

# Balanced Binary Search Trees



EECS2011 N & Z:  
Fundamentals of Data Structures  
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# Learning Outcomes of this Lecture

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This module is designed to help you understand:

- When the **Worst-Case RT** of a **BST Search** Occurs
- **Height-Balance** Property
- Performing **Rotations** to Restore Tree **Balance**

# Balanced Binary Search Trees: Motivation

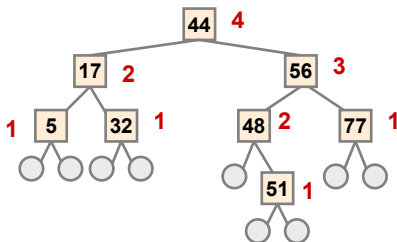
- After *insertions* into a BST, the **worst-case RT** of a *search* occurs when the *height*  $h$  is at its *maximum*:  **$O(n)$** :
  - e.g., Entries were inserted in an decreasing order of their keys  
 $\langle 100, 75, 68, 60, 50, 1 \rangle$   
⇒ **One-path, left-slanted** BST
  - e.g., Entries were inserted in an increasing order of their keys  
 $\langle 1, 50, 60, 68, 75, 100 \rangle$   
⇒ **One-path, right-slanted** BST
  - e.g., Last entry's key is in-between keys of the previous two entries  
 $\langle 1, 100, 50, 75, 60, 68 \rangle$   
⇒ **One-path, side-alternating** BST
- To avoid the worst-case RT ( $\therefore$  a *ill-balanced tree*), we need to take actions **as soon as** the tree becomes *unbalanced*.

# Balanced Binary Search Trees: Definition

- Given a node  $p$ , the **height** of the subtree rooted at  $p$  is:

$$\text{height}(p) = \begin{cases} 0 & \text{if } p \text{ is external} \\ 1 + \text{MAX} (\{ \text{height}(c) \mid \text{parent}(c) = p \}) & \text{if } p \text{ is internal} \end{cases}$$

- A **balanced** BST  $T$  satisfies the **height-balance property**:  
For every **internal node**  $n$ , **heights** of  $n$ 's child nodes differ  $\leq 1$ .



Q: Is the above tree a **balanced BST**? ✓

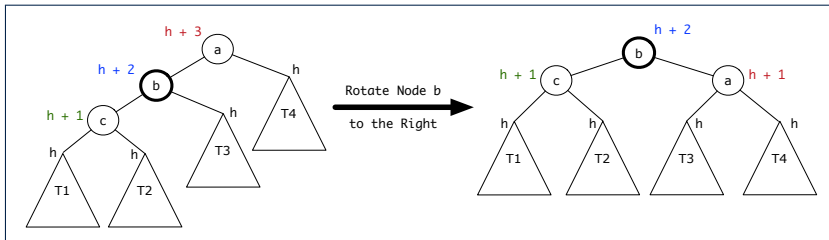
Q: Will the tree remain **balanced** after inserting 55? ✗

Q: Will the tree remain **balanced** after inserting 63? ✓

# Fixing Unbalanced BST: Rotations

A tree **rotation** is performed:

- When the latest **insertion/deletion** creates **unbalanced** nodes, along the **ancestor path** of the node being inserted/deleted.
- To change the **shape** of tree, **restoring** the **height-balance property**



**Q.** An **in-order traversal** on the resulting tree?

**A.** Still produces a sequence of **sorted keys**  $\langle T_1, c, T_2, b, T_3, a, T_4 \rangle$

- After **rotating** node **b** to the **right**:
    - Heights of **descendants** ( $b, c, T_1, T_2, T_3$ ) and **sibling** ( $T_4$ ) stay **unchanged**.
    - Height of **parent** ( $a$ ) is **decreased by 1**.
- $\Rightarrow$  **Balance** of node  $a$  was **restored** by the **rotation**.

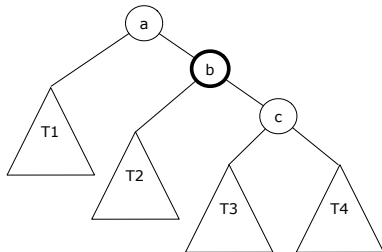
# After Insertions: Trinode Restructuring via Rotation(s)

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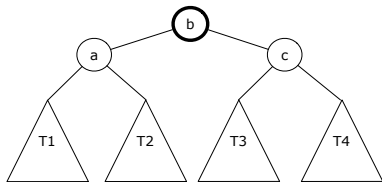
After *inserting* a new node  $n$ :

- **Case 1:** Nodes on  $n$ 's *ancestor path* remain *balanced*.  
⇒ No rotations needed
- **Case 2:** At least one of  $n$ 's *ancestors* becomes *unbalanced*.
  1. Get the first/lowest *unbalanced* node **a** on  $n$ 's *ancestor path*.
  2. Get  $a$ 's child node **b** in  $n$ 's *ancestor path*.
  3. Get  $b$ 's child node **c** in  $n$ 's *ancestor path*.
  4. Perform rotation(s) based on the *alignment* of  $a$ ,  $b$ , and  $c$ :
    - Slanted the *same* way ⇒ **single rotation** on the middle node **b**
    - Slanted *different* ways ⇒ **double rotations** on the lower node **c**

# Trinode Restructuring: Single, Left Rotation



After a **left rotation** on the middle node *b*:



**BST property** maintained?

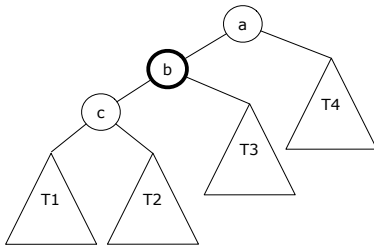
$\langle T_1, a, T_2, b, T_3, c, T_4 \rangle$

# Left Rotation

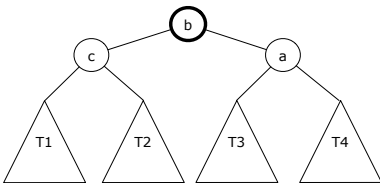
- **Insert** the following sequence of nodes into an empty BST:  
 $\langle 44, 17, 78, 32, 50, 88, 95 \rangle$
- Is the BST now **balanced**?
- **Insert** 100 into the BST.
- Is the BST still **balanced**?
- Perform a **left rotation** on the appropriate node.
- Is the BST again **balanced**?



# Trinode Restructuring: Single, Right Rotation



After a *right rotation* on the middle node *b*:



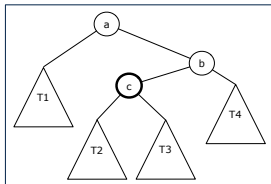
*BST property* maintained?

$\langle T_1, a, T_2, b, T_3, c, T_4 \rangle$

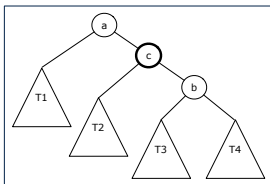
# Right Rotation

- **Insert** the following sequence of nodes into an empty BST:  
 $\langle 44, 17, 78, 32, 50, 88, 48 \rangle$
- Is the BST now **balanced**?
- **Insert** 46 into the BST.
- Is the BST still **balanced**?
- Perform a **right rotation** on the appropriate node.
- Is the BST again **balanced**?

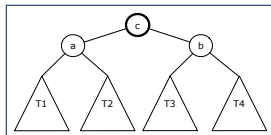
# Trinode Restructuring: Double, R-L Rotations



Perform a **Right Rotation** on Node *c*



Perform a **Left Rotation** on Node *c*



After Right-Left Rotations

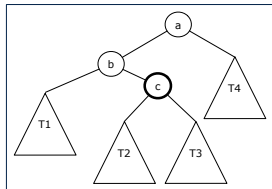
**BST property** maintained?

$\langle T_1, a, T_2, c, T_3, b, T_4 \rangle$

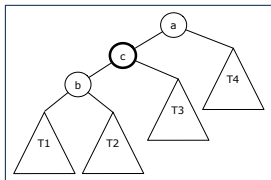
# R-L Rotations

- **Insert** the following sequence of nodes into an empty BST:  
 $\langle 44, 17, 78, 32, 50, 88, 82, 95 \rangle$
- Is the BST now **balanced**?
- **Insert** 85 into the BST.
- Is the BST still **balanced**?
- Perform the **R-L rotations** on the appropriate node.
- Is the BST again **balanced**?

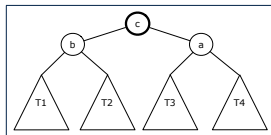
# Trinode Restructuring: Double, L-R Rotations



Perform a **Left Rotation** on Node c



Perform a **Right Rotation** on Node c



After Left-Right Rotations

**BST property** maintained?

$\langle T_1, b, T_2, c, T_3, a, T_4 \rangle$

# L-R Rotations

- **Insert** the following sequence of nodes into an empty BST:  
 $\langle 44, 17, 78, 32, 50, 88, 48, 62 \rangle$
- Is the BST now **balanced**?
- **Insert** 54 into the BST.
- Is the BST still **balanced**?
- Perform the **L-R rotations** on the appropriate node.
- Is the BST again **balanced**?

# After Deletions: Continuous Trinode Restructuring

- **Recall**: **Deletion** from a BST results in removing a node with zero or one **internal** child node.
- After **deleting** an existing node, say its child is  $n$ :
  - Case 1**: Nodes on  $n$ 's **ancestor path** remain **balanced**.  $\Rightarrow$  No rotations
  - Case 2**: At least one of  $n$ 's **ancestors** becomes **unbalanced**.
    1. Get the **first/lowest unbalanced** node  $a$  on  $n$ 's **ancestor path**.
    2. Get  $a$ 's **taller** child node  $b$ . [  $b \notin n$ 's **ancestor path** ]
    3. Choose  $b$ 's child node  $c$  as follows:
      - $b$ 's two child nodes have **different** heights  $\Rightarrow c$  is the **taller** child
      - $b$ 's two child nodes have **same** height  $\Rightarrow a, b, c$  slant the **same** way
    4. Perform rotation(s) based on the **alignment** of  $a, b$ , and  $c$ :
      - Slanted the **same** way  $\Rightarrow$  **single rotation** on the **middle** node  $b$
      - Slanted **different** ways  $\Rightarrow$  **double rotations** on the **lower** node  $c$
- As  $n$ 's **unbalanced ancestors** are found, keep applying **Case 2**, until **Case 1** is satisfied. [  $O(h) = O(\log n)$  **rotations** ]

# Single Trinode Restructuring Step

- **Insert** the following sequence of nodes into an empty BST:  
 $\langle 44, 17, 62, 32, 50, 78, 48, 54, 88 \rangle$
- Is the BST now **balanced**?
- **Delete** 32 from the BST.
- Is the BST still **balanced**?
- Perform a **left rotation** on the appropriate node.
- Is the BST again **balanced**?

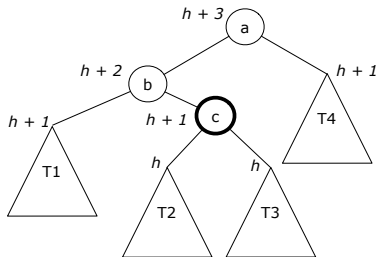


# Multiple Trinode Restructuring Steps

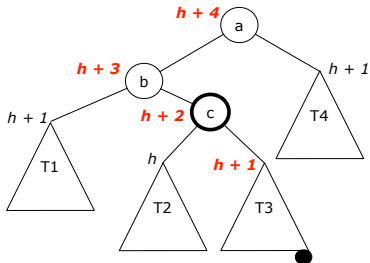
- **Insert** the following sequence of nodes into an empty BST:  
    {50, 25, 10, 30, 5, 15, 27, 1, 75, 60, 80, 55}
- Is the BST now **balanced**?
- **Delete** 80 from the BST.
- Is the BST still **balanced**?
- Perform a **right rotation** on the appropriate node.
- Is the BST now **balanced**?
- Perform another **right rotation** on the appropriate node.
- Is the BST again **balanced**?

# Restoring Balance from Insertions

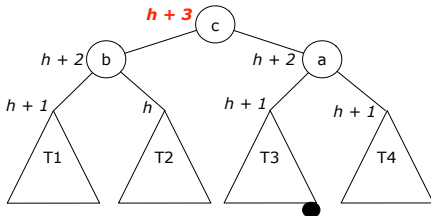
Before Insertion into T3



After Insertion into T3

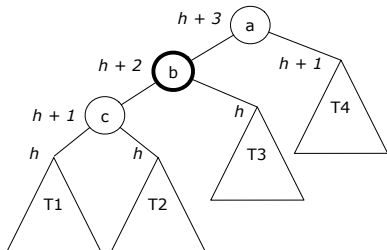


After Performing L-R Rotations on Node c: Height of Subtree Being Fixed Remains  $h + 3$

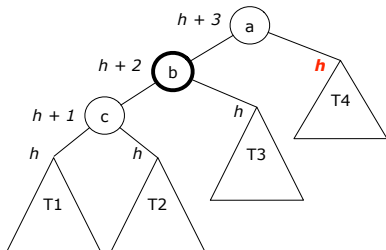


# Restoring Balance from Deletions

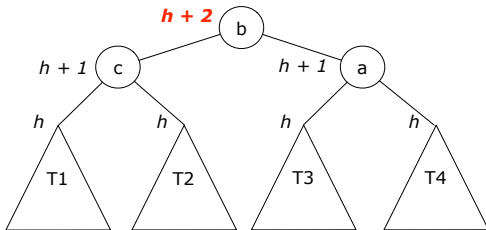
Before Deletion from T4



After Deletion from T4



After Performing Right Rotation on Node  $b$ : Height of Subtree Being Fixed **Reduces its Height by 1!**



# Restoring Balance: Insertions vs. Deletions

- Each **rotation** involves only **POs** of setting parent-child references.
  - ⇒  **$O(1)$**  running time for each tree **rotation**
- After each **insertion**, a **trinode restructuring** step can **restore the balance** of the subtree rooted at the first **unbalanced** node.
  - ⇒  **$O(1)$**  rotations suffices to restore the balance of tree
- After each **deletion**, one or more **trinode restructuring** steps may **restore the balance** of the subtree rooted at the first **unbalanced** node.
  - ⇒ May take  **$O(\log n)$**  rotations to restore the balance of tree

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