Specifying & Refining a File Transfer Protocol

MEB: Chapter 4



EECS3342 Z: System Specification and Refinement Winter 2022

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Learning Outcomes



This module is designed to help you review:

- What a Requirement Document (RD) is
- What a refinement is
- Writing <u>formal</u> specifications
 - o (Static) contexts: constants, axioms, theorems
 - o (Dynamic) machines: variables, invariants, events, guards, actions
- Proof Obligations (POs) associated with proving:
 - refinements
 - system properties
- Applying inference rules of the sequent calculus



A Different Application Domain

- The bridge controller we specified, refined, and proved exemplifies
 a reactive system, working with the physical world via:
 - sensorsactuators[a, b, c, ml_pass, il_pass][ml_tl, il_tl]
- We now study an example exemplifying a distributed program:
 - A protocol followed by two agents, residing on distinct geographical locations, on a computer network
 - Each file is transmitted asynchronously:
 bytes of the file do <u>not</u> arrive at the receiver all at one go.
 - Language of predicates, sets, and relations required
 - The <u>same</u> principles of generating *proof obligations* apply.



Requirements Document: File Transfer Protocol (FTP)

You are required to implement a system for transmitting files between *agents* over a computer network.



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Requirements Document: R-Descriptions

Each *R-Description* is an <u>atomic</u> *specification* of an intended *functionality* or a desired *property* of the working system.

REQ1	The protocol ensures the copy of a file from the sender to the receiver.
REQ2	The file is supposed to be made of a sequence of items.
REQ3	The file is sent piece by piece between the two sites.

Refinement Strategy



- Recall the <u>design</u> strategy of progressive <u>refinements</u>.
 - initial model (m₀): a file is transmitted from the sender to the receiver. [REQ1]
 However, at this most abstract model:
 - file transmitted from sender to receiver synchronously & instantaneously
 - transmission process abstracted away
 - **1.** 1st refinement $(m_1 \text{ refining } m_0)$:

transmission is done asynchronously

[REQ2, REQ3]

However, at this more concrete model:

- no communication between sender and receiver
- exchanges of messages and acknowledgements abstracted away
- **2. 2nd refinement** (m_2 **refining** m_1): communication mechanism elaborated

[REQ2, REQ3]

3. <u>final</u>, 3rd refinement (m₃ refining m₂): communication mechanism optimized

[REQ2, REQ3]

• Recall *Correct by Construction*:

From each *model* to its *refinement*, only a <u>manageable</u> amount of details are added, making it *feasible* to conduct **analysis** and **proofs**.

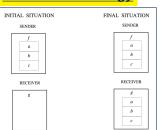




- In this most abstract perception of the protocol, we do not consider the sender and receiver:
 - residing in geographically distinct locations
 - communicating via message exchanges
- Instead, we focus on this single requirement:

REQ1 The protocol ensures the copy of a file from the sender to the receiver.

Abstraction Strategy:



- Observe the system with the process of transmission abstracted away
- <u>only</u> meant to inform what the protocol is supposed to achieve
- <u>not</u> meant to detail <u>how</u> the transmission is achieved

Math Background Review



Refer to LECTURE 1 for reviewing:

- Predicates
- Sets
- Relations and Operations
- Functions

[e.g., ∀]

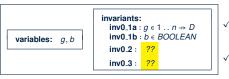


Model m_0 : Abstract State Space

 The <u>static</u> part formulates the *file* (from the *sender*'s end) as a sequence of data items:

```
sets: D,BOOLEAN constants: n,f axioms: axm0.1: n>0 axm0.2: f \in 1... n \rightarrow D axm0.3: BOOLEAN = {TRUE, FALSE}
```

2. The **dynamic** part of the state consists of two **variables**:



- ✓ g: file from the receiver's end
- √ b: whether or not the transmission is completed
- ✓ inv0_1a and inv0_1b are typing constraints.
- √ inv0_2 specifies what happens before the transmission
- ✓ inv0_3 specifies what happens after the transmission

LASSONDE

Model m_0 : State Transitions via Events

- The system acts as an ABSTRACT STATE MACHINE (ASM): it evolves as actions of enabled events change values of variables, subject to invariants.
- Initially, before the transmission:



- Nothing has been transmitted to the receiver.
- The transmission process has not been completed.
- Finally, <u>after</u> the transmission:



- The entire file f has been transmitted to the receiver.
- The *transmission* process has been completed.
- o In this abstract model:
 - Think of the transmission being <u>instantaneous</u>.
 - A later refinement specifies how f is transmitted asynchronously.

PO of Invariant Establishment



• How many **sequents** to be proved?

- [# invariants]
- We have <u>four</u> **sequents** generated for **event** init of model m_0 :

```
n > 0
       f \in 1 ... n \rightarrow D
       BOOLEAN = { TRUE, FALSE }
                                                init/inv0 1a/INV
1.
       \emptyset \in 1 ... n \rightarrow D
      n > 0
       f \in 1 ... n \rightarrow D
2.
       BOOLEAN = {TRUE, FALSE}
                                                init/inv0_1b/INV
       FALSE ∈ BOOLEAN
       n > 0
       f \in 1 ... n \rightarrow D
3.
       BOOLEAN = {TRUE, FALSE}
                                                init/inv0_2/INV
       FALSE = FALSE \Rightarrow \emptyset = \emptyset
       n > 0
       f \in 1 n \rightarrow D
       BOOLEAN = {TRUE, FALSE}
                                                init/inv0 3/INV
       FALSE = TRUE \Rightarrow \emptyset = f
```

Exercises: Prove the above sequents related to invariant establishment.

PO of Invariant Preservation



- How many **sequents** to be proved? [# non-init events × # invariants]
- We have four sequents generated for event final of model m₀:

```
n > 0
f \in 1 n \rightarrow D
BOOLEAN = {TRUE, FALSE}
q \in 1 ... n \rightarrow D
b ∈ BOOLEAN
b = FALSE \Rightarrow a = \emptyset
b = TRUE \Rightarrow q = f
b = FALSE
f \in 1 ... n \rightarrow D
```

final/inv0 1a/INV

```
n > 0
f \in 1 n \rightarrow D
BOOLEAN = { TRUE, FALSE}
q \in 1 ... n \rightarrow D
b ∈ BOOLEAN
b = FALSE \Rightarrow a = \emptyset
b = TRUE \Rightarrow q = f
b = FALSE
TRUE ∈ BOOLEAN
```

final/inv0 1b/INV

```
n > 0
f \in 1 ... n \rightarrow D
BOOLEAN = { TRUE, FALSE}
a \in 1... n \rightarrow D
b ∈ BOOLEAN
b = FALSE \Rightarrow a = \emptyset
b = TRUE \Rightarrow a = f
b = FALSE
TRUE = FALSE \Rightarrow f = \emptyset
```

final/inv0_2/INV

```
n > 0
f \in 1 ... n \rightarrow D
BOOLEAN = { TRUE, FALSE }
a \in 1...n \rightarrow D
b ∈ BOOLEAN
b = FALSE \Rightarrow a = \emptyset
b = TRUE \Rightarrow a = f
b = FALSE
TRIJF = TRIJF \Rightarrow f = f
```

final/inv0_3/INV

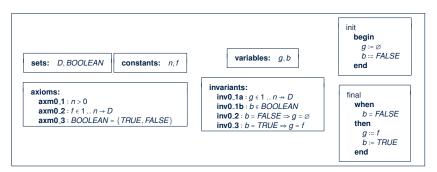
Exercises: Prove the above sequents related to invariant preservation.





- Our *initial model m*₀ is *provably correct* w.r.t.:
 - Establishment of *Invariants*
 - Preservation of *Invariants*
 - o Deadlock Freedom
- Here is the **specification** of m_0 :

[EXERCISE]

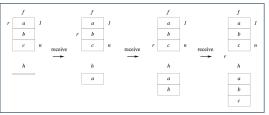


Model m_1 : "More Concrete" Abstraction



- In m₀, the transmission (evt. final) is synchronous and instantaneous.
- The 1st refinement has a more concrete perception of the file transmission:
 - The sender's file is coped gradually, *element by element*, to the receiver.
 - \rightarrow Such progress is denoted by occurrences of a *new event* receive.

h: elements transmitted so far
r: index of element to be sent
abstract variable g is replaced
by concrete variables h and r.



- o Nonetheless, communication between two agents remain abstracted away!
- That is, we focus on these two *intended functionalities*:

REQ2	The file is supposed to be made of a sequence of items.
REQ3	The file is sent piece by piece between the two sites.

• We are **obliged to prove** this **added concreteness** is **consistent** with m_0 .



Model m_1 : Refined, Concrete State Space

1. The **static** part remains the same as m_0 :

```
sets: D,BOOLEAN constants: n,f axioms: axm0.1: n > 0 axm0.2: f \in 1...n \rightarrow D axm0.3: BOOLEAN = \{TRUE, FALSE\}
```

- 2. The **dynamic** part formulates the **gradual** transmission process:
 - ♦ inv1_1: typing constraint
 - inv2_2: elements up to index r 1 have been transmitted



- inv2.3: transmission completed <u>means</u> <u>no</u> more elements to be transmitted
- thm1_1: transmission completed <u>means</u> receiver has a complete copy of sender's file
- A theorem, once proved as derivable from invariants, needs <u>not</u> be proved for preservation by events.

Model m_1 : Property Provable from Invariants

To prove that a theorem can be derived from the invariants:

variables: b, h, r

invariants:

*inv1*_1: $r \in 1 ... n + 1$ *inv1*_2: $h = (1 ... r - 1) \triangleleft f$ *inv1*_3: $b = TRUE \Rightarrow r = n + 1$ *thm1*_1: $b = TRUE \Rightarrow h = f$

We need to prove the following sequent:

$$r \in 1 ... n + 1$$

 $h = (1 ... r - 1) \lhd f$
 $b = TRUE \Rightarrow r = n + 1$
 \vdash
 $b = TRUE \Rightarrow h = f$

• Exercise: Prove the above sequent.



Model m_1 : Old and New Concrete Events

Initially, before the transmission:



- ♦ The *transmission* process has not been completed.
- Nothing has been transmitted to the *receiver*.
- ⋄ First file element is available for transmission.
- While the transmission is <u>ongoing</u>:



- While sender has more file elements available for transmission:
 - Next file element is received and accumulated to the receiver's copy.
 - Sender's next available file element is updated.
- ♦ In this concrete model:
 - Receiver having access to sender's private variable r is <u>unrealistic</u>.
 - A later refinement specifies how two agents communicate.
- Finally, <u>after</u> the transmission:



- When sender has no more file element available for transmission:
 - The transmission process is marked as completed.

PO of Invariant Establishment



• How many **sequents** to be proved?

- [# invariants]
- We have three **sequents** generated for **event** init of model m_1 :

n>0

$$f \in 1 ... n \rightarrow D$$

BOOLEAN = {TRUE, FALSE}
 \vdash
 $1 \in 1 ... n + 1$

init/inv1_1/INV

2.
$$| f \in 1 ... n \rightarrow D$$

$$BOOLEAN = \{TRUE, FALSE\}$$

$$\vdash \\ \emptyset \in (1...1-1) \triangleleft f$$

init/inv1_2/INV

$$n > 0$$

$$f \in 1 ... n \rightarrow D$$
3. BOOLEAN = {TRUE, FALSE}
$$FALSE = TRUE \Rightarrow 1 = n + 1$$

init/inv1_3/INV

• Exercises: Prove the above sequents related to invariant establishment.

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PO of Invariant Preservation - final

- We have three **sequents** generated for **old event** final of model m_1 .
- Here is one of the sequents:

```
n > 0
f \in 1 ... n \rightarrow D
BOOLEAN = {TRUE, FALSE}
g \in 1 ... n \rightarrow D
b ∈ BOOLEAN
b = FALSE \Rightarrow g = \emptyset
b = TRUE \Rightarrow q = f
r \in 1...n + 1
h = (1 \dots r - 1) \triangleleft f
b = TRUE \Rightarrow r = n + 1
b = FALSE
r = n + 1
r \in 1 \dots n+1
```

final/inv1_1/INV

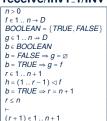
• Exercises: Formulate & prove other sequents of *invariant preservation*.



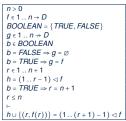
PO of Invariant Preservation - receive

• We have three **sequents** generated for **new event** receive of model m_1 :

receive/inv1_1/INV



receive/inv1_2/INV



receive/inv1_3/INV

```
n > 0

f \in 1 ... n \rightarrow D

BOOLEAN = \{TRUE, FALSE\}

g \in 1 ... n \rightarrow D

b \in BOOLEAN

b = FALSE \Rightarrow g = \emptyset

b = TRUE \Rightarrow g = f

r \in 1 ... n + 1

h = (1 ... r - 1) \lor f

b = TRUE \Rightarrow r = n + 1

r \le n

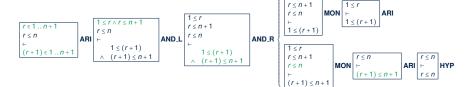
b = TRUE \Rightarrow (r + 1) = n + 1
```

Exercises: Prove the above sequents of invariant preservation.



Proving Refinement: receive/inv1_1/INV

MON





Proving Refinement: receive/inv1_2/INV

```
 \begin{array}{c} n > 0 \\ f \in 1 \dots n \to D \\ f \in 1 \dots n \to D \\ BOOLEAN = \{TRUE, FALSE\} \\ g \in 1 \dots n \to D \\ b \in BOOLEAN \\ b = FALSE \to g = \emptyset \\ b = TRUE \Rightarrow g = f \\ r \in 1 \dots n + 1 \\ h = (1 \dots r - 1) \lhd f \\ b = TRUE \Rightarrow r = n + 1 \\ r \le n \\ \vdash U \cup \{(r, f(r))\} = (1 \dots (r + 1) - 1) \lhd f \\ \end{array}
```

MON

```
\begin{cases} f \in 1 \dots n \to D \\ r \in 1 \dots n+1 \\ h = (1 \dots r-1) \lhd f \\ r \le n \\ \vdash h \cup \{(r,f(r))\} = (1 \dots (r+1)-1) \lhd f \end{cases}
```

```
|\mathbf{R}| \begin{cases} f \in 1 ... n \to D \\ 1 \le r \\ h = (1 ... r - 1) \lhd f \\ r \le n \\ h \cup \{(r, f(r))\} = (1 ... (r+1) - 1) \lhd f \end{cases}
```

```
EQ_LR,
MON,
ARI
```

```
ARI

\begin{cases}
f \in 1 ... n \to D \\
1 \le r \\
r \le n
\end{cases}

(1 ... r - 1) \lhd f \cup \{(r, f(r))\} = (1 ... r) \lhd f
```

Proving Refinement: receive/inv1_3/INV



```
 \begin{array}{ll} n\!>\!0\\ f\in\!1...n\to D\\ BOOLEAN=\{TRUE,FALSE\}\\ g\in\!1...n\to D\\ b\in\!BOOLEAN\\ b\in\!BOOLEAN\\ b\in\!FALSE\Rightarrow g=\varnothing\\ b=TRUE\Rightarrow g=f\\ r\in\!1...n+1\\ h=\{1...r-1\}\lor df\\ b=TRUE\Rightarrow r=n+1\\ r\le n\\ \vdash\\ b=TRUE\Rightarrow (r+1)=n+1\\ \end{array}
```

MON



IMP_R $b = TRUE \Rightarrow r = n + 1$ $r \le n$ b = TRUE(r + 1) = n + 1



EQ_LR, MON



ARI, MON

```
II.  \begin{bmatrix} \bot \\ \vdash \\ ((n+1)+1) = n+1 \end{bmatrix}  FALSE_L
```



m_1 : PO of Convergence of New Events

- Recall:
 - Interleaving of new events charactered as an integer expression: variant.
 - \circ A variant V(c, w) may refer to constants and/or *concrete* variables.
 - For m_1 , let's try **variants**: n + 1 r
- Accordingly, for the <u>new</u> event <u>receive</u>:

```
n > 0

f \in 1 ... n \rightarrow D

BOOLEAN = \{TRUE, FALSE\}

g \in 1 ... n \rightarrow D

b \in BOOLEAN

b = FALSE \Rightarrow g = \emptyset

b = TRUE \Rightarrow g = f

r \in 1 ... n + 1

h = (1 ... r - 1) \lhd f

b = TRUE \Rightarrow r = n + 1

r \le n

h = (r + 1) \lhd (r + 1) \lhd (r + 1) = r
```

receive/VAR

Exercises: Prove receive/VAR and Formulate/Prove receive/NAT.



[init]

First Refinement: Summary

- The *first refinement m*₁ is *provably correct* w.r.t.:
 - Establishment of Concrete Invariants
 - Preservation of Concrete Invariants
 - Strengthening of *quards*
 - Convergence (a.k.a. livelock freedom, non-divergence) [new events, EXERCISE]
 - Relative **Deadlock** Freedom
- [old events, EXERCISE] [EXERCISE]

[old & new events]

Here is the **specification** of m₁:

```
sets: D. BOOLEAN
                                constants: n.f
   axioms:
     axm0.1: n > 0
     axm0 2: f \in 1...n \rightarrow D
     axm0_3: BOOLEAN = {TRUE, FALSE}
                 invariants:
variables:
                   inv1 1: r \in 1...n + 1
                   inv1_2: h = (1...r-1) \triangleleft f
  b.h.r
                   inv1 3: h = TRLIF \Rightarrow r = n + 1
                   thm1 1: b = TRUE \Rightarrow h = f
```

init begin h := FALSE $h := \emptyset$ r := 1end

final when r = n + 1h = FALSE then b := TRUE end

receive when r < nthen $h := h \cup \{(r, f(r))\}$ r := r + 1end

> variants: n + 1 - r





Learning Outcomes

A Different Application Domain

Requirements Document:

File Transfer Protocol (FTP)

Requirements Document: R-Descriptions

Refinement Strategy

Model m_0 : Abstraction

Math Background Review

Model m_0 : Abstract State Space

Model m_0 : State Transitions via Events

PO of Invariant Establishment



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PO of Invariant Preservation

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Model m_1 : Refined, Concrete State Space

Model m_1 : Property Provable from Invariants

Model m_1 : Old and New Concrete Events

PO of Invariant Establishment

PO of Invariant Preservation - final

PO of Invariant Preservation - receive

Proving Refinement: receive/inv1_1/INV

Proving Refinement: receive/inv1_2/INV



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Proving Refinement: receive/inv1_3/INV

 m_1 : PO of Convergence of New Events

First Refinement: Summary