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EECS 1028 Z. Problem Set No3 —Solutions

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1. (4 MARKS) Show that if $\mathbb{F}$ is a function and $\text{dom}(\mathbb{F})$ is a set then $\mathbb{F}$ is a set.

Proof. Note that $\text{ran}(\mathbb{F})$ is also a set —by the function (forward) image theorem 5.1.10— since

$$\text{ran}(\mathbb{F}) = \mathbb{F}[\text{dom}(\mathbb{F})]$$

We are done by the subclass theorem and $\mathbb{F} \subseteq \text{dom}(\mathbb{F}) \times \text{ran}(\mathbb{F})$. $\square$

2. (3 MARKS) Show by an easy counterexample that “if $\mathbb{F}$ is a function and $\text{ran}(\mathbb{F})$ is a set then $\mathbb{F}$ is a set” is false.

Answer. Define the function $\mathbb{F}$ by $\mathbb{F}(x) = 0$ for all $x \in \mathbb{U}$. Then $\text{ran}(\mathbb{F}) = \{0\}$, a set, but $\text{dom}(\mathbb{F}) = \mathbb{U}$, a proper class. $\square$

3. (4 MARKS) Define a **DIFFERENT implementation**

$$\langle x, y \rangle \stackrel{\text{Def}}{=} \left\{ \{x\}, \{x, y\} \right\}$$

for **ordered pair** where this time we denote the latter as “ $\langle x, y \rangle$ ” (angular brackets).

Prove that

$$\langle x, y \rangle = \langle x', y' \rangle \rightarrow x = x' \wedge y = y' \quad (1)$$

Caution. This does NOT require arguments via “**set formation by stages**”.

Proof. Assume **LHS of (1)**. Then, applying the function “ $\cap$ ” to the same input $\langle x, y \rangle = \langle x', y' \rangle$, that is,

$$\cap \left\{ \{x\}, \{x, y\} \right\} = \cap \left\{ \{x'\}, \{x', y'\} \right\}$$

we get $\{x\} = \{x'\}$, hence

$$x = x' \quad (2)$$

.

Next, applying the function “ $\cup$ ” to the same input $\langle x, y \rangle = \langle x', y' \rangle$, that is,

$$\bigcup \{ \{x\}, \{x, y\} \} = \bigcup \{ \{x'\}, \{x', y'\} \}$$

we get (mindful of (2))

$$\{x, y\} = \{x, y'\} \quad (3)$$

.

We now have two cases:

- Case $x = y$. Then $y' \in \{x, y\} = \{y\}$, hence $y = y'$ which along with (2) settles RHS of (1).
- Case $x \neq y$. Now, $y \in \{x, y'\}$ by (3). Since the Case forbids $y = x$ we must have the only alternative, $y = y'$. This, along with (2), proves the RHS of (1) one last time.

□

4. (a) (2 MARKS) State the definition given in Class/NOTES/Text for

$$\boxed{A \text{ is countable}} \quad (2)$$

Definition: A is countable **iff** we have some **onto** (not necessarily total) $g : \mathbb{N} \rightarrow A$.

- (b) (4 MARKS) Prove that the Definition from Class/NOTES/Text is **equivalent to**

$$\boxed{A \text{ is countable } \textcolor{blue}{\text{iff}}, \text{ there is a } \underline{\text{total}} \text{ and } \underline{1-1} \ f : A \rightarrow \mathbb{N}} \quad (3)$$

Proof.

- i. Suppose A is **countable** according to Definition 4a. That means some (not necessarily total) onto function $g : \mathbb{N} \rightarrow A$ exists.

By Theorem 5.1.30 of the Notes, we also have a **total** and **1-1** $f : A \rightarrow \mathbb{N}$. (The theorem from Notes adds: “such that $gf = 1_A$ ” but this part is **irrelevant** to this problem).

- ii. Conversely, suppose that we have $f : A \rightarrow \mathbb{N}$ that is total and 1-1. By Theorem 5.1.29 of the Notes, there is an onto function $g : \mathbb{N} \rightarrow A$, so A is countable by 4a. (The quoted theorem from Notes also observes: “such that $gf = 1_A$ ” but this part is irrelevant to this problem). $\square$

5. (4 MARKS) Prove transitivity of $\sim$, that is, if $A \sim B \sim C$, then $A \sim C$.

Proof. Hypothesis says that

- (a) We have a total, 1-1 and onto $f : A \rightarrow B$.
- (b) We have a total, 1-1 and onto $g : B \rightarrow C$.

To prove transitivity we need to prove that

$$gf : A \rightarrow C$$

is a 1-1 correspondence.

- (1) gf is onto C . Indeed, let $z \in C$. By onto-ness of g there is a $y \in B$ such that $g(y) = z$. By onto-ness of f there is a $x \in A$ such that $f(x) = y$.
So, $(gf)(x) = g(f(x)) = g(y) = z$. So gf is onto.
- (2) Prove gf is total on A . OK, show $(gf)(a)$ is defined for any $a \in A$.
But $(gf)(a) = g(f(a))$. Since f is total, we have $f(a)$ is an object in B . But g is total, so $g(f(a))$ is an object in C .
- (3) Prove (gf) is 1-1. Let

$$(gf)(x) = (gf)(y) = z \tag{\dagger}$$

We shall show $x = y$.

Now, $(\dagger)$ says $g(f(x)) = g(f(y)) = z$. But g is 1-1, hence

$$f(x) = f(y) = w \tag{\ddagger}$$

the w since f is total. Now we remember that f is 1-1 too. So $x = y$ by $(\ddagger)$ and we proved 1-1ness of (gf) . $\square$