuition: REAL
class RESIDENT_STUDENT
/* new features */
discount_rate: REAL
set_dr (r: REAL)
/* redefined features */
tuition: REAL
class NON_RESIDENT_STUDENT

```

```

s1, s2, s3: STUDENT ; rs: RESIDENT_STUDENT ; nrs: NON_RESIDENT_STUDENT
create {STUDENT} s1.make ("S1")
create {RESIDENT_STUDENT} s2.make ("S2")
create {NON_RESIDENT_STUDENT} s3.make ("S3")
create {RESIDENT_STUDENT} rs.make ("RS")
create {NON_RESIDENT_STUDENT} nrs.make ("NRS")

```

| | name | courses | reg | tuition | pr | set_pr | dr | set_dr |
|------|------|---------|-----|---------|----|--------|----|--------|
| s1. | | ✓ | | | | | × | |
| s2. | | ✓ | | | | | × | |
| s3. | | ✓ | | | | | × | |
| rs. | | ✓ | | | ✓ | | | × |
| nrs. | | ✓ | | | × | | | ✓ |

CRASH

Polymorphism: Intuition (2)

```

1 local s: STUDENT ; rs: RESIDENT_STUDENT
2 do create {STUDENT} s.make ("Stella")
3   create {RESIDENT_STUDENT} rs.make ("Rachael")
4   rs.set_pr (1.25)
5   s := rs /* Is this valid? */
6   rs := s /* Is this valid? */

```

- rs := s (L6) should be **invalid**:
- rs declared of type RESIDENT_STUDENT ∴ calling rs.set_pr(1.50) can be expected.
- rs is now pointing to a STUDENT object.
- Then, what would happen to rs.set_pr(1.50)?

CRASH ∴ rs.premium_rate is **undefined**!!

CRASH

Polymorphism: Intuition (1)

```

1 local
2   s: STUDENT
3   rs: RESIDENT_STUDENT
4 do
5   create s.make ("Stella")
6   create rs.make ("Rachael")
7   rs.set_pr (1.25)
8   s := rs /* Is this valid? */
9   rs := s /* Is this valid? */

```

- Which one of L8 and L9 is **valid**? Which one is **invalid**?
 - L8: What **kind** of address can s store? [STUDENT]
 - ∴ The context object s is **expected** to be used as:
 - s.register(eecs3311) and s.tuition
 - L9: What **kind** of address can rs store? [RESIDENT_STUDENT]
 - ∴ The context object rs is **expected** to be used as:
 - rs.register(eecs3311) and rs.tuition
 - rs.set_pr (1.50) [increase premium rate]

CRASH

Polymorphism: Intuition (3)

```

1 local s: STUDENT ; rs: RESIDENT_STUDENT
2 do create {STUDENT} s.make ("Stella")
3   create {RESIDENT_STUDENT} rs.make ("Rachael")
4   rs.set_pr (1.25)
5   s := rs /* Is this valid? */
6   rs := s /* Is this valid? */

```

- s := rs (L5) should be **valid**:
- Since s is declared of type STUDENT, a subsequent call s.set_pr(1.50) is **never** expected.
- s is now pointing to a RESIDENT_STUDENT object.
- Then, what would happen to s.tuition?

OK ∴ s.premium_rate is just **never used**!!

CRASH

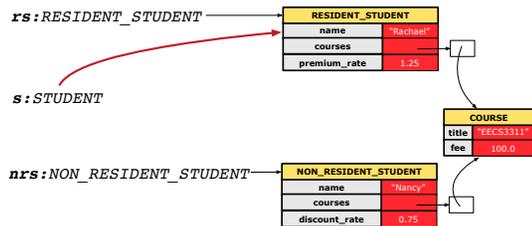
Dynamic Binding: Intuition (1)

```

1 local c : COURSE ; s : STUDENT
2 do crate c.make ("EECS3311", 100.0)
3   create {RESIDENT_STUDENT} rs.make("Rachael")
4   create {NON_RESIDENT_STUDENT} nrs.make("Nancy")
5   rs.set_pr(1.25); rs.register(c)
6   nrs.set_dr(0.75); nrs.register(c)
7   s := rs; ; check s.tuition = 125.0 end
8   s := nrs; ; check s.tuition = 75.0 end

```

After `s := rs` (L7), `s` points to a `RESIDENT_STUDENT` object.
 ⇒ Calling `s.tuition` applies the premium-rate.



DbC: Contract View of Supplier

Any potential **client** who is interested in learning about the kind of services provided by a **supplier** can look through the **contract view** (without showing any implementation details):

```

class ACCOUNT
create
  make
feature -- Attributes
  owner : STRING
  balance : INTEGER
feature -- Constructors
  make(nn: STRING; nb: INTEGER)
  require -- precondition
    positive_balance: nb > 0
  end
feature -- Commands
  withdraw(amount: INTEGER)
  require -- precondition
    non_negative_amount: amount > 0
    affordable_amount: amount <= balance -- problematic, why?
  ensure -- postcondition
    balance_deducted: balance = old balance - amount
  end
invariant -- class invariant
  positive_balance: balance > 0
end

```

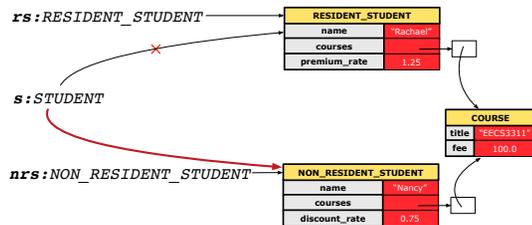
Dynamic Binding: Intuition (2)

```

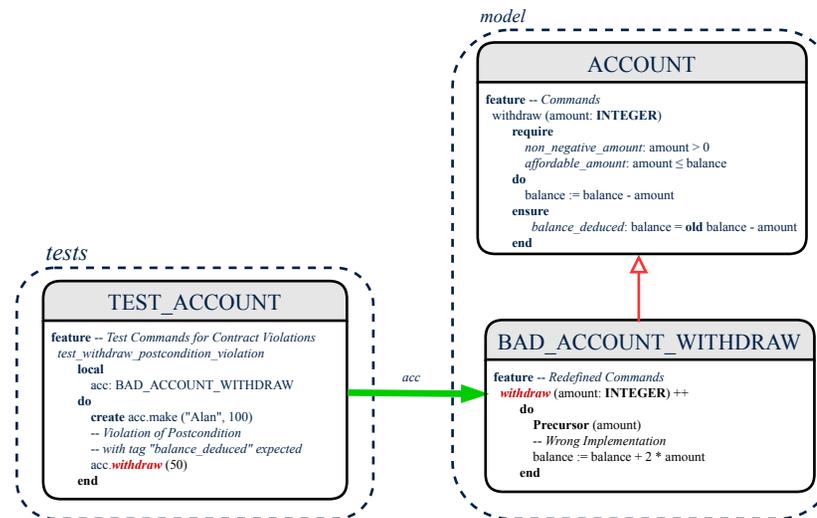
1 local c : COURSE ; s : STUDENT
2 do crate c.make ("EECS3311", 100.0)
3   create {RESIDENT_STUDENT} rs.make("Rachael")
4   create {NON_RESIDENT_STUDENT} nrs.make("Nancy")
5   rs.set_pr(1.25); rs.register(c)
6   nrs.set_dr(0.75); nrs.register(c)
7   s := rs; ; check s.tuition = 125.0 end
8   s := nrs; ; check s.tuition = 75.0 end

```

After `s := nrs` (L8), `s` points to a `NON_RESIDENT_STUDENT` object.
 ⇒ Calling `s.tuition` applies the discount-rate.



ES_TEST: Expecting to Fail Postcondition (1)



ES_TEST: Expecting to Fail Postcondition (2.1)



```
1 class
2   BAD_ACCOUNT_WITHDRAW
3 inherit
4   ACCOUNT
5   redefine withdraw end
6 create
7   make
8 feature -- redefined commands
9   withdraw(amount: INTEGER)
10  do
11    Precursor(amount)
12    -- Wrong implementation
13    balance := balance + 2 * amount
14  end
15 end
```

- L3-5: BAD_ACCOUNT_WITHDRAW.withdraw inherits postcondition from ACCOUNT.withdraw: balance = old balance - amount.
- L11 calls *correct* implementation from parent class ACCOUNT.
- L13 makes overall implementation *incorrect*.

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Exercise



Recall from the “Writing Complete Postconditions” lecture:

```
class BANK
  deposit_on_v5 (n: STRING; a: INTEGER)
  do ... -- Put Correct Implementation Here.
  ensure
  ...
  others_unchanged:
  across old accounts.deep.twin as cursor
  all cursor.item.owner /~ n implies
    cursor.item ~ account_of (cursor.item.owner)
  end
end
```

How do you create a “bad” descendant of BANK that violates this postcondition?

```
class BAD_BANK_DEPOSIT
  inherit BANK redefine deposit end
  feature -- redefined feature
  deposit_on_v5 (n: STRING; a: INTEGER)
  do Precursor (n, a)
    accounts[accounts.lower].deposit(a)
  end
end
```

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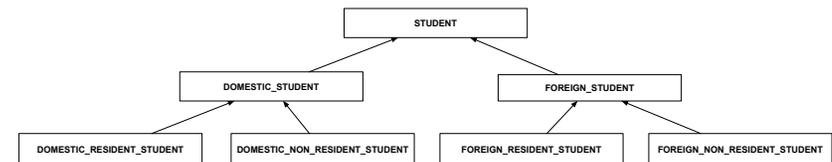
ES_TEST: Expecting to Fail Postcondition (2.2)



```
1 class TEST_ACCOUNT
2 inherit ES_TEST
3 create make
4 feature -- Constructor for adding tests
5   make
6   do
7     add_violation_case_with_tag ("balance_deducted",
8     agent test_withdraw_postcondition_violation)
9   end
10 feature -- Test commands (test to fail)
11   test_withdraw_postcondition_violation
12   local
13     acc: BAD_ACCOUNT_WITHDRAW
14   do
15     comment ("test: expected postcondition violation of withdraw")
16     create acc.make ("Alan", 100)
17     -- Postcondition Violation with tag "balance_deducted" to occur.
18     acc.withdraw (50)
19   end
20 end
```

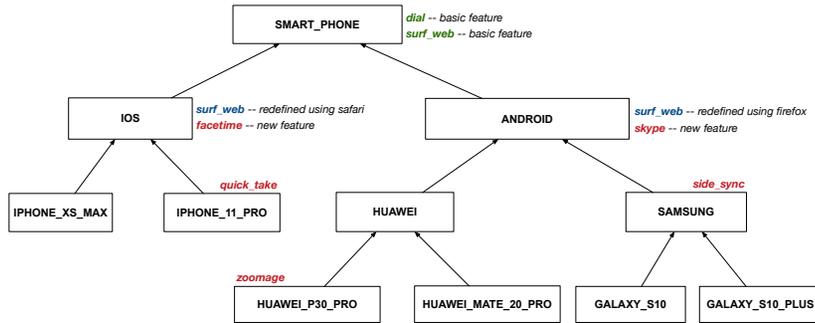
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Multi-Level Inheritance Architecture (1)



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Multi-Level Inheritance Architecture (2)



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Inheritance Forms a Type Hierarchy



- A (data) **type** denotes a set of related *runtime values*.
 - Every *class* can be used as a type: the set of runtime *objects*.
- Use of *inheritance* creates a **hierarchy** of classes:
 - (Implicit) Root of the hierarchy is ANY.
 - Each inherit declaration corresponds to an upward arrow.
 - The inherit relationship is *transitive*: when A inherits B and B inherits C, we say A *indirectly* inherits C.
 - e.g., Every class implicitly inherits the ANY class.
- Ancestor vs. Descendant** classes:
 - The **ancestor classes** of a class A are: A itself and all classes that A directly, or indirectly, inherits.
 - A inherits all features from its *ancestor classes*.
 - ∴ A's instances have a **wider range of expected usages** (i.e., attributes, queries, commands) than instances of its *ancestor* classes.
 - The **descendant classes** of a class A are: A itself and all classes that directly, or indirectly, inherits A.
 - Code defined in A is inherited to all its *descendant classes*.

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Inheritance Accumulates Code for Reuse



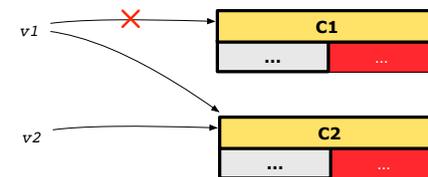
- The *lower* a class is in the type hierarchy, the *more code* it accumulates from its *ancestor classes*:
 - A *descendant class* inherits all code from its *ancestor classes*.
 - A *descendant class* may also:
 - Declare new attributes.
 - Define new queries or commands.
 - Redefine** inherited queries or commands.
- Consequently:
 - When being used as **context objects**, instances of a class' *descendant classes* have a **wider range of expected usages** (i.e., attributes, commands, queries).
 - When expecting an object of a particular class, we may **substitute** it with an object of any of its *descendant classes*.
 - e.g., When expecting a STUDENT object, substitute it with either a RESIDENT_STUDENT or a NON_RESIDENT_STUDENT object.
 - Justification:** A *descendant class* contains **at least as many** features as defined in its *ancestor classes* (but **not vice versa!**).

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Substitutions via Assignments



- By declaring $v1 : C1$, *reference variable* v1 will store the *address* of an object of class C1 at runtime.
- By declaring $v2 : C2$, *reference variable* v2 will store the *address* of an object of class C2 at runtime.
- Assignment $v1 := v2$ *copies the address* stored in v2 into v1.
 - v1 will instead point to wherever v2 is pointing to. [**object alias**]



- In such assignment $v1 := v2$, we say that we **substitute** an object of type C1 with an object of type C2.
- Substitutions** are subject to *rules!*

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Rules of Substitution

Given an inheritance hierarchy:

1. When expecting an object of class A, it is *safe* to **substitute** it with an object of any **descendant class** of A (including A).
 - e.g., When expecting an IOS phone, you *can* substitute it with either an IPHONE_XS_MAX or IPHONE_11_PRO.
 - ∴ Each **descendant class** of A is guaranteed to contain all code of (non-private) attributes, commands, and queries defined in A.
 - ∴ All features defined in A are *guaranteed to be available* in the new substitute.
2. When expecting an object of class A, it is *unsafe* to **substitute** it with an object of any **ancestor class of A's parent**.
 - e.g., When expecting an IOS phone, you *cannot* substitute it with just a SMART_PHONE, because the facetime feature is not supported in an ANDROID phone.
 - ∴ Class A may have defined new features that do not exist in any of its **parent's ancestor classes**.

REMARK



Reference Variable: Static Type

- A reference variable's **static type** is what we declare it to be.
 - e.g., `jim:STUDENT` declares jim's static type as STUDENT.
 - e.g., `my_phone:SMART_PHONE` declares a variable my_phone of static type SmartPhone.
 - The **static type** of a reference variable *never changes*.
- For a reference variable v, its **static type** `C` defines the **expected usages of v as a context object**.
- A feature call `v.m(...)` is **compilable** if `m` is defined in `C`.
 - e.g., After declaring `jim:STUDENT`, we
 - **may** call `register` and `tuition` on jim
 - **may not** call `set_pr` (specific to a resident student) or `set_dr` (specific to a non-resident student) on jim
 - e.g., After declaring `my_phone:SMART_PHONE`, we
 - **may** call `dial` and `surf_web` on my_phone
 - **may not** call `facetime` (specific to an IOS phone) or `skype` (specific to an Android phone) on my_phone

REMARK



Reference Variable: Dynamic Type

A reference variable's **dynamic type** is the type of object that it is currently pointing to at runtime.

- The **dynamic type** of a reference variable **may change** whenever we **re-assign** that variable to a different object.
- There are two ways to re-assigning a reference variable.

REMARK



Reference Variable: Changing Dynamic Type (1)

Re-assigning a reference variable to a newly-created object:

- **Substitution Principle**: the new object's class must be a **descendant class** of the reference variable's **static type**.
- e.g., Given the declaration `jim:STUDENT`:
 - `create {RESIDENT_STUDENT} jim.make("Jim")` changes the **dynamic type** of jim to RESIDENT_STUDENT.
 - `create {NON_RESIDENT_STUDENT} jim.make("Jim")` changes the **dynamic type** of jim to NON_RESIDENT_STUDENT.
- e.g., Given an alternative declaration `jim:RESIDENT_STUDENT`:
 - e.g., `create {STUDENT} jim.make("Jim")` is illegal because STUDENT is **not** a **descendant class** of the **static type** of jim (i.e., RESIDENT_STUDENT).

REMARK



Reference Variable: Changing Dynamic Type (2)

Re-assigning a reference variable v to an existing object that is referenced by another variable $other$ (i.e., $v := other$):

- **Substitution Principle**: the static type of $other$ must be a *descendant class* of v 's *static type*.
- e.g.,

```
jim: STUDENT ; rs: RESIDENT_STUDENT; nrs: NON_RESIDENT_STUDENT
create {STUDENT} jim.make (...)
create {RESIDENT_STUDENT} rs.make (...)
create {NON_RESIDENT_STUDENT} nrs.make (...)
```

- $rs := jim$ ×
- $nrs := jim$ ×
- $jim := rs$ ✓
changes the *dynamic type* of jim to the dynamic type of rs
- $jim := nrs$ ✓
changes the *dynamic type* of jim to the dynamic type of nrs

4.07.68



Polymorphism and Dynamic Binding (2.1)

```
1 test_polymorphism_students
2 local
3   jim: STUDENT
4   rs: RESIDENT_STUDENT
5   nrs: NON_RESIDENT_STUDENT
6 do
7   create {STUDENT} jim.make ("J. Davis")
8   create {RESIDENT_STUDENT} rs.make ("J. Davis")
9   create {NON_RESIDENT_STUDENT} nrs.make ("J. Davis")
10  jim := rs ✓
11  rs := jim ×
12  jim := nrs ✓
13  rs := jim ×
14 end
```

In (L3, L7), (L4, L8), (L5, L9), **ST = DT**, so we may abbreviate:

L7: `create jim.make ("J. Davis")`

L8: `create rs.make ("J. Davis")`

L9: `create nrs.make ("J. Davis")`

4.07.68



Polymorphism and Dynamic Binding (1)

- **Polymorphism**: An object variable may have “multiple possible shapes” (i.e., allowable *dynamic types*).
 - Consequently, there are *multiple possible versions* of each feature that may be called.
 - e.g., 3 possibilities of tuition on a *STUDENT* reference variable:
 - In *STUDENT*: base amount
 - In *RESIDENT_STUDENT*: base amount with `premium_rate`
 - In *NON_RESIDENT_STUDENT*: base amount with `discount_rate`
- **Dynamic binding**: When a feature m is called on an object variable, the version of m corresponding to its “current shape” (i.e., one defined in the *dynamic type* of m) will be called.

```
jim: STUDENT; rs: RESIDENT_STUDENT; nrs: NON_STUDENT
create {RESIDENT_STUDENT} rs.make (...)
create {NON_RESIDENT_STUDENT} nrs.nrs (...)
jim := rs
jim.tuition; /* version in RESIDENT_STUDENT */
jim := nrs
jim.tuition; /* version in NON_RESIDENT_STUDENT */
```

4.07.68



Polymorphism and Dynamic Binding (2.2)

```
test_dynamic_binding_students: BOOLEAN
local
  jim: STUDENT
  rs: RESIDENT_STUDENT
  nrs: NON_RESIDENT_STUDENT
  c: COURSE
do
  create c.make ("ECS3311", 500.0)
  create {STUDENT} jim.make ("J. Davis")
  create {RESIDENT_STUDENT} rs.make ("J. Davis")
  rs.register (c)
  rs.set_pr (1.5)
  jim := rs
  Result := jim.tuition = 750.0
  check Result end
  create {NON_RESIDENT_STUDENT} nrs.make ("J. Davis")
  nrs.register (c)
  nrs.set_dr (0.5)
  jim := nrs
  Result := jim.tuition = 250.0
end
```

4.07.68

Reference Type Casting: Motivation

```

1 local jim: STUDENT; rs: RESIDENT_STUDENT
2 do create {RESIDENT_STUDENT} jim.make ("J. Davis")
3   rs := jim
4   rs.setPremiumRate(1.5)

```

- **Line 2** is *legal*: `RESIDENT_STUDENT` is a *descendant class* of the static type of `jim` (i.e., `STUDENT`).
- **Line 3** is *illegal*: `jim`'s static type (i.e., `STUDENT`) is *not* a *descendant class* of `rs`'s static type (i.e., `RESIDENT_STUDENT`).
- Eiffel compiler is *unable to infer* that `jim`'s **dynamic type** in **Line 4** is `RESIDENT_STUDENT`. [**Undecidable**]
- Force the Eiffel compiler to believe so, by replacing **L3, L4** by a **type cast** (which **temporarily** changes the **ST** of `jim`):

```

check attached {RESIDENT_STUDENT} jim as rs_jim then
  rs := rs_jim
  rs.set_pr (1.5)
end

```

4.5.0.68

Notes on Type Cast (1)

- `check attached {C} y then ... end` **always compiles**
- What if `C` is not an **ancestor** of `y`'s **DT**?
⇒ A **runtime** assertion violation occurs!
∴ `y`'s **DT** cannot fulfill the expectation of `C`.

4.5.0.68

Reference Type Casting: Syntax

```

1 check attached {RESIDENT_STUDENT} jim as rs_jim then
2   rs := rs_jim
3   rs.set_pr (1.5)
4 end

```

L1 is an assertion:

- `attached RESIDENT_STUDENT jim` is a Boolean expression that is to be evaluated at **runtime**.
 - If it evaluates to **true**, then the `as rs_jim` expression has the effect of assigning "the cast version" of `jim` to a new variable `rs_jim`.
 - If it evaluates to **false**, then a runtime assertion violation occurs.
- **Dynamic Binding**: **Line 4** executes the correct version of `set_pr`.
- It is approximately the same as following Java code:

```

if(jim instanceof ResidentStudent) {
  ResidentStudent rs = (ResidentStudent) jim;
  rs.set_pr(1.5);
}
else { throw new Exception("Cast Not Done."); }

```

4.5.0.68

Notes on Type Cast (2)

- Given `v` of static type `ST`, it is **violation-free** to cast `v` to `C`, as long as `C` is a **descendant or ancestor class** of `ST`.
- Why Cast?
 - Without cast, we can **only** call features defined in `ST` on `v`.
 - By casting `v` to `C`, we create an **alias** of the object pointed by `v`, with the new **static type** `C`.
⇒ All features that are defined in `C` can be called.

```

my_phone: IOS
create {IPHONE_11_PRO} my_phone.make
-- can only call features defined in IOS on myPhone
-- dial, surf_web, facetime ✓ quick_take, skype, side_sync, zoomage ×
check attached {SMART_PHONE} my_phone as sp then
-- can now call features defined in SMART_PHONE on sp
-- dial, surf_web ✓ facetime, quick_take, skype, side_sync, zoomage ×
end
check attached {IPHONE_11_PRO} my_phone as ip11_pro then
-- can now call features defined in IPHONE_11_PRO on ip11_pro
-- dial, surf_web, facetime, quick_take ✓ skype, side_sync, zoomage ×
end

```

4.5.0.68

Notes on Type Cast (3)



A cast `check attached {C} v as ...` triggers an **assertion violation** if C is **not** along the **ancestor path** of v's **DT**.

```
test_smart_phone_type_cast_violation
local mine: ANDROID
do create {HUAWEI} mine.make
  -- ST of mine is ANDROID; DT of mine is HUAWEI
  check attached {SMART_PHONE} mine as sp then ... end
  -- ST of sp is SMART_PHONE; DT of sp is HUAWEI
  check attached {HUAWEI} mine as huawei then ... end
  -- ST of huawei is HUAWEI; DT of huawei is HUAWEI
  check attached {SAMSUNG} mine as samsung then ... end
  -- Assertion violation
  -- ∴ SAMSUNG is not ancestor of mine's DT (HUAWEI)
  check attached {HUAWEI_P30_PRO} mine as p30_pro then ... end
  -- Assertion violation
  -- ∴ HUAWEI_P30_PRO is not ancestor of mine's DT (HUAWEI)
end
```

START

Polymorphism: Feature Call Arguments (2)



```
test_polymorphism_feature_arguments
local
  s1, s2, s3: STUDENT
  rs: RESIDENT_STUDENT ; nrs: NON_RESIDENT_STUDENT
  sms: STUDENT_MANAGEMENT_SYSTEM
do
  create sms.make
  create {STUDENT} s1.make ("s1")
  create {RESIDENT_STUDENT} s2.make ("s2")
  create {NON_RESIDENT_STUDENT} s3.make ("s3")
  create {RESIDENT_STUDENT} rs.make ("rs")
  create {NON_RESIDENT_STUDENT} nrs.make ("nrs")
  sms.add_s (s1) ✓ sms.add_s (s2) ✓ sms.add_s (s3) ✓
  sms.add_s (rs) ✓ sms.add_s (nrs) ✓
  sms.add_rs (s1) × sms.add_rs (s2) × sms.add_rs (s3) ×
  sms.add_rs (rs) ✓ sms.add_rs (nrs) ×
  sms.add_nrs (s1) × sms.add_nrs (s2) × sms.add_nrs (s3) ×
  sms.add_nrs (rs) × sms.add_nrs (nrs) ✓
end
```

START

Polymorphism: Feature Call Arguments (1)



```
1 class STUDENT_MANAGEMENT_SYSTEM {
2   ss : ARRAY[STUDENT] -- ss[i] has static type Student
3   add_s (s: STUDENT) do ss[0] := s end
4   add_rs (rs: RESIDENT_STUDENT) do ss[0] := rs end
5   add_nrs (nrs: NON_RESIDENT_STUDENT) do ss[0] := nrs end
```

- **L4:** `ss[0] := rs` is valid. ∴ RHS's ST **RESIDENT_STUDENT** is a **descendant class** of LHS's ST **STUDENT**.
- Say we have a **STUDENT_MANAGEMENT_SYSTEM** object `sms`:
 - ∴ **call by value**, `sms.add_rs(o)` attempts the following assignment (i.e., replace parameter `rs` by a copy of argument `o`):

```
rs := o
```

- Whether this argument passing is valid depends on `o`'s **static type**.
- Rule:** In the signature of a feature `m`, if the type of a parameter is class `C`, then we may call feature `m` by passing objects whose **static types** are `C`'s **descendants**.

START

Why Inheritance: A Polymorphic Collection of Students



How do you define a class **STUDENT_MANAGEMENT_SYSTEM** that contains a list of **resident** and **non-resident** students?

```
class STUDENT_MANAGEMENT_SYSTEM
  students: LINKED_LIST[STUDENT]
  add_student(s: STUDENT)
  do
    students.extend (s)
  end
  registerAll (c: COURSE)
  do
    across
      students as s
    loop
      s.item.register (c)
    end
  end
end
```

START



Polymorphism and Dynamic Binding: A Polymorphic Collection of Students

```
test_sms_polymorphism: BOOLEAN
local
  rs: RESIDENT_STUDENT
  nrs: NON_RESIDENT_STUDENT
  c: COURSE
  sms: STUDENT_MANAGEMENT_SYSTEM
do
  create rs.make ("Jim")
  rs.set_pr (1.5)
  create nrs.make ("Jeremy")
  nrs.set_dr (0.5)
  create sms.make
  sms.add_s (rs)
  sms.add_s (nrs)
  create c.make ("EECS3311", 500)
  sms.register_all (c)
Result := sms.ss[1].tuition = 750 and sms.ss[2].tuition = 250
end
```

START



Polymorphism: Return Values (2)

```
1 test_sms_polymorphism: BOOLEAN
2 local
3   rs: RESIDENT_STUDENT ; nrs: NON_RESIDENT_STUDENT
4   c: COURSE ; sms: STUDENT_MANAGEMENT_SYSTEM
5 do
6   create rs.make ("Jim") ; rs.set_pr (1.5)
7   create nrs.make ("Jeremy") ; nrs.set_dr (0.5)
8   create sms.make ; sms.add_s (rs) ; sms.add_s (nrs)
9   create c.make ("EECS3311", 500) ; sms.register_all (c)
10  Result :=
11    get_student(1).tuition = 750
12    and get_student(2).tuition = 250
13 end
```

- L11: get_student (1) 's dynamic type? [RESIDENT_STUDENT]
- L11: Version of tuition? [RESIDENT_STUDENT]
- L12: get_student (2) 's dynamic type? [NON_RESIDENT_STUDENT]
- L12: Version of tuition? [NON_RESIDENT_STUDENT]

START



Polymorphism: Return Values (1)

```
1 class STUDENT_MANAGEMENT_SYSTEM {
2   ss: LINKED_LIST[STUDENT]
3   add_s (s: STUDENT)
4   do
5     ss.extend (s)
6   end
7   get_student(i: INTEGER): STUDENT
8   require 1 <= i and i <= ss.count
9   do
10    Result := ss[i]
11  end
12 end
```

- L2: **ST** of each stored item (ss[i]) in the list: [STUDENT]
- L3: **ST** of input parameter s: [STUDENT]
- L7: **ST** of return value (Result) of get_student: [STUDENT]
- L11: ss[i]'s **ST** is *descendant* of Result' **ST**.

Question: What can be the *dynamic type* of s after Line 11?

Answer: All descendant classes of Student.

START



Design Principle: Polymorphism

- When declaring an attribute $a: T$
 - ⇒ Choose **static type** T which "accumulates" all features that you predict you will want to call on a.
 - e.g., Choose $s: STUDENT$ if you do not intend to be specific about which kind of student s might be.
 - ⇒ Let **dynamic binding** determine at runtime which version of tuition will be called.
- What if after declaring $s: STUDENT$ you find yourself often needing to **cast** s to RESIDENT_STUDENT in order to access premium_rate?

```
check attached {RESIDENT_STUDENT} s as rs then rs.set_pr(...) end
```

⇒ Your design decision should have been: $s: RESIDENT_STUDENT$

- Same design principle applies to:
 - Type of feature parameters: $f(a: T)$
 - Type of queries: $q(...): T$

START



Static Type vs. Dynamic Type: When to consider which?

- **Whether or not an OOP code compiles** depends only on the **static types** of relevant variables.
 - ∴ Inferring the **dynamic type** statically is an **undecidable** problem that is inherently impossible to solve.
- **The behaviour of Eiffel code being executed at runtime**
 - e.g., which version of method is called
 - e.g., if a **check attached {...} as ... then ... end** assertion error will occur
 depends on the **dynamic types** of relevant variables.
 - ⇒ Best practice is to visualize how objects are created (by drawing boxes) and variables are re-assigned (by drawing arrows).

summary



Summary: Type Checking Rules

| CODE | CONDITION TO BE TYPE CORRECT |
|--|---|
| <code>x := y</code> | y's ST a descendant of x's ST |
| <code>x.f(y)</code> | Feature f defined in x's ST y's ST a descendant of f's parameter's ST |
| <code>z := x.f(y)</code> | Feature f defined in x's ST y's ST a descendant of f's parameter's ST ST of m's return value a descendant of z's ST |
| <code>check attached {C} y</code> | Always compiles |
| <code>check attached {C} y as temp then x := temp end</code> | C a descendant of x's ST |
| <code>check attached {C} y as temp then x.f(temp) end</code> | Feature f defined in x's ST C a descendant of f's parameter's ST |

Even if `check attached {C} y then ... end` always compiles,
a runtime assertion error occurs if C is not an **ancestor** of y's **DT!**

summary



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The State Design Pattern

Readings: OOSC2 Chapter 20



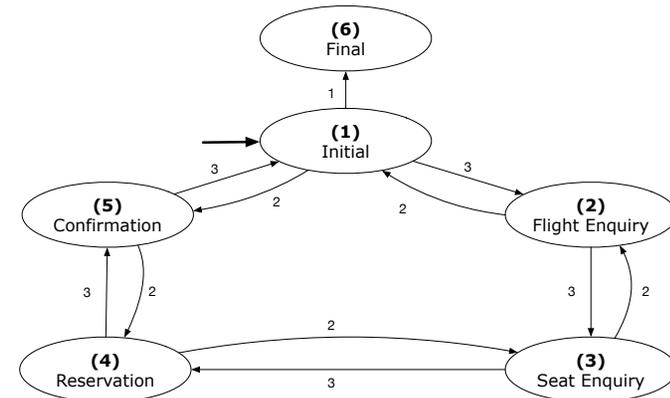
EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG

State Transition Diagram



Characterize **interactive system** as: **1)** A set of **states**; and **2)** For each state, its list of **applicable transitions** (i.e., actions).
e.g., Above reservation system as a **finite state machine**:



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Motivating Problem



Consider the reservation panel of an **online booking system**:

-- Enquiry on Flights --

Flight sought from: To:

Departure on or after: On or before:

Preferred airline (s):

Special requirements:

AVAILABLE FLIGHTS: 1
Flt#AA 42 Dep 8:25 Arr 7:45 Thru: Chicago

Choose next action:
0 - Exit
1 - Help
2 - Further enquiry
3 - Reserve a seat

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Design Challenges



1. The state-transition graph may **large** and **sophisticated**.
A large number N of states has $O(N^2)$ transitions
2. The graph structure is subject to **extensions/modifications**.
e.g., To merge "(2) Flight Enquiry" and "(3) Seat Enquiry":
Delete the state "(3) Seat Enquiry".
Delete its 4 incoming/outgoing transitions.
e.g., Add a new state "Dietary Requirements"
3. A **general solution** is needed for such **interactive systems**.
e.g., **taobao**, **eBay**, **amazon**, etc.

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A First Attempt



```

1.Initial_panel:
-- Actions for Label 1.
2.Flight_Enquiry_panel:
-- Actions for Label 2.
3.Seat_Enquiry_panel:
-- Actions for Label 3.
4.Reservation_panel:
-- Actions for Label 4.
5.Confirmation_panel:
-- Actions for Label 5.
6.Final_panel:
-- Actions for Label 6.
    
```

```

3.Seat_Enquiry_panel:
from
  Display Seat Enquiry Panel
until
  not (wrong answer or wrong choice)
do
  Read user's answer for current panel
  Read user's choice [C] for next step
  if wrong answer or wrong choice then
    Output error messages
  end
end
end
Process user's answer
case [C] in
  2: goto 2.Flight_Enquiry_panel
  3: goto 4.Reservation_panel
end
    
```

not

A First Attempt: Good Design?



- Runtime execution \approx a **“bowl of spaghetti”**.
 \Rightarrow The system's behaviour is hard to predict, trace, and debug.
- **Transitions** hardwired as system's **central control structure**.
 \Rightarrow The system is vulnerable to changes/additions of states/transitions.
- All labelled blocks are largely similar in their code structures.
 \Rightarrow This design **“smells”** due to duplicates/repetitions!
- The branching structure of the design exactly corresponds to that of the specific **transition graph**.
 \Rightarrow The design is **application-specific** and **not reusable** for other interactive systems.

not

A Top-Down, Hierarchical Solution



- **Separation of Concern** Declare the **transition table** as a feature the system, rather than its central control structure:

```

transition (src: INTEGER; choice: INTEGER): INTEGER
-- Return state by taking transition 'choice' from 'src' state.
require valid_source_state: 1 ≤ src ≤ 6
       valid_choice: 1 ≤ choice ≤ 3
ensure valid_target_state: 1 ≤ Result ≤ 6
    
```

- We may implement transition via a 2-D array.

| SRC STATE | CHOICE | | |
|--------------------|--------|---|---|
| | 1 | 2 | 3 |
| 1 (Initial) | 6 | 5 | 2 |
| 2 (Flight Enquiry) | – | 1 | 3 |
| 3 (Seat Enquiry) | – | 2 | 4 |
| 4 (Reservation) | – | 3 | 5 |
| 5 (Confirmation) | – | 4 | 1 |
| 6 (Final) | – | – | – |

| state | choice | | |
|-------|--------|---|---|
| | 1 | 2 | 3 |
| 1 | 6 | 5 | 2 |
| 2 | | 1 | 3 |
| 3 | | 2 | 4 |
| 4 | | 3 | 5 |
| 5 | | 4 | 1 |
| 6 | | | |

not

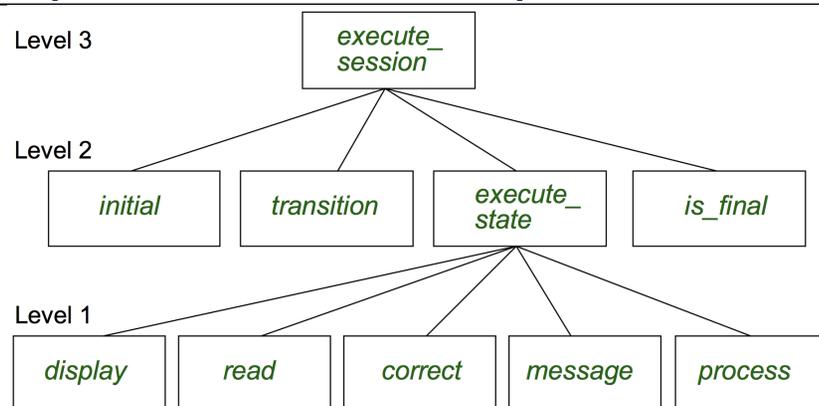
Hierarchical Solution: Good Design?



- This is a more general solution.
 \therefore **State transitions** are **separated** from the system's **central control structure**.
 \Rightarrow **Reusable** for another interactive system by making changes only to the **transition** feature.
- How does the **central control structure** look like in this design?

not

Hierarchical Solution: Top-Down Functional Decomposition



Modules of `execute_session` and `execute_state` are general enough on their *control structures*. ⇒ reusable

12/23/20

Hierarchical Solution: State Handling (1)

The following *control pattern* handles all states:

```

execute_state (current_state : INTEGER) : INTEGER
-- Handle interaction at the current state.
-- Return user's exit choice.
local
  answer: ANSWER; valid_answer: BOOLEAN; choice: INTEGER
do
  from
  until
    valid_answer
  do
    display(current_state)
    answer := read_answer(current_state)
    choice := read_choice(current_state)
    valid_answer := correct(current_state, answer)
    if not valid_answer then message(current_state, answer)
  end
  process(current_state, answer)
  Result := choice
end
  
```

12/23/20

Hierarchical Solution: System Control

All interactive sessions share the following *control pattern*:

- Start with some *initial state*.
- Repeatedly make *state transitions* (based on *choices* read from the user) until the state is *final* (i.e., the user wants to exit).

```

execute_session
-- Execute a full interactive session.
local
  current_state, choice: INTEGER
do
  from
  until
    is_final(current_state)
  do
    choice := execute_state(current_state)
    current_state := transition(current_state, choice)
  end
end
  
```

12/23/20

Hierarchical Solution: State Handling (2)

| FEATURE CALL | FUNCTIONALITY |
|---------------------------------|--|
| <code>display(s)</code> | Display screen outputs associated with <i>state s</i> |
| <code>read_answer(s)</code> | Read user's input for answers associated with <i>state s</i> |
| <code>read_choice(s)</code> | Read user's input for exit choice associated with <i>state s</i> |
| <code>correct(s, answer)</code> | Is the user's <i>answer</i> valid w.r.t. <i>state s</i> ? |
| <code>process(s, answer)</code> | Given that user's <i>answer</i> is valid w.r.t. <i>state s</i> , process it accordingly. |
| <code>message(s, answer)</code> | Given that user's <i>answer</i> is not valid w.r.t. <i>state s</i> , display an error message accordingly. |

Q: How similar are the code structures of the above state-dependant commands or queries?

12/23/20

Hierarchical Solution: State Handling (3)



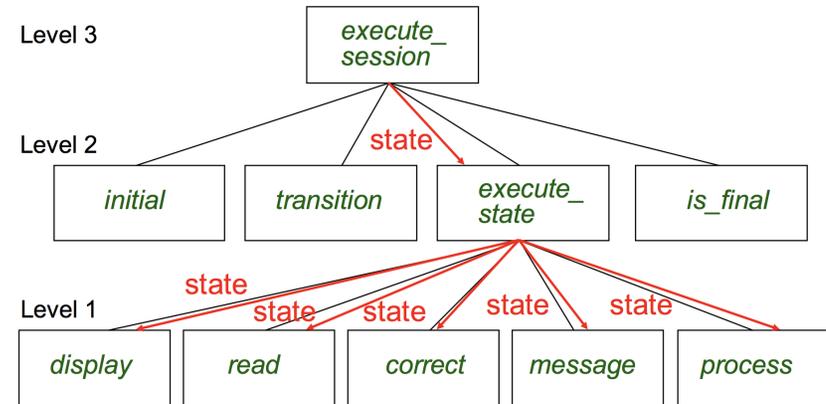
A: Actions of all such state-dependant features must **explicitly discriminate** on the input state argument.

```
display(current_state: INTEGER)
  require
    valid_state: 1 ≤ current_state ≤ 6
  do
    if current_state = 1 then
      -- Display Initial Panel
    elseif current_state = 2 then
      -- Display Flight Enquiry Panel
    ...
    else
      -- Display Final Panel
    end
  end
end
```

- Such design **smells!**
∴ Same list of conditional repeats for **all** state-dependant features.
- Such design **violates** the **Single Choice Principle**.
e.g., To add/delete a state ⇒ Add/delete a branch in all such features.

13/04/00

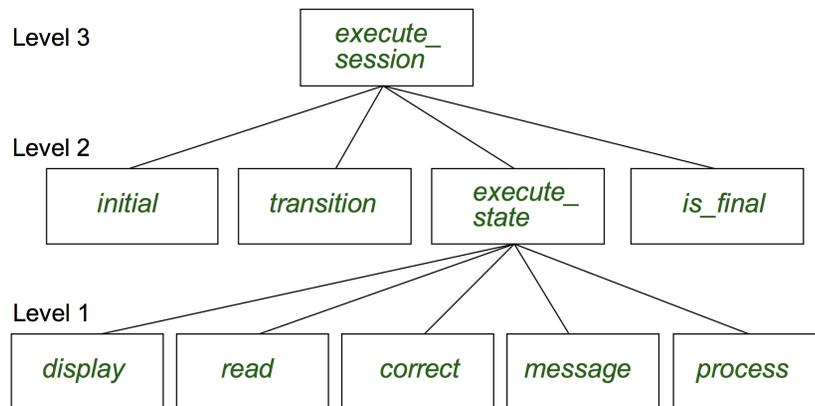
Hierarchical Solution: Pervasive States



- Too much data transmission: `current_state` is passed
- From `execute_session` (Level 3) to `execute_state` (Level 2)
 - From `execute_state` (Level 2) to all features at Level 1

13/04/00

Hierarchical Solution: Visible Architecture



13/04/00

Law of Inversion



If your routines exchange too many data, then put your routines in your data.

e.g.,

`execute_state` (Level 2) and all features at Level 1:

- Pass around (as **inputs**) the notion of **current_state**
- Build upon (via **discriminations**) the notion of **current_state**

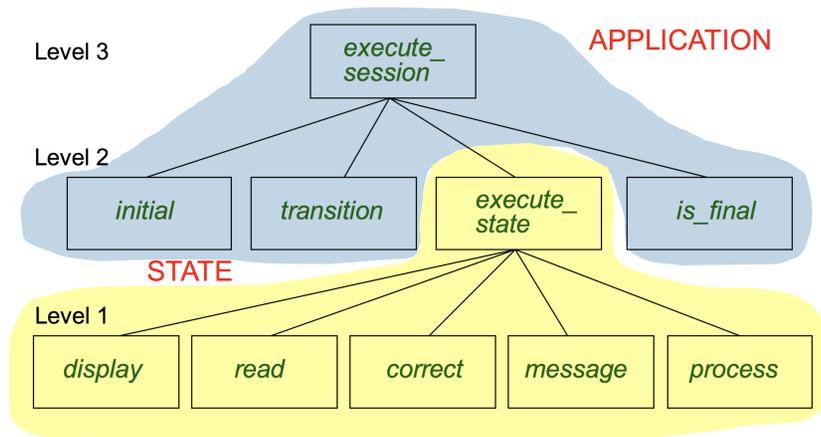
```

execute_state (s: INTEGER)
display (s: INTEGER)
read_answer (s: INTEGER)
read_choice (s: INTEGER)
correct (s: INTEGER; answer: ANSWER)
process (s: INTEGER; answer: ANSWER)
message (s: INTEGER; answer: ANSWER)
  
```

- ⇒ **Modularize** the notion of state as **class STATE**.
- ⇒ **Encapsulate** state-related information via a **STATE** interface.
- ⇒ Notion of **current_state** becomes **implicit**: the `Current` class.

13/04/00

Grouping by Data Abstractions



2013/01

The STATE ADT

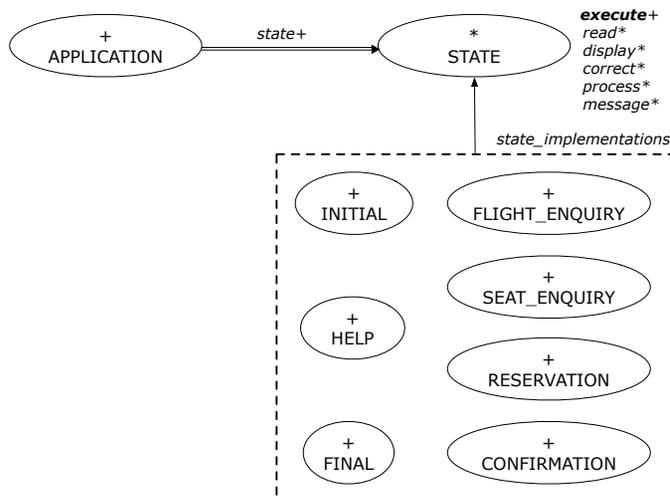


```
deferred class STATE
  read
  -- Read user's inputs
  -- Set 'answer' and 'choice'
  deferred end
  answer: ANSWER
  -- Answer for current state
  choice: INTEGER
  -- Choice for next step
  display
  -- Display current state
  deferred end
  correct: BOOLEAN
  deferred end
  process
  require correct
  deferred end
  message
  require not correct
  deferred end
```

```
execute
  local
    good: BOOLEAN
  do
    from
    until
      good
    loop
      display
      -- set answer and choice
      read
      good := correct
      if not good then
        message
      end
    end
  end
  process
end
end
```

2013/01

Architecture of the State Pattern



2013/01

The Template Design Pattern



Consider the following fragment of Eiffel code:

```
1 s: STATE
2 create {SEAT_ENQUIRY} s.make
3 s.execute
4 create {CONFIRMATION} s.make
5 s.execute
```

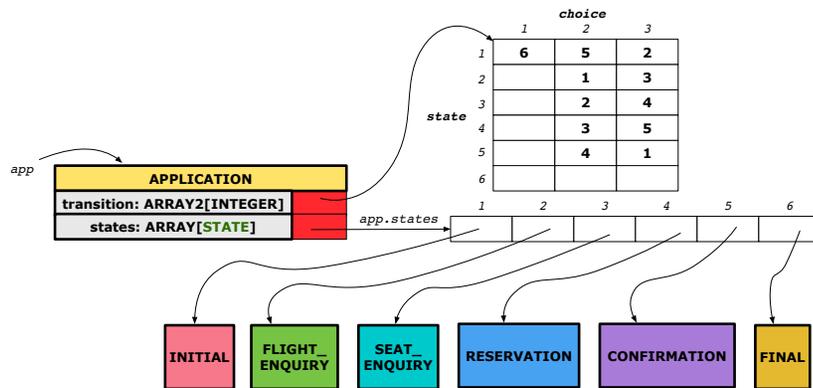
L2 and **L4**: the same version of effective feature execute (from the deferred class STATE) is called. [**template**]

L2: specific version of effective features display, process, etc., (from the effective descendant class SEAT_ENQUIRY) is called. [**template instantiated for SEAT_ENQUIRY**]

L4: specific version of effective features display, process, etc., (from the effective descendant class CONFIRMATION) is called. [**template instantiated for CONFIRMATION**]

2013/01

APPLICATION Class: Array of STATE



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APPLICATION Class (2)



```
class APPLICATION
feature {NONE} -- Implementation of Transition Graph
transition: ARRAY2[INTEGER]
states: ARRAY[STATE]
feature
put_state(s: STATE; index: INTEGER)
require 1 ≤ index ≤ number_of_states
do states.force(s, index) end
choose_initial(index: INTEGER)
require 1 ≤ index ≤ number_of_states
do initial := index end
put_transition(tar, src, choice: INTEGER)
require
1 ≤ src ≤ number_of_states
1 ≤ tar ≤ number_of_states
1 ≤ choice ≤ number_of_choices
do
transition.put(tar, src, choice)
end
end
```

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APPLICATION Class (1)



```
class APPLICATION create make
feature {NONE} -- Implementation of Transition Graph
transition: ARRAY2[INTEGER]
-- State transitions: transition[state, choice]
states: ARRAY[STATE]
-- State for each index, constrained by size of 'transition'
feature
initial: INTEGER
number_of_states: INTEGER
number_of_choices: INTEGER
make(n, m: INTEGER)
do number_of_states := n
number_of_choices := m
create transition.make_filled(0, n, m)
create states.make_empty
end
invariant
transition.height = number_of_states
transition.width = number_of_choices
end
```

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Example Test: Non-Interactive Session



```
test_application: BOOLEAN
local
app: APPLICATION ; current_state: STATE ; index: INTEGER
do
create app.make (6, 3)
app.put_state (create {INITIAL}.make, 1)
-- Similarly for other 5 states.
app.choose_initial (1)
-- Transit to FINAL given current state INITIAL and choice 1.
app.put_transition (6, 1, 1)
-- Similarly for other 10 transitions.

index := app.initial
current_state := app.states [index]
Result := attached {INITIAL} current_state
check Result end
-- Say user's choice is 3: transit from INITIAL to FLIGHT_STATUS
index := app.transition.item (index, 3)
current_state := app.states [index]
Result := attached {FLIGHT_ENQUIRY} current_state
end
```

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APPLICATION Class (3): Interactive Session



```
class APPLICATION
feature {NONE} -- Implementation of Transition Graph
  transition: ARRAY2[INTEGER]
  states: ARRAY[STATE]
feature
  execute_session
  local
    current_state: STATE
    index: INTEGER
  do
    from
      index := initial
    until
      is_final (index)
    loop
      current_state := states[index] -- polymorphism
      current_state.execute -- dynamic binding
      index := transition.item (index, current_state.choice)
    end
  end
end
end
```

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Top-Down, Hierarchical vs. OO Solutions



- In the second (top-down, hierarchy) solution, it is required for every state-related feature to *explicitly* and *manually* discriminate on the argument value, via a list of conditionals. e.g., Given `display(current_state: INTEGER)`, the calls `display(1)` and `display(2)` behave differently.
- The third (OO) solution, called the State Pattern, makes such conditional *implicit* and *automatic*, by making STATE as a deferred class (whose descendants represent all types of states), and by delegating such conditional actions to *dynamic binding*. e.g., Given `s: STATE`, behaviour of the call `s.display` depends on the *dynamic type* of `s` (such as INITIAL vs. FLIGHT_ENQUIRY).

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Building an Application



- Create instances of STATE.

```
s1: STATE
create {INITIAL} s1.make
```

- Initialize an APPLICATION.

```
create app.make(number_of_states, number_of_choices)
```

- Perform polymorphic assignments on app.states.

```
app.put_state(initial, 1)
```

- Choose an initial state.

```
app.choose_initial(1)
```

- Build the transition table.

```
app.put_transition(6, 1, 1)
```

- Run the application.

```
app.execute_session
```

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The Composite Design Pattern

EECS3311 A: Software Design
Winter 2020



CHEN-WEI WANG

Motivating Problem (1)



- Many manufactured systems, such as computer systems or stereo systems, are composed of **individual components** and **sub-systems** that contain components.
 - e.g., A computer system is composed of:
 - Individual pieces of equipment (*hard drives, cd-rom drives*)
Each equipment has **properties**: e.g., power consumption and cost.
 - Composites such as *cabinets, busses, and chassis*
Each *cabinet* contains various types of *chassis*, each of which in turn containing components (*hard-drive, power-supply*) and *busses* that contain *cards*.
- Design a system that will allow us to easily **build** systems and **calculate** their total cost and power consumption.

2 of 3

Motivating Problem (2)

Design for *tree structures* with whole-part *hierarchies*.

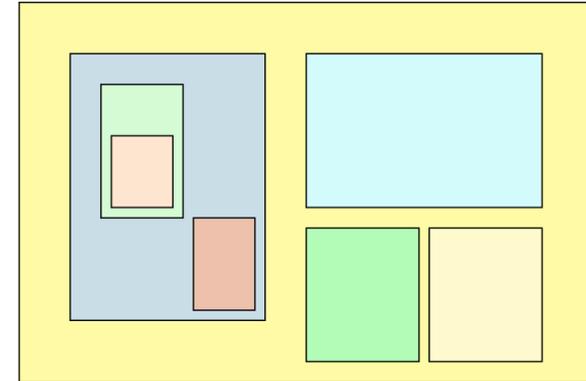


Challenge: There are *base* and *recursive* modelling artifacts.

not

MI: Combining Abstractions (2.1)

Q: How do you design class(es) for nested windows?

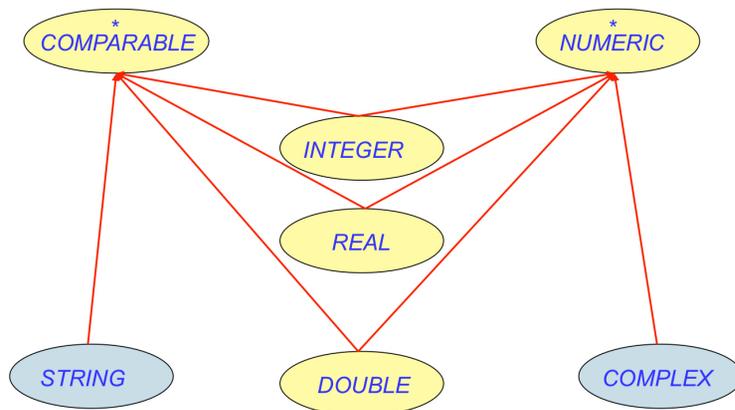


Hints: height, width, xpos, ypos, change width, change height, move, parent window, descendant windows, add child window

not

Multiple Inheritance: Combining Abstractions (1)

A class may have two more parent classes.



not

MI: Combining Abstractions (2)

A: Separating *Graphical* features and *Hierarchical* features

```
class RECTANGLE
  feature -- Queries
    width, height: REAL
    xpos, ypos: REAL
  feature -- Commands
    make (w, h: REAL)
    change_width
    change_height
    move
end
```

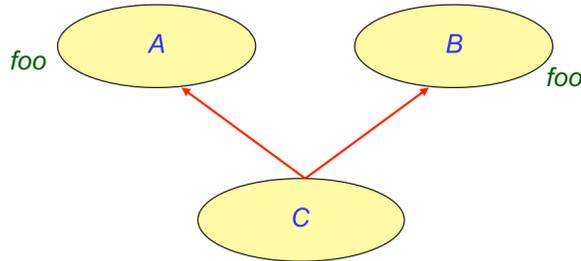
```
class TREE[G]
  feature -- Queries
    descendants: ITERABLE[G]
  feature -- Commands
    add (c: G)
    -- Add a child 'c'.
end
```

```
class WINDOW
  inherit
    RECTANGLE
    TREE[WINDOW]
end
```

```
test_window: BOOLEAN
local w1, w2, w3, w4: WINDOW
do
  create w1.make(8, 6) ; create w2.make(4, 3)
  create w3.make(1, 1) ; create w4.make(1, 1)
  w2.add(w4) ; w1.add(w2) ; w1.add(w3)
  Result := w1.descendants.count = 2
end
```

not

MI: Name Clashes



In class C, feature `foo` inherited from ancestor class A clashes with feature `foo` inherited from ancestor class B.

10:11

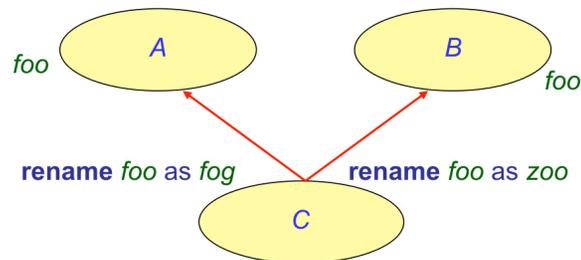
Solution: The Composite Pattern



- **Design**: Categorize into *base* artifacts or *recursive* artifacts.
- **Programming**: Build a *tree structure* representing the whole-part *hierarchy*.
- **Runtime**: Allow clients to treat *base* objects (leaves) and *recursive* compositions (nodes) *uniformly*.
 - ⇒ **Polymorphism**: *leaves* and *nodes* are “substitutable”.
 - ⇒ **Dynamic Binding**: Different versions of the same operation is applied on *individual objects* and *composites*. e.g., Given `e: EQUIPMENT`:
 - `e.price` may return the unit price of a *DISK_DRIVE*.
 - `e.price` may sum prices of a *CHASSIS*' containing equipments.

10:11

MI: Resolving Name Clashes

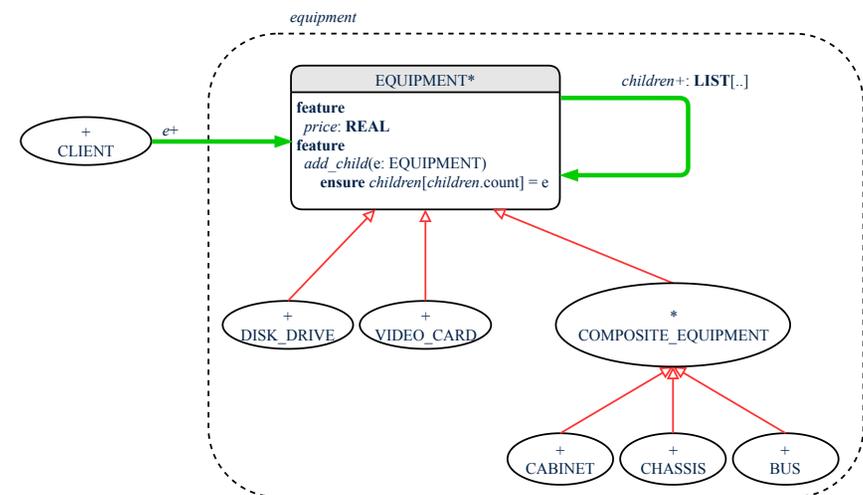


```
class C
  inherit
  A rename foo as fog end
  B rename foo as zoo end
  ...
```

| | o.foo | o.fog | o.zoo |
|------|-------|-------|-------|
| o: A | ✓ | ✗ | ✗ |
| o: B | ✓ | ✗ | ✗ |
| o: C | ✗ | ✓ | ✓ |

10:11

Composite Architecture: Design (1.1)



10:11

Composite Architecture: Design (1.2)



Q: Any flaw of this first design?

A: Two “composite” features defined at the EQUIPMENT level:

- children: LIST[EQUIPMENT]
- add(child: EQUIPMENT)

⇒ Inherited to all *base* equipments (e.g., HARD_DRIVE) that do not apply to such features.

1.018

Implementing the Composite Pattern (1)

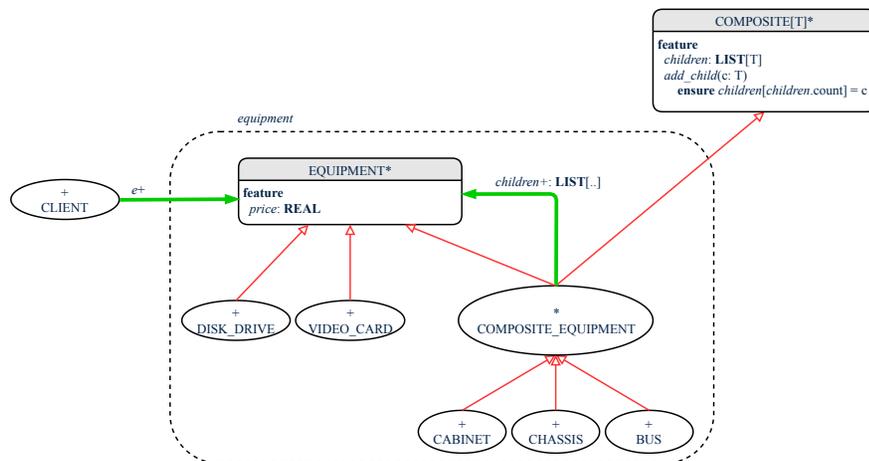


```
deferred class
  EQUIPMENT
  feature
    name: STRING
    price: REAL -- uniform access principle
  end
```

```
class
  CARD
  inherit
    EQUIPMENT
  feature
    make (n: STRING; p: REAL)
    do
      name := n
      price := p -- price is an attribute
    end
  end
```

1.018

Composite Architecture: Design (2.1)



1.018

Implementing the Composite Pattern (2.1)



```
deferred class
  COMPOSITE[T]
  feature
    children: LINKED_LIST[T]

  add (c: T)
  do
    children.extend (c) -- Polymorphism
  end
end
```

Exercise: Make the COMPOSITE class *iterable*.

1.018

Implementing the Composite Pattern (2.2)



```
class
  COMPOSITE_EQUIPMENT
inherit
  EQUIPMENT
  COMPOSITE [EQUIPMENT]
create
  make
feature
  make (n: STRING)
  do name := n ; create children.make end
  price : REAL -- price is a query
  -- Sum the net prices of all sub-equipments
  do
    across
      children as cursor
    loop
      Result := Result + cursor.item.price -- dynamic binding
    end
  end
end
```

15/018

Testing the Composite Pattern



```
test_composite_equipment: BOOLEAN
local
  card, drive: EQUIPMENT
  cabinet: CABINET -- holds a CHASSIS
  chassis: CHASSIS -- contains a BUS and a DISK_DRIVE
  bus: BUS -- holds a CARD
do
  create {CARD} card.make("16Mbs Token Ring", 200)
  create {DISK_DRIVE} drive.make("500 GB harddrive", 500)
  create bus.make("MCA Bus")
  create chassis.make("PC Chassis")
  create cabinet.make("PC Cabinet")

  bus.add(card)
  chassis.add(bus)
  chassis.add(drive)
  cabinet.add(chassis)
  Result := cabinet.price = 700
end
```

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The Visitor Design Pattern



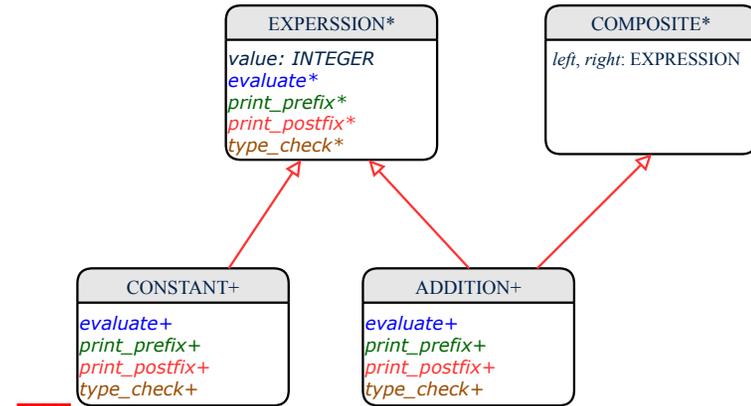
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Motivating Problem (2)

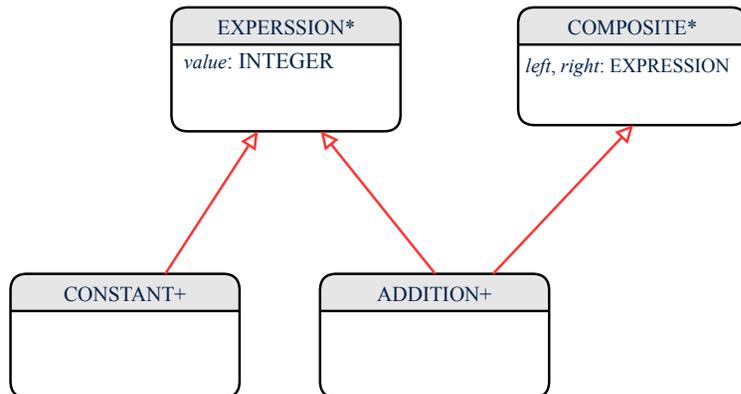
Extend the **composite pattern** to support **operations** such as evaluate, pretty printing (print_prefix, print_postfix), and type-check.



Motivating Problem (1)



Based on the **composite pattern** you learned, design classes to model **structures** of arithmetic expressions (e.g., 341 , 2 , $341 + 2$).



Problems of Extended Composite Pattern



- Distributing the various **unrelated operations** across nodes of the **abstract syntax tree** violates the **single-choice principle**:
To add/delete/modify an operation
⇒ Change of all descendants of EXPRESSION
- Each node class lacks in **cohesion**:
A **class** is supposed to group **relevant** concepts in a **single** place.
⇒ Confusing to mix codes for evaluation, pretty printing, and type checking.
⇒ We want to avoid “polluting” the classes with these various unrelated operations.

Open/Closed Principle



Software entities (classes, features, etc.) should be **open** for **extension**, but **closed** for **modification**.

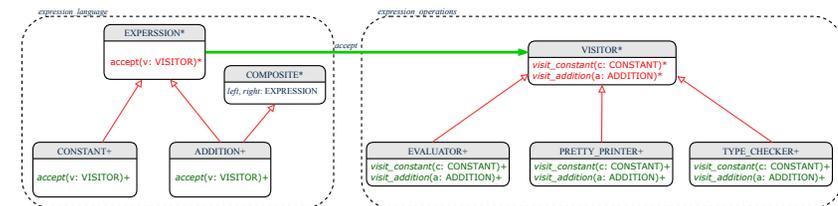
⇒ When **extending** the behaviour of a system, we:

- May add/modify the **open** (unstable) part of system.
- May not add/modify the **closed** (stable) part of system.

e.g., In designing the application of an expression language:

- **ALTERNATIVE 1:**
Syntactic constructs of the language may be **open**, whereas operations on the language may be **closed**.
- **ALTERNATIVE 2:**
Syntactic constructs of the language may be **closed**, whereas operations on the language may be **open**.

Visitor Pattern: Architecture



Visitor Pattern



• Separation of concerns :

- Set of language constructs
- Set of operations

⇒ Classes from these two sets are **decoupled** and organized into two separate clusters.

• Open-Closed Principle (OCP) : [ALTERNATIVE 2]

- **Closed**, stable part of system: set of language constructs
- **Open**, unstable part of system: set of operations

⇒ **OCP** helps us determine if Visitor Pattern is **applicable**.

⇒ If it was decided that language constructs are **open** and operations are **closed**, then do **not** use Visitor Pattern.

Visitor Pattern Implementation: Structures



Cluster **expression language**

- Declare **deferred** feature `accept(v: VISITOR)` in EXPRESSION.
- Implement `accept` feature in each of the descendant classes.

```
class CONSTANT inherit EXPRESSION
...
accept(v: VISITOR)
do
  v.visit_constant(Current)
end
end
```

```
class ADDITION
inherit EXPRESSION COMPOSITE
...
accept(v: VISITOR)
do
  v.visit_addition(Current)
end
end
```

Visitor Pattern Implementation: Operations



Cluster *expression_operations*

- For each descendant class C of EXPRESSION, declare a *deferred* feature `visit_c (e: C)` in the *deferred* class VISITOR.

```
deferred class VISITOR
  visit_constant(c: CONSTANT) deferred end
  visit_addition(a: ADDITION) deferred end
end
```

- Each descendant of VISITOR denotes a kind of operation.

```
class EVALUATOR inherit VISITOR
  value: INTEGER
  visit_constant(c: CONSTANT) do value := c.value end
  visit_addition(a: ADDITION)
    local eval_left, eval_right: EVALUATOR
    do a.left.accept(eval_left)
      a.right.accept(eval_right)
      value := eval_left.value + eval_right.value
    end
end
```

MARK

To Use or Not to Use the Visitor Pattern



- In the architecture of visitor pattern, what kind of *extensions* is easy and hard? Language structure? Language Operation?
 - Adding a new kind of *operation* element is easy.
 - To introduce a new operation for generating C code, we only need to introduce a new descendant class `C_CODE_GENERATOR` of VISITOR, then implement how to handle each language element in that class.
 - ⇒ *Single Choice Principle* is *obeyed*.
 - Adding a new kind of *structure* element is hard.
 - After adding a descendant class MULTIPLICATION of EXPRESSION, every concrete visitor (i.e., descendant of VISITOR) must be amended to provide a new `visit_multiplication` operation.
 - ⇒ *Single Choice Principle* is *violated*.
- The applicability of the visitor pattern depends on to what extent the *structure* will change.
 - ⇒ Use visitor if *operations* applied to *structure* change often.
 - ⇒ Do not use visitor if the *structure* changes often.

MARK

Testing the Visitor Pattern



```
1 test_expression_evaluation: BOOLEAN
2   local add, c1, c2: EXPRESSION ; v: VISITOR
3   do
4     create {CONSTANT} c1.make (1) ; create {CONSTANT} c2.make (2)
5     create {ADDITION} add.make (c1, c2)
6     create {EVALUATOR} v.make
7     add.accept (v)
8     check_attached {EVALUATOR} v as eval then
9       Result := eval.value = 3
10    end
11  end
```

Double Dispatch in Line 7:

- DT* of add is *ADDITION* ⇒ Call accept in *ADDITION*

```
v.visit_addition (add)
```

- DT* of v is *EVALUATOR* ⇒ Call visit_addition in *EVALUATOR*

```
visiting result of add.left + visiting result of add.right
```

MARK

Beyond this Lecture . . .



Learn about implementing the Composite and Visitor Patterns, from scratch, in this tutorial series:

<https://www.youtube.com/playlist?list=PL5dxAmCmjv4z5eXGW-ZBqsS2WZTyBHY2>

MARK

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12/12/20

Observer Design Pattern Event-Driven Design

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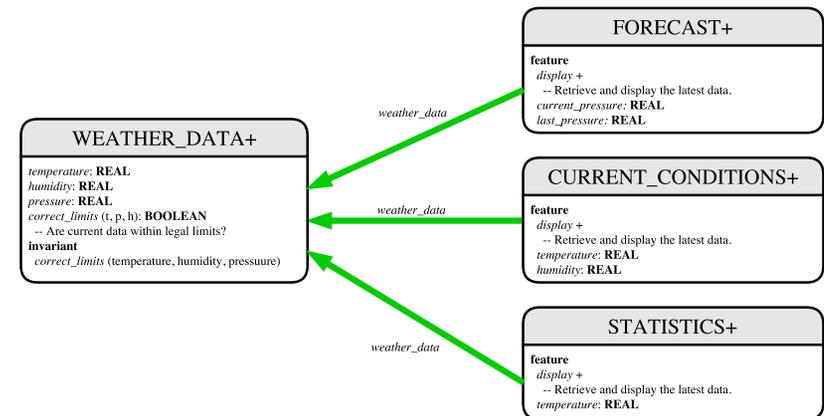
Motivating Problem



- A *weather station* maintains *weather data* such as *temperature*, *humidity*, and *pressure*.
- Various kinds of applications on these *weather data* should regularly update their *displays*:
 - *Forecast*: if expecting for rainy weather due to reduced *pressure*.
 - *Condition*: *temperature* in celsius and *humidity* in percentages.
 - *Statistics*: minimum/maximum/average measures of *temperature*.

12/12/20

First Design: Weather Station



Whenever the display feature is called, **retrieve** the current values of temperature, humidity, and/or pressure via the weather_data reference.

12/12/20

Implementing the First Design (1)



```
class WEATHER_DATA create make
feature -- Data
  temperature: REAL
  humidity: REAL
  pressure: REAL
feature -- Queries
  correct_limits(t,p,h: REAL): BOOLEAN
  ensure
    Result implies -36 <= t and t <= 60
    Result implies 50 <= p and p <= 110
    Result implies 0.8 <= h and h <= 100
feature -- Commands
  make (t, p, h: REAL)
  require
    correct_limits(temperature, pressure, humidity)
  ensure
    temperature = t and pressure = p and humidity = h
invariant
  correct_limits(temperature, pressure, humidity)
end
```

not

Implementing the First Design (2.2)



```
class CURRENT_CONDITIONS create make
feature -- Attributes
  temperature: REAL
  humidity: REAL
  weather_data: WEATHER_DATA
feature -- Commands
  make (wd: WEATHER_DATA)
  ensure weather_data = wd
  update
  do temperature := weather_data.temperature
     humidity := weather_data.humidity
  end
  display
  do update
     io.put_string("Current Conditions: ")
     io.put_real (temperature) ; io.put_string (" degrees C and ")
     io.put_real (humidity) ; io.put_string (" percent humidity%N")
  end
end
```

not

Implementing the First Design (2.1)



```
class FORECAST create make
feature -- Attributes
  current_pressure: REAL
  last_pressure: REAL
  weather_data: WEATHER_DATA
feature -- Commands
  make (wd: WEATHER_DATA)
  ensure weather_data = wd
  update
  do last_pressure := current_pressure
     current_pressure := weather_data.pressure
  end
  display
  do update
     if current_pressure > last_pressure then
       print("Improving weather on the way!%N")
     elseif current_pressure = last_pressure then
       print("More of the same%N")
     else print("Watch out for cooler, rainy weather%N") end
  end
end
```

not

Implementing the First Design (2.3)



```
class STATISTICS create make
feature -- Attributes
  weather_data: WEATHER_DATA
  current_temp: REAL
  max, min, sum_so_far: REAL
  num_readings: INTEGER
feature -- Commands
  make (wd: WEATHER_DATA)
  ensure weather_data = wd
  update
  do current_temp := weather_data.temperature
     -- Update min, max if necessary.
  end
  display
  do update
     print("Avg/Max/Min temperature = ")
     print(sum_so_far / num_readings + "/" + max + "/" min + "%N")
  end
end
```

not

Implementing the First Design (3)



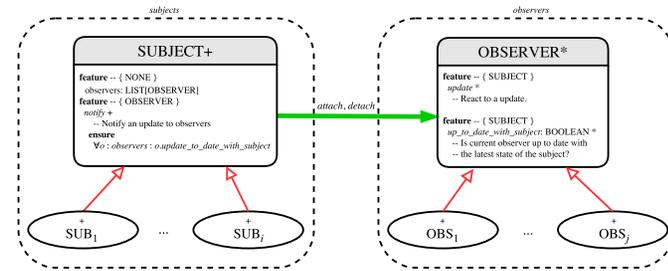
```

1 class WEATHER_STATION create make
2 feature -- Attributes
3   cc: CURRENT_CONDITIONS ; fd: FORECAST ; sd: STATISTICS
4   wd: WEATHER_DATA
5 feature -- Commands
6   make
7   do create wd.make (9, 75, 25)
8     create cc.make (wd) ; create fd.make (wd) ; create sd.make (wd)
9
10    wd.set_measurements (15, 60, 30.4)
11    cc.display ; fd.display ; sd.display
12    cc.display ; fd.display ; sd.display
13
14    wd.set_measurements (11, 90, 20)
15    cc.display ; fd.display ; sd.display
16  end
17 end
    
```

L14: Updates occur on cc, fd, sd even with the same data.

10/13/16

Observer Pattern: Architecture



- Observer (publish-subscribe) pattern: **one-to-many** relation.
 - Observers (*subscribers*) are attached to a subject (*publisher*).
 - The subject notify its attached observers about changes.
- Some interchangeable vocabulary:
 - subscribe \approx attach \approx register
 - unsubscribe \approx detach \approx unregister
 - publish \approx notify
 - handle \approx update

10/13/16

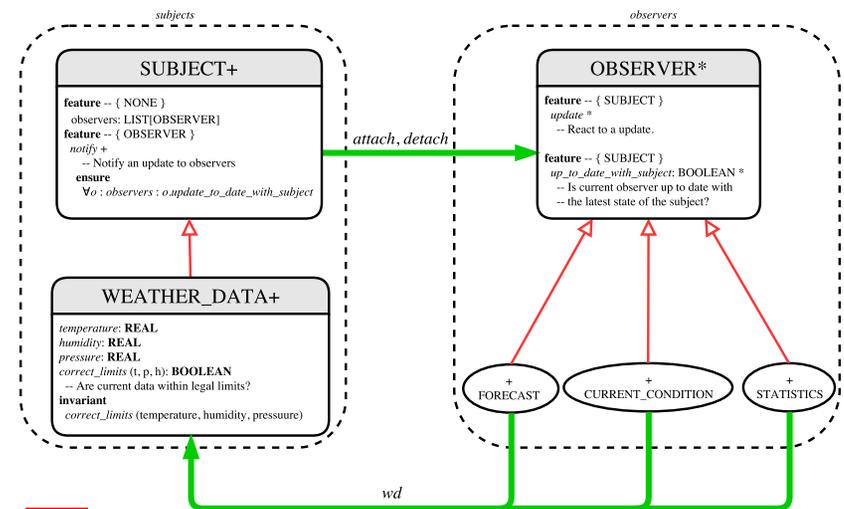
First Design: Good Design?



- Each application (CURRENT_CONDITION, FORECAST, STATISTICS) *cannot know* when the weather data change.
 - ⇒ All applications have to periodically initiate updates in order to keep the display results up to date.
 - ∴ Each inquiry of current weather data values is *a remote call*.
 - ∴ Waste of computing resources (e.g., network bandwidth) when there are actually no changes on the weather data.
- To avoid such overhead, it is better to let:
 - Each application is *subscribed/attached/registered* to the weather data.
 - The weather data *publish/notify* new changes.
 - ⇒ Updates on the application side occur only **when necessary**.

10/13/16

Observer Pattern: Weather Station



10/13/16

Implementing the Observer Pattern (1.1)



```
class SUBJECT create make
feature -- Attributes
  observers: LIST[OBSERVER]
feature -- Commands
  make
    do create {LINKED_LIST[OBSERVER]} observers.make
    ensure no_observers: observers.count = 0 end
feature -- Invoked by an OBSERVER
  attach (o: OBSERVER) -- Add 'o' to the observers
    require not_yet_attached: not observers.has (o)
    ensure is_attached: observers.has (o) end
  detach (o: OBSERVER) -- Add 'o' to the observers
    require currently_attached: observers.has (o)
    ensure is_attached: not observers.has (o) end
feature -- invoked by a SUBJECT
  notify -- Notify each attached observer about the update.
    do across observers as cursor loop cursor.item.update end
    ensure all_views_updated:
      across observers as o all o.item.up_to_date_with_subject end
    end
end
```

12/18/18

Implementing the Observer Pattern (2.1)



```
deferred class
  OBSERVER
feature -- To be effected by a descendant
  up_to_date_with_subject: BOOLEAN
  -- Is this observer up to date with its subject?
  deferred
  end

  update
  -- Update the observer's view of 's'
  deferred
  ensure
    up_to_date_with_subject: up_to_date_with_subject
  end
end
```

Each effective descendant class of OBSERVER should:

- Define what weather data are required to be up-to-date.
- Define how to update the required weather data.

12/18/18

Implementing the Observer Pattern (1.2)



```
class WEATHER_DATA
inherit SUBJECT rename make as make_subject end
create make
feature -- data available to observers
  temperature: REAL
  humidity: REAL
  pressure: REAL
  correct_limits(t,p,h: REAL): BOOLEAN
feature -- Initialization
  make (t, p, h: REAL)
  do
    make_subject -- initialize empty observers
    set_measurements (t, p, h)
  end
feature -- Called by weather station
  set_measurements(t, p, h: REAL)
  require correct_limits(t,p,h)
invariant
  correct_limits(temperature, pressure, humidity)
end
```

12/18/18

Implementing the Observer Pattern (2.2)



```
class FORECAST
inherit OBSERVER
feature -- Commands
  make(a_weather_data: WEATHER_DATA)
  do weather_data := a_weather_data
    weather_data.attach (Current)
  ensure weather_data = a_weather_data
    weather_data.observers.has (Current)
  end
feature -- Queries
  up_to_date_with_subject: BOOLEAN
  ensure then
    Result = current_pressure = weather_data.pressure
  update
  do -- Same as 1st design; Called only on demand
  end
  display
  do -- No need to update; Display contents same as in 1st design
  end
end
```

12/18/18

Implementing the Observer Pattern (2.3)



```
class CURRENT_CONDITIONS
inherit OBSERVER
feature -- Commands
  make(a_weather_data: WEATHER_DATA)
  do weather_data := a_weather_data
    weather_data.attach (Current)
  ensure weather_data = a_weather_data
    weather_data.observers.has (Current)
  end
feature -- Queries
  up_to_date_with_subject: BOOLEAN
  ensure then Result = temperature = weather_data.temperature and
    humidity = weather_data.humidity
  update
  do -- Same as 1st design; Called only on demand
  end
  display
  do -- No need to update; Display contents same as in 1st design
  end
end
```

L12/13

Implementing the Observer Pattern (3)



```
1 class WEATHER_STATION create make
2 feature -- Attributes
3   cc: CURRENT_CONDITIONS ; fd: FORECAST ; sd: STATISTICS
4   wd: WEATHER_DATA
5 feature -- Commands
6   make
7   do create wd.make (9, 75, 25)
8     create cc.make (wd) ; create fd.make (wd) ; create sd.make (wd)
9
10    wd.set_measurements (15, 60, 30.4)
11    wd.notify
12    cc.display ; fd.display ; sd.display
13    cc.display ; fd.display ; sd.display
14
15    wd.set_measurements (11, 90, 20)
16    wd.notify
17    cc.display ; fd.display ; sd.display
18  end
19 end
```

L13: cc, fd, sd make use of “cached” data values.

L12/13

Implementing the Observer Pattern (2.4)



```
class STATISTICS
inherit OBSERVER
feature -- Commands
  make(a_weather_data: WEATHER_DATA)
  do weather_data := a_weather_data
    weather_data.attach (Current)
  ensure weather_data = a_weather_data
    weather_data.observers.has (Current)
  end
feature -- Queries
  up_to_date_with_subject: BOOLEAN
  ensure then
  Result = current_temperature = weather_data.temperature
  update
  do -- Same as 1st design; Called only on demand
  end
  display
  do -- No need to update; Display contents same as in 1st design
  end
end
```

L12/13

Observer Pattern: Limitation? (1)

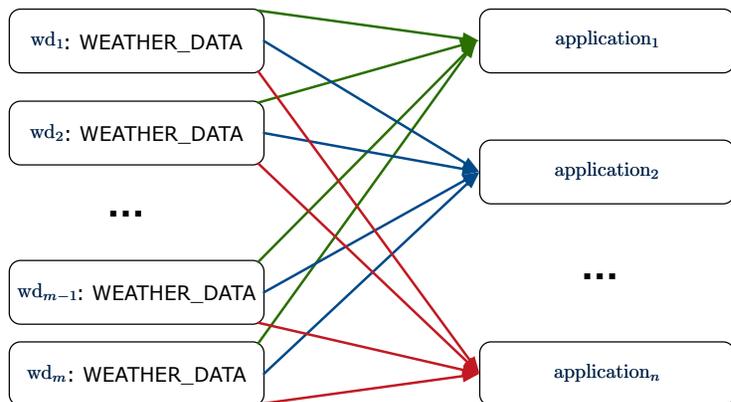


- The *observer design pattern* is a reasonable solution to building a *one-to-many* relationship: one subject (publisher) and multiple observers (subscribers).
- But what if a *many-to-many* relationship is required for the application under development?
 - *Multiple weather data* are maintained by weather stations.
 - Each application observes *all* these *weather data*.
 - But, each application still stores the *latest* measure only. e.g., the statistics app stores one copy of temperature
 - Whenever some weather station updates the temperature of its associated *weather data*, all *relevant* subscribed applications (i.e., current conditions, statistics) should update their temperatures.
- How can the observer pattern solve this general problem?
 - Each *weather data* maintains a list of subscribed *applications*.
 - Each *application* is subscribed to *multiple weather data*.

L12/13

Observer Pattern: Limitation? (2)

What happens at runtime when building a **many-to-many** relationship using the **observer pattern**?



Graph complexity, with m subjects and n observers? $[O(m \cdot n)]$

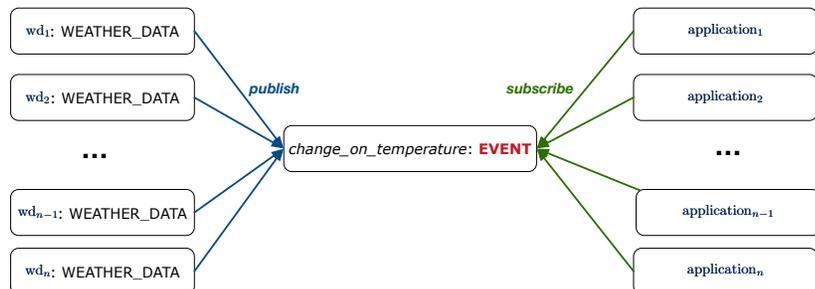
Event-Driven Design (2)

In an **event-driven design**:

- Each variable being observed (e.g., temperature, humidity, pressure) is called a **monitored variable**.
e.g., A **nuclear power plant** (i.e., the **subject**) has its temperature and pressure being **monitored** by a shutdown system (i.e., an **observer**): as soon as values of these **monitored variables** exceed the normal threshold, the SDS will be notified and react by shutting down the plant.
- Each **monitored variable** is declared as an **event**:
 - An **observer** is **attached/subscribed** to the relevant events.
 - CURRENT_CONDITION attached to events for temperature, humidity.
 - FORECAST only subscribed to the event for pressure.
 - STATISTICS only subscribed to the event for temperature.
 - A **subject notifies/publishes** changes to the relevant events.

Event-Driven Design (1)

Here is what happens at runtime when building a **many-to-many** relationship using the **event-driven design**.



Graph complexity, with m subjects and n observers? $[O(m + n)]$

Additional cost by adding a new subject? $[O(1)]$

Additional cost by adding a new observer? $[O(1)]$

Additional cost by adding a new event type? $[O(m + n)]$

Event-Driven Design: Example

Darlington Nuclear Generating System (by Ontario Power Generation)



Image Source: <https://www.opg.com>

Event-Driven Design: Implementation



- Requirements for implementing an **event-driven design** are:
 - When an **observer** object is **subscribed to** an **event**, it attaches:
 - The **reference/pointer** to an update operation
Such reference/pointer is used for **delayed** executions.
 - Itself (i.e., the **context object** for invoking the update operation)
 - For the **subject** object to **publish** an update to the **event**, it:
 - Iterates over its (previously) attached operation **references/pointers**
 - Invokes these operations, which update the corresponding observers
- Both requirements can be satisfied by Eiffel and Java.
- We will compare how an **event-driven design** for the weather station problems is implemented in Eiffel and Java.
⇒ It is more convenient to implement such design in Eiffel.

25/06/18

Event-Driven Design in Java (2)



```
1 public class WeatherData {
2     private double temperature;
3     private double pressure;
4     private double humidity;
5     public WeatherData(double t, double p, double h) {
6         setMeasurements(t, h, p);
7     }
8     public static Event changeOnTemperature = new Event();
9     public static Event changeOnHumidity = new Event();
10    public static Event changeOnPressure = new Event();
11    public void setMeasurements(double t, double h, double p) {
12        temperature = t;
13        humidity = h;
14        pressure = p;
15        changeOnTemperature.publish(temperature);
16        changeOnHumidity.publish(humidity);
17        changeOnPressure.publish(pressure);
18    }
19 }
```

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Event-Driven Design in Java (1)



```
1 public class Event {
2     Hashtable<Object, MethodHandle> listenersActions;
3     Event() { listenersActions = new Hashtable<>(); }
4     void subscribe(Object listener, MethodHandle action) {
5         listenersActions.put(listener, action);
6     }
7     void publish(Object arg) {
8         for (Object listener : listenersActions.keySet()) {
9             MethodHandle action = listenersActions.get(listener);
10            try {
11                action.invokeWithArguments(listener, arg);
12            } catch (Throwable e) { }
13        }
14    }
15 }
```

- L5:** Both the delayed action reference and its context object (or call target) listener are stored into the table.
- L11:** An invocation is made from retrieved listener and action.

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Event-Driven Design in Java (3)



```
1 public class CurrentConditions {
2     private double temperature; private double humidity;
3     public void updateTemperature(double t) { temperature = t; }
4     public void updateHumidity(double h) { humidity = h; }
5     public CurrentConditions() {
6         MethodHandles.Lookup lookup = MethodHandles.lookup();
7         try {
8             MethodHandle ut = lookup.findVirtual(
9                 this.getClass(), "updateTemperature",
10                MethodType.methodType(void.class, double.class));
11            WeatherData.changeOnTemperature.subscribe(this, ut);
12            MethodHandle uh = lookup.findVirtual(
13                this.getClass(), "updateHumidity",
14                MethodType.methodType(void.class, double.class));
15            WeatherData.changeOnHumidity.subscribe(this, uh);
16        } catch (Exception e) { e.printStackTrace(); }
17    }
18    public void display() {
19        System.out.println("Temperature: " + temperature);
20        System.out.println("Humidity: " + humidity); } }
```

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Event-Driven Design in Java (4)



```
1 public class WeatherStation {
2     public static void main(String[] args) {
3         WeatherData wd = new WeatherData(9, 75, 25);
4         CurrentConditions cc = new CurrentConditions();
5         System.out.println("=====");
6         wd.setMeasurements(15, 60, 30.4);
7         cc.display();
8         System.out.println("=====");
9         wd.setMeasurements(11, 90, 20);
10        cc.display();
11    } }

```

L4 invokes

```
WeatherData.changeOnTemperature.subscribe(
    cc, "updateTemperature handle")
```

L6 invokes

```
WeatherData.changeOnTemperature.publish(15)
```

which in turn invokes

```
"updateTemperature handle".invokeWithArguments(cc, 15)
```

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Event-Driven Design in Eiffel (2)



```
1 class WEATHER_DATA
2 create make
3 feature -- Measurements
4     temperature: REAL ; humidity: REAL ; pressure: REAL
5     correct_limits(t,p,h: REAL): BOOLEAN do ... end
6     make (t, p, h: REAL) do ... end
7 feature -- Event for data changes
8     change_on_temperature: EVENT[TUPLE[REAL]]once create Result end
9     change_on_humidity: EVENT[TUPLE[REAL]]once create Result end
10    change_on_pressure: EVENT[TUPLE[REAL]]once create Result end
11 feature -- Command
12    set_measurements(t, p, h: REAL)
13    require correct_limits(t,p,h)
14    do temperature := t ; pressure := p ; humidity := h
15        change_on_temperature.publish ([t])
16        change_on_humidity.publish ([p])
17        change_on_pressure.publish ([h])
18    end
19 invariant correct_limits(temperature, pressure, humidity) end

```

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Event-Driven Design in Eiffel (1)



```
1 class EVENT [ARGUMENTS -> TUPLE ]
2 create make
3 feature -- Initialization
4     actions: LINKED_LIST[PROCEDURE[ARGUMENTS]]
5     make do create actions.make end
6 feature
7     subscribe (an_action: PROCEDURE[ARGUMENTS])
8         require action_not_already_subscribed: not actions.has(an_action)
9         do actions.extend (an_action)
10        ensure action_subscribed: action.has(an_action) end
11    publish (args: ARGUMENTS)
12    do from actions.start until actions.after
13        loop actions.item.call (args) ; actions.forth end
14    end
15 end

```

- L1 constrains the generic parameter ARGUMENTS: any class that instantiates ARGUMENTS must be a **descendant** of TUPLE.
- L4: The type **PROCEDURE** encapsulates both the context object and the reference/pointer to some update operation.

28/06/16

Event-Driven Design in Eiffel (3)



```
1 class CURRENT_CONDITIONS
2 create make
3 feature -- Initialization
4     make(wd: WEATHER_DATA)
5     do
6         wd.change_on_temperature.subscribe (agent update_temperature)
7         wd.change_on_humidity.subscribe (agent update_humidity)
8     end
9 feature
10    temperature: REAL
11    humidity: REAL
12    update_temperature (t: REAL) do temperature := t end
13    update_humidity (h: REAL) do humidity := h end
14    display do ... end
15 end

```

- **agent cmd** retrieves the pointer to cmd and its context object.
- L6 ≈ ... (agent **Current**.update_temperature)
- Contrast L6 with L8–11 in Java class CurrentConditions.

28/06/16



Event-Driven Design in Eiffel (4)

```

1 class WEATHER_STATION create make
2 feature
3   cc: CURRENT_CONDITIONS
4   make
5   do create wd.make (9, 75, 25)
6     create cc.make (wd)
7     wd.set_measurements (15, 60, 30.4)
8     cc.display
9     wd.set_measurements (11, 90, 20)
10    cc.display
11  end
12 end

```

L6 invokes

```
wd.change_on_temperature.subscribe(
  agent cc.update_temperature)
```

L7 invokes

```
wd.change_on_temperature.publish([15])
```

which in turn invokes `cc.update_temperature(15)`

START



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Event-Driven Design: Eiffel vs. Java

Storing observers/listeners of an event

- Java, in the Event class:

```
Hashtable<Object, MethodHandle> listenersActions;
```

- Eiffel, in the EVENT class:

```
actions: LINKED_LIST[PROCEDURE[ARGUMENTS]]
```

Creating and passing function pointers

- Java, in the CurrentConditions class constructor:

```
MethodHandle ut = lookup.findVirtual(
  this.getClass(), "updateTemperature",
  MethodType.methodType(void.class, double.class));
WeatherData.changeOnTemperature.subscribe(this, ut);
```

- Eiffel, in the CURRENT_CONDITIONS class construction:

```
wd.change_on_temperature.subscribe (agent update_temperature)
```

⇒ Eiffel's type system has been better thought-out for **design**.

START



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Event-Driven Design: Eiffel vs. Java

Not 4:

Program Correctness

OOSC2 Chapter 11



EECS3311 A: Software Design
Winter 2020

CHEN-WEI WANG



Assertions: Weak vs. Strong

- Describe each assertion as **a set of satisfying value**.
 $x > 3$ has satisfying values $\{ x \mid x > 3 \} = \{ 4, 5, 6, 7, \dots \}$
 $x > 4$ has satisfying values $\{ x \mid x > 4 \} = \{ 5, 6, 7, \dots \}$
- An assertion p is **stronger** than an assertion q if p 's set of satisfying values is a subset of q 's set of satisfying values.
 - Logically speaking, p being stronger than q (or, q being weaker than p) means $p \Rightarrow q$.
 - e.g., $x > 4 \Rightarrow x > 3$
- What's the weakest assertion? [TRUE]
- What's the strongest assertion? [FALSE]
- In **Design by Contract** :
 - A weaker **invariant** has more acceptable object states
e.g., $balance > 0$ vs. $balance > 100$ as an invariant for ACCOUNT
 - A weaker **precondition** has more acceptable input values
 - A weaker **postcondition** has more acceptable output values

Not 4:



Assertions: Preconditions

Given **preconditions** P_1 and P_2 , we say that

P_2 **requires less** than P_1 if

P_2 is **less strict** on (thus **allowing more**) inputs than P_1 does.

$$\{ x \mid P_1(x) \} \supseteq \{ x \mid P_2(x) \}$$

More concisely:

$$P_1 \Rightarrow P_2$$

e.g., For command `withdraw(amount: INTEGER)`,

$P_2: amount \geq 0$ **requires less** than $P_1: amount > 0$

What is the **precondition** that **requires the least**? [true]

Not 4:



Assertions: Postconditions

Given **postconditions** or **invariants** Q_1 and Q_2 , we say that

Q_2 **ensures more** than Q_1 if Q_2 is **stricter** on (thus **allowing less**) outputs than Q_1 does.

$$\{x \mid Q_2(x)\} \subseteq \{x \mid Q_1(x)\}$$

More concisely:

$$Q_2 \Rightarrow Q_1$$

e.g., For query $q(i: \text{INTEGER}): \text{BOOLEAN}$,

$Q_2: \text{Result} = (i > 0) \wedge (i \bmod 2 = 0)$ **ensures more** than

$Q_1: \text{Result} = (i > 0) \vee (i \bmod 2 = 0)$

What is the **postcondition** that **ensures the most**? [**false**]

not 4:



Motivating Examples (2)

Is this feature correct?

```
class FOO
  i: INTEGER
  increment_by_9
  require
    i > 5
  do
    i := i + 9
  ensure
    i > 13
end
```

Q: Is $i > 5$ too weak or too strong?

A: Maybe too strong

\therefore assertion $i > 5$ disallows 5 which would not fail postcondition.
Whether 5 should be allowed depends on the requirements.

not 4:



Motivating Examples (1)

Is this feature correct?

```
class FOO
  i: INTEGER
  increment_by_9
  require
    i > 3
  do
    i := i + 9
  ensure
    i > 13
end
```

Q: Is $i > 3$ is too weak or too strong?

A: Too weak

\therefore assertion $i > 3$ allows value 4 which would fail postcondition.

not 4:



Software Correctness

- Correctness is a **relative** notion: **consistency** of **implementation** with respect to **specification**.
 \Rightarrow This assumes there is a specification!
- We introduce a formal and systematic way for formalizing a program **S** and its **specification** (pre-condition **Q** and post-condition **R**) as a **Boolean predicate**: $\{Q\} S \{R\}$
 - e.g., $\{i > 3\} i := i + 9 \{i > 13\}$
 - e.g., $\{i > 5\} i := i + 9 \{i > 13\}$
 - If $\{Q\} S \{R\}$ **can** be proved **TRUE**, then the **S** is **correct**.
e.g., $\{i > 5\} i := i + 9 \{i > 13\}$ **can** be proved TRUE.
 - If $\{Q\} S \{R\}$ **cannot** be proved **TRUE**, then the **S** is **incorrect**.
e.g., $\{i > 3\} i := i + 9 \{i > 13\}$ **cannot** be proved TRUE.

not 4:



Hoare Logic

- Consider a program **S** with precondition **Q** and postcondition **R**.
 - $\{Q\} S \{R\}$ is a **correctness predicate** for program **S**
 - $\{Q\} S \{R\}$ is TRUE if program **S** starts executing in a state satisfying the precondition **Q**, and then:
 - The program **S** terminates.
 - Given that program **S** terminates, then it terminates in a state satisfying the postcondition **R**.
- Separation of concerns
 - requires a proof of **termination**.
 - requires a proof of **partial correctness**.
 Proofs of (a) + (b) imply **total correctness**.

not 48



Proof of Hoare Triple using wp

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

- $wp(S, R)$ is the **weakest precondition for S to establish R**.
- S can be:
 - Assignments ($x := y$)
 - Alternations (**if ... then ... else ... end**)
 - Sequential compositions ($S_1 ; S_2$)
 - Loops (**from ... until ... loop ... end**)
- We will learn how to calculate the **wp** for the above programming constructs.

not 48



Hoare Logic and Software Correctness

Consider the **contract view** of a feature f (whose body of implementation is **S**) as a **Hoare Triple**:

$$\{Q\} S \{R\}$$

Q is the **precondition** of f .
S is the implementation of f .
R is the **postcondition** of f .

- $\{true\} S \{R\}$
All input values are valid [Most-user friendly]
- $\{false\} S \{R\}$
All input values are invalid [Most useless for clients]
- $\{Q\} S \{true\}$
All output values are valid [Most risky for clients; Easiest for suppliers]
- $\{Q\} S \{false\}$
All output values are invalid [Most challenging coding task]
- $\{true\} S \{true\}$
All inputs/outputs are valid (No contracts) [Least informative]

not 48



Hoare Logic A Simple Example

Given $\{??\} n := n + 9 \{n > 13\}$:

- $n > 4$ is the **weakest precondition (wp)** for the given implementation ($n := n + 9$) to start and establish the postcondition ($n > 13$).
- Any precondition that is **equal to or stronger than** the wp ($n > 4$) will result in a correct program.
e.g., $\{n > 5\} n := n + 9 \{n > 13\}$ can be proved **TRUE**.
- Any precondition that is **weaker than** the wp ($n > 4$) will result in an incorrect program.
e.g., $\{n > 3\} n := n + 9 \{n > 13\}$ cannot be proved **TRUE**.
Counterexample: $n = 4$ satisfies precondition $n > 3$ but the output $n = 13$ fails postcondition $n > 13$.

not 48

Denoting New and Old Values



In the **postcondition**, for a program variable x :

- We write $\boxed{x_0}$ to denote its **pre-state (old)** value.
 - We write \boxed{x} to denote its **post-state (new)** value.
- Implicitly, in the **precondition**, all program variables have their **pre-state** values.

e.g., $\{b_0 > a\} b := b - a \{b = b_0 - a\}$

- Notice that:
 - We may choose to write “ b ” rather than “ b_0 ” in preconditions
∴ All variables are pre-state values in preconditions
 - We don't write “ b_0 ” in program
∴ there might be **multiple intermediate values** of a variable due to sequential composition

13/07/18

wp Rule: Assignments (1)



$$wp(x := e, R) = R[x := e]$$

$R[x := e]$ means to substitute all **free occurrences** of variable x in postcondition R by expression e .

13/07/18

wp Rule: Assignments (2)



Recall:

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

How do we prove $\{Q\} x := e \{R\}$?

$$\{Q\} x := e \{R\} \iff Q \Rightarrow \underbrace{R[x := e]}_{wp(x := e, R)}$$

13/07/18

wp Rule: Assignments (3) Exercise



What is the weakest precondition for a program $x := x + 1$ to establish the postcondition $x > x_0$?

$$\{??\} x := x + 1 \{x > x_0\}$$

For the above Hoare triple to be **TRUE**, it must be that $?? \Rightarrow wp(x := x + 1, x > x_0)$.

$$\begin{aligned} & wp(x := x + 1, x > x_0) \\ &= \{Rule\ of\ wp:\ Assignment\} \\ & \quad x > x_0[x := x_0 + 1] \\ &= \{Replacing\ x\ by\ x_0 + 1\} \\ & \quad x_0 + 1 > x_0 \\ &= \{1 > 0\ always\ true\} \\ & \quad True \end{aligned}$$

Any precondition is OK.

False is valid but not useful.

13/07/18



wp Rule: Assignments (4) Exercise

What is the weakest precondition for a program $x := x + 1$ to establish the postcondition $x = 23$?

$$\{??\} x := x + 1 \{x = 23\}$$

For the above Hoare triple to be **TRUE**, it must be that $?? \Rightarrow wp(x := x + 1, x = 23)$.

$$\begin{aligned}
& wp(x := x + 1, x = 23) \\
&= \{Rule\ of\ wp:\ Assignments\} \\
&\quad x = 23[x := x_0 + 1] \\
&= \{Replacing\ x\ by\ x_0 + 1\} \\
&\quad x_0 + 1 = 23 \\
&= \{arithmetic\} \\
&\quad x_0 = 22
\end{aligned}$$

Any precondition weaker than $x = 22$ is not OK.

FALSE



wp Rule: Alternations (2)

Recall: $\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$

How do we prove that $\{Q\} \text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end } \{R\}$?

```

{Q}
if B then
  {Q ∧ B} S1 {R}
else
  {Q ∧ ¬B} S2 {R}
end
{R}

```

$$\begin{aligned}
& \{Q\} \text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end } \{R\} \\
& \iff \left(\begin{array}{l} \{Q \wedge B\} S_1 \{R\} \\ \wedge \\ \{Q \wedge \neg B\} S_2 \{R\} \end{array} \right) \iff \left(\begin{array}{l} (Q \wedge B) \Rightarrow wp(S_1, R) \\ \wedge \\ (Q \wedge \neg B) \Rightarrow wp(S_2, R) \end{array} \right)
\end{aligned}$$

FALSE



wp Rule: Alternations (1)

$$wp(\text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end}, R) = \left(\begin{array}{l} B \Rightarrow wp(S_1, R) \\ \wedge \\ \neg B \Rightarrow wp(S_2, R) \end{array} \right)$$

The wp of an alternation is such that **all branches** are able to establish the postcondition R .

FALSE



wp Rule: Alternations (3) Exercise

Is this program correct?

```

{x > 0 ∧ y > 0}
if x > y then
  bigger := x ; smaller := y
else
  bigger := y ; smaller := x
end
{bigger ≥ smaller}

```

$$\begin{aligned}
& \left(\begin{array}{l} \{(x > 0 \wedge y > 0) \wedge (x > y)\} \\ \quad bigger := x ; smaller := y \\ \{bigger \geq smaller\} \end{array} \right) \\
& \wedge \\
& \left(\begin{array}{l} \{(x > 0 \wedge y > 0) \wedge \neg(x > y)\} \\ \quad bigger := y ; smaller := x \\ \{bigger \geq smaller\} \end{array} \right)
\end{aligned}$$

FALSE

wp Rule: Sequential Composition (1)



$$wp(S_1 ; S_2, R) = wp(S_1, wp(S_2, R))$$

The *wp* of a sequential composition is such that the **first phase** establishes the *wp* for the **second phase** to establish the postcondition *R*.

20/07/18

wp Rule: Sequential Composition (2)



Recall:

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

How do we prove $\{Q\} S_1 ; S_2 \{R\}$?

$$\{Q\} S_1 ; S_2 \{R\} \iff Q \Rightarrow \underbrace{wp(S_1, wp(S_2, R))}_{wp(S_1 ; S_2, R)}$$

20/07/18

wp Rule: Sequential Composition (3) Exercise



Is $\{True\} tmp := x; x := y; y := tmp \{x > y\}$ correct?
 If and only if $True \Rightarrow wp(tmp := x; x := y; y := tmp, x > y)$

$$\begin{aligned} & wp(tmp := x; x := y; y := tmp, x > y) \\ = & \{wp \text{ rule for seq. comp.}\} \\ & wp(tmp := x, wp(x := y; y := tmp, x > y)) \\ = & \{wp \text{ rule for seq. comp.}\} \\ & wp(tmp := x, wp(x := y, wp(y := tmp, x > y))) \\ = & \{wp \text{ rule for assignment}\} \\ & wp(tmp := x, wp(x := y, x > tmp)) \\ = & \{wp \text{ rule for assignment}\} \\ & wp(tmp := x, y > tmp) \\ = & \{wp \text{ rule for assignment}\} \\ & y > x \end{aligned}$$

$\therefore True \Rightarrow y > x$ does not hold in general.

\therefore The above program is not correct.

20/07/18

Loops



- A loop is a way to compute a certain result by *successive approximations*.

e.g. computing the maximum value of an array of integers

- Loops are needed and powerful
- But loops **very hard** to get right:

- Infinite loops [termination]
- "off-by-one" error [partial correctness]
- Improper handling of borderline cases [partial correctness]
- Not establishing the desired condition [partial correctness]

20/07/18

Loops: Binary Search

| BS1 | BS2 |
|---|---|
| <pre> from i := 1; j := n until i = j loop m := (i + j) // 2 if t @ m <= x then i := m else j := m end end Result := (x = t @ i) </pre> | <pre> from i := 1; j := n; found := false until i = j and not found loop m := (i + j) // 2 if t @ m < x then i := m + 1 elseif t @ m = x then found := true else j := m - 1 end end Result := found </pre> |
| BS3 | BS4 |
| <pre> from i := 0; j := n until i = j loop m := (i + j) // 2 if t @ m <= x then i := m + 1 else j := m end end if i >= j and i <= n then Result := (x = t @ i) else Result := false end </pre> | <pre> from i := 0; j := n + 1 until i = j loop m := (i + j) // 2 if t @ m <= x then i := m + 1 else j := m end end if i >= j and i <= n then Result := (x = t @ i) else Result := false end </pre> |

4 implementations for binary search: published, but *wrong!*

See page 381 in *Object Oriented Software Construction*

25.07.18

Contracts for Loops: Syntax

```

from
  Sinit
invariant
  invariant_tag: I -- Boolean expression for partial correctness
until
  B
loop
  Sbody
variant
  variant_tag: V -- Integer expression for termination
end

```

25.07.18

Correctness of Loops

How do we prove that the following loops are correct?

```

{Q}
from
  Sinit
until
  B
loop
  Sbody
end
{R}

```

```

{Q}
Sinit
while (¬ B) {
  Sbody
}
{R}

```

- In case of C/Java, $\neg B$ denotes the **stay condition**.
- In case of Eiffel, B denotes the **exit condition**.
There is native, syntactic support for checking/proving the **total correctness** of loops.

25.07.18

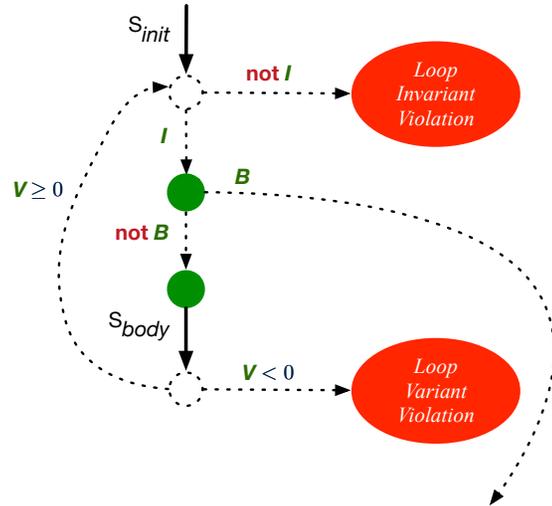
Contracts for Loops

- Use of **loop invariants (LI)** and **loop variants (LV)**.
 - Invariants:** Boolean expressions for **partial correctness**.
 - Typically a special case of the postcondition.
e.g., Given postcondition “**Result is maximum of the array**”:
LI can be “**Result is maximum of the part of array scanned so far**”.
 - Established before the very first iteration.
 - Maintained TRUE after each iteration.
 - Variants:** Integer expressions for **termination**
 - Denotes the **number of iterations remaining**
 - Decreased** at the end of each subsequent iteration
 - Maintained **non-negative** at the end of each iteration.
 - As soon as value of LV reaches **zero**, meaning that no more iterations remaining, the loop must exit.
- Remember:

total correctness = **partial correctness** + **termination**

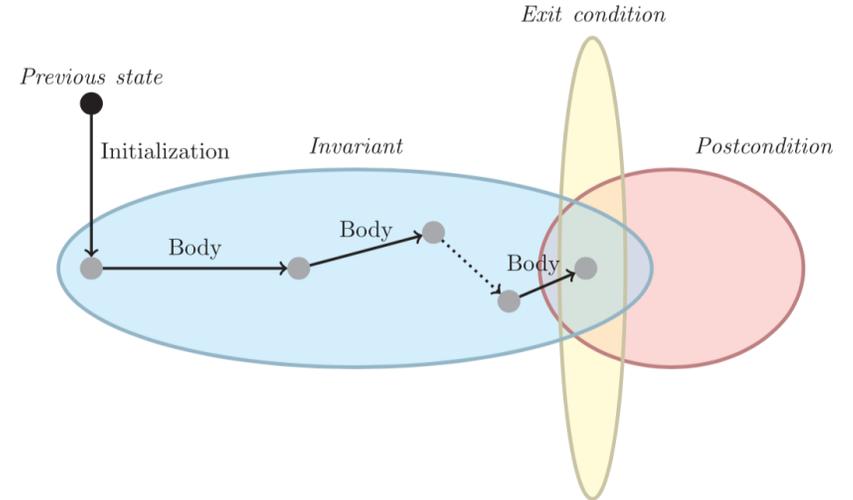
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Contracts for Loops: Runtime Checks (1)



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Contracts for Loops: Visualization



Digram Source: page 5 in *Loop Invariants: Analysis, Classification, and Examples*

28/07/18

Contracts for Loops: Runtime Checks (2)



```

1 test
2   local
3     i: INTEGER
4   do
5     from
6       i := 1
7     invariant
8       1 <= i and i <= 6
9     until
10      i > 5
11    loop
12      io.put_string ("iteration " + i.out + "%N")
13      i := i + 1
14    variant
15      6 - i
16  end
17 end
    
```

L8: Change to $1 \leq i$ and $i \leq 5$ for a **Loop Invariant Violation**.

L10: Change to $i > 0$ to bypass the body of loop.

L15: Change to $5 - i$ for a **Loop Variant Violation**.

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Contracts for Loops: Example 1.1



```

find_max (a: ARRAY [INTEGER]): INTEGER
  local i: INTEGER
  do
    from
      i := a.lower ; Result := a[i]
    invariant
      loop_invariant: --  $\forall j | a.lower \leq j \leq i \bullet Result \geq a[j]$ 
      across a.lower |..| i as j all Result >= a [j.item] end
    until
      i > a.upper
    loop
      if a [i] > Result then Result := a [i] end
      i := i + 1
    variant
      loop_variant: a.upper - i + 1
    end
  ensure
    correct_result: --  $\forall j | a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
    across a.lower |..| a.upper as j all Result >= a [j.item]
  end
end
    
```

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Contracts for Loops: Example 1.2

Consider the feature call `find_max(⟨⟨20, 10, 40, 30⟩⟩)`, given:

- **Loop Invariant:** $\forall j \mid a.lower \leq j \leq i \bullet Result \geq a[j]$
- **Loop Variant:** $a.upper - i + 1$

| AFTER ITERATION | i | Result | LI | EXIT ($i > a.upper$)? | LV |
|-----------------|----------|--------|----|-------------------------|----|
| Initialization | 1 | 20 | ✓ | × | – |
| 1st | 2 | 20 | ✓ | × | 3 |
| 2nd | 3 | 20 | × | – | – |

Loop invariant violation at the end of the 2nd iteration:

$$\forall j \mid a.lower \leq j \leq 3 \bullet 20 \geq a[j]$$

evaluates to **false** $\because 20 \not\geq a[3] = 40$

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Contracts for Loops: Example 2.2

Consider the feature call `find_max(⟨⟨20, 10, 40, 30⟩⟩)`, given:

- **Loop Invariant:** $\forall j \mid a.lower \leq j < i \bullet Result \geq a[j]$
- **Loop Variant:** $a.upper - i$

| AFTER ITERATION | i | Result | LI | EXIT ($i > a.upper$)? | LV |
|-----------------|---|--------|----|-------------------------|-----------|
| Initialization | 1 | 20 | ✓ | × | – |
| 1st | 2 | 20 | ✓ | × | 2 |
| 2nd | 3 | 20 | ✓ | × | 1 |
| 3rd | 4 | 40 | ✓ | × | 0 |
| 4th | 5 | 40 | ✓ | ✓ | -1 |

Loop variant violation at the end of the 2nd iteration

$\because a.upper - i = 4 - 5$ evaluates to **non-zero**.

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Contracts for Loops: Example 2.1

```

find_max (a: ARRAY [INTEGER]): INTEGER
local i: INTEGER
do
  from
    i := a.lower ; Result := a[i]
  invariant
    loop_invariant: --  $\forall j \mid a.lower \leq j < i \bullet Result \geq a[j]$ 
    across a.lower |..| (i - 1) as j all Result >= a [j.item] end
  until
    i > a.upper
  loop
    if a [i] > Result then Result := a [i] end
    i := i + 1
  variant
    loop_variant: a.upper - i
  end
ensure
  correct_result: --  $\forall j \mid a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
  across a.lower |..| a.upper as j all Result >= a [j.item]
end
end

```

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Contracts for Loops: Example 3.1

```

find_max (a: ARRAY [INTEGER]): INTEGER
local i: INTEGER
do
  from
    i := a.lower ; Result := a[i]
  invariant
    loop_invariant: --  $\forall j \mid a.lower \leq j < i \bullet Result \geq a[j]$ 
    across a.lower |..| (i - 1) as j all Result >= a [j.item] end
  until
    i > a.upper
  loop
    if a [i] > Result then Result := a [i] end
    i := i + 1
  variant
    loop_variant: a.upper - i + 1
  end
ensure
  correct_result: --  $\forall j \mid a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
  across a.lower |..| a.upper as j all Result >= a [j.item]
end
end

```

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Contracts for Loops: Example 3.2

Consider the feature call `find_max(⟨⟨20, 10, 40, 30⟩⟩)`, given:

- **Loop Invariant**: $\forall j \mid a.lower \leq j < i \bullet Result \geq a[j]$
- **Loop Variant**: $a.upper - i + 1$
- **Postcondition**: $\forall j \mid a.lower \leq j \leq a.upper \bullet Result \geq a[j]$

| AFTER ITERATION | i | Result | LI | EXIT ($i > a.upper$)? | LV |
|-----------------|---|--------|----|-------------------------|----|
| Initialization | 1 | 20 | ✓ | × | — |
| 1st | 2 | 20 | ✓ | × | 3 |
| 2nd | 3 | 20 | ✓ | × | 2 |
| 3rd | 4 | 40 | ✓ | × | 1 |
| 4th | 5 | 40 | ✓ | ✓ | 0 |

RECAP

Proving Correctness of Loops (1)

```
{Q} from
  S_init
invariant
  I
until
  B
loop
  S_body
variant
  V
end {R}
```

- A loop is **partially correct** if:
 - Given precondition Q , the initialization step S_{init} establishes LI .
 - At the end of S_{body} , if not yet to exit, LI is maintained.
 - If ready to exit and LI maintained, postcondition R is established.
- A loop **terminates** if:
 - Given LI , and not yet to exit, S_{body} maintains LV as non-negative.
 - Given LI , and not yet to exit, S_{body} decrements LV .

RECAP

Contracts for Loops: Exercise

```
class DICTIONARY[V, K]
feature {NONE} -- Implementations
  values: ARRAY[K]
  keys: ARRAY[K]
feature -- Abstraction Function
  model: FUN[K, V]
feature -- Queries
  get_keys(v: V): ITERABLE[K]
  local i: INTEGER; ks: LINKED_LIST[K]
  do
    from i := keys.lower ; create ks.make_empty
    invariant ??
    until i > keys.upper
    do if values[i] ~ v then ks.extend(keys[i]) end
    end
  Result := ks.new_cursor
ensure
  result_valid:  $\forall k \in \text{Result} \bullet \text{model.item}(k) \sim v$ 
  no_missing_keys:  $\forall k \in \text{model.domain} \bullet \text{model.item}(k) \sim v \Rightarrow k \in \text{Result}$ 
end
```

RECAP

Proving Correctness of Loops (2)

```
{Q} from S_init invariant I until B loop S_body variant V end {R}
```

- A loop is **partially correct** if:
 - Given precondition Q , the initialization step S_{init} establishes LI .
 $\{Q\} S_{init} \{I\}$
 - At the end of S_{body} , if not yet to exit, LI is maintained.
 $\{I \wedge \neg B\} S_{body} \{I\}$
 - If ready to exit and LI maintained, postcondition R is established.
 $I \wedge B \Rightarrow R$
- A loop **terminates** if:
 - Given LI , and not yet to exit, S_{body} maintains LV as non-negative.
 $\{I \wedge \neg B\} S_{body} \{V \geq 0\}$
 - Given LI , and not yet to exit, S_{body} decrements LV .
 $\{I \wedge \neg B\} S_{body} \{V < V_0\}$

RECAP

Proving Correctness of Loops: Exercise (1.1)



Prove that the following program is correct:

```

find_max (a: ARRAY [INTEGER]): INTEGER
local i: INTEGER
do
  from
    i := a.lower ; Result := a[i]
  invariant
    loop_invariant:  $\forall j | a.lower \leq j < i \bullet Result \geq a[j]$ 
  until
    i > a.upper
  loop
    if a [i] > Result then Result := a [i] end
    i := i + 1
  variant
    loop_variant: a.upper - i + 1
end
ensure
  correct_result:  $\forall j | a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
end
end
    
```

QWERTZ

Proving Correctness of Loops: Exercise (1.3)



Prove that each of the following **Hoare Triples** is TRUE.

4. Loop Variant Stays Non-Negative Before Exit:

$$\{ (\forall j | a.lower \leq j < i \bullet Result \geq a[j]) \wedge \neg(i > a.upper) \}$$

```

  if a [i] > Result then Result := a [i] end
  i := i + 1
{ a.upper - i + 1 ≥ 0 }
    
```

5. Loop Variant Keeps Decrementing before Exit:

$$\{ (\forall j | a.lower \leq j < i \bullet Result \geq a[j]) \wedge \neg(i > a.upper) \}$$

```

  if a [i] > Result then Result := a [i] end
  i := i + 1
{ a.upper - i + 1 < (a.upper - i + 1)0 }
    
```

where $(a.upper - i + 1)_0 \equiv a.upper_0 - i_0 + 1$

QWERTZ

Proving Correctness of Loops: Exercise (1.2)



Prove that each of the following **Hoare Triples** is TRUE.

1. Establishment of Loop Invariant:

$$\{ True \}$$

```

  i := a.lower
  Result := a[i]
{  $\forall j | a.lower \leq j < i \bullet Result \geq a[j]$  }
    
```

2. Maintenance of Loop Invariant:

$$\{ (\forall j | a.lower \leq j < i \bullet Result \geq a[j]) \wedge \neg(i > a.upper) \}$$

```

  if a [i] > Result then Result := a [i] end
  i := i + 1
{  $\forall j | a.lower \leq j < i \bullet Result \geq a[j]$  }
    
```

3. Establishment of Postcondition upon Termination:

$$(\forall j | a.lower \leq j < i \bullet Result \geq a[j]) \wedge i > a.upper$$

$$\Rightarrow \forall j | a.lower \leq j \leq a.upper \bullet Result \geq a[j]$$

QWERTZ

Proof Tips (1)



$$\{ Q \} S \{ R \} \Rightarrow \{ Q \wedge P \} S \{ R \}$$

In order to prove $\{ Q \wedge P \} S \{ R \}$, it is sufficient to prove a version with a **weaker** precondition: $\{ Q \} S \{ R \}$.

Proof:

- Assume: $\{ Q \} S \{ R \}$

It's equivalent to assuming: $\boxed{Q} \Rightarrow wp(S, R)$

(A1)

- To prove: $\{ Q \wedge P \} S \{ R \}$

- It's equivalent to proving: $Q \wedge P \Rightarrow wp(S, R)$
- Assume: $Q \wedge P$, which implies \boxed{Q}
- According to (A1), we have $wp(S, R)$. ■

QWERTZ

Proof Tips (2)



When calculating $wp(s, R)$, if either program s or postcondition R involves array indexing, then R should be augmented accordingly.

e.g., Before calculating $wp(s, a[i] > 0)$, augment it as

$$wp(s, a.lower \leq i \leq a.upper \wedge a[i] > 0)$$

e.g., Before calculating $wp(x := a[i], R)$, augment it as

$$wp(x := a[i], a.lower \leq i \leq a.upper \wedge R)$$

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