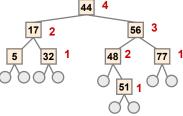
• After *insertions* into a BST, the *worst-case RT* of a *search* **Balanced Binary Search Trees** occurs when the *height* h is at its *maximum*: **O(n)**: • e.g., Entries were inserted in an decreasing order of their keys (100, 75, 68, 60, 50, 1)⇒ One-path, left-slanted BST • e.g., Entries were inserted in an increasing order of their keys EECS2011 N & Z: $\langle 1, 50, 60, 68, 75, 100 \rangle$ Fundamentals of Data Structures ⇒ One-path, right-slanted BST • e.g., Last entry's key is in-between keys of the previous two entries Winter 2022 (1, 100, 50, 75, 60, 68)CHEN-WEI WANG ⇒ One-path, side-alternating BST • To avoid the worst-case RT (:: a *ill-balanced tree*), we need to take actions as soon as the tree becomes unbalanced. 3 of 22 **Balanced Binary Search Trees: Definition** Learning Outcomes of this Lecture LASSONDE LASSONDE • Given a node p, the *height* of the subtree rooted at p is:

This module is designed to help you understand:

- When the Worst-Case RT of a BST Search Occurs
- Height-Balance Property
- Performing *Rotations* to Restore Tree *Balance*

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- Q: Is the above tree a *balanced BST*?
- Q: Will the tree remain balanced after inserting 55? Q: Will the tree remain *balanced* after inserting 63?



Balanced Binary Search Trees: Motivation

 $height(p) = \begin{cases} 0 & \text{if } p \text{ is external} \\ 1 + MAX \left(\{ height(c) \mid parent(c) = p \} \right) & \text{if } p \text{ is internal} \end{cases}$

• A *balanced* BST T satisfies the *height-balance property* :

For every *internal node n*, *heights* of *n*'s <u>child nodes</u> differ \leq **1**.

 \checkmark

if *p* is *external*

LASSONDE

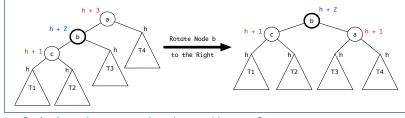
Fixing Unbalanced BST: Rotations



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A tree **rotation** is performed:

- When the latest insertion/deletion creates unbalanced nodes, along the ancestor path of the node being inserted/deleted.
- To change the shape of tree, restoring the height-balance property



- Q. An in-order traversal on the resulting tree?
- A. Still produces a sequence of *sorted keys* $(T_1, c, T_2, b, T_3, a, T_4)$
- After **rotating** node b to the right:
 - Heights of *descendants* (b, c, T₁, T₂, T₃) and *sibling* (T₄) stay *unchanged*.
 - Height of *parent* (a) is *decreased by 1*.
 - ⇒ Balance of node a was restored by the rotation.

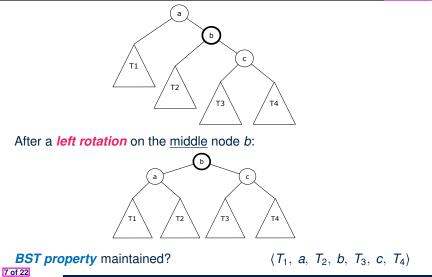
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After Insertions: Trinode Restructuring via Rotation(s)

After *inserting* a new node *n*:

- Case 1: Nodes on *n*'s ancestor path remain balanced.
 - \Rightarrow No rotations needed
- Case 2: At least one of n's ancestors becomes unbalanced. 1. Get the first/lowest unbalanced node a on n's ancestor path.
 - 2. Get a's child node *b* in n's ancestor path.
 - **3.** Get *b*'s child node *c* in *n*'s *ancestor path*.
 - **4.** Perform rotation(s) based on the *alignment* of *a*, *b*, and *c*:
 - Slanted the *same* way \Rightarrow *single rotation* on the *middle* node *b*
 - Slanted *different* ways \Rightarrow *double* rotations on the lower node c

Trinode Restructuring: Single, Left Rotation

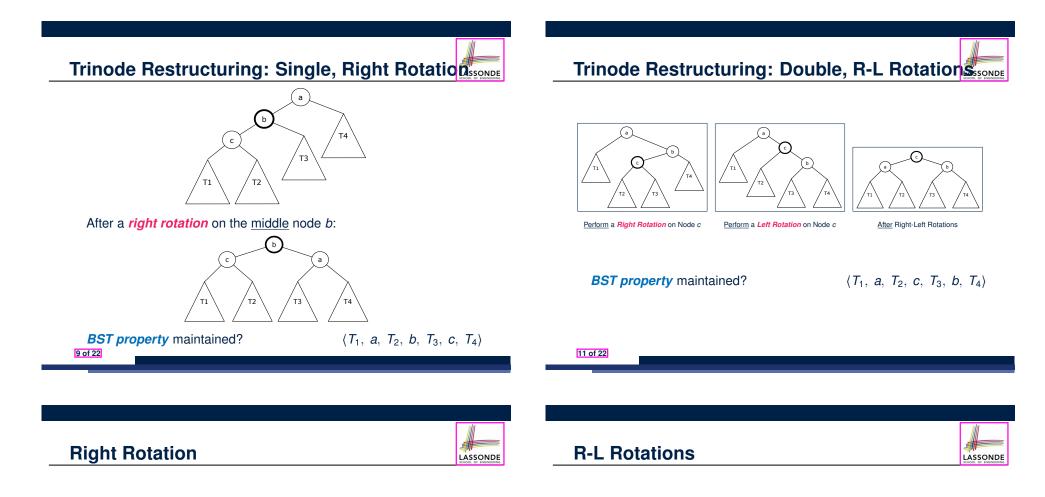




Left Rotation



- Insert the following sequence of nodes into an empty BST: (44, 17, 78, 32, 50, 88, 95)
- Is the BST now balanced?
- Insert 100 into the BST.
- Is the BST still balanced?
- Perform a *left rotation* on the appropriate node.
- Is the BST again balanced?



- Insert the following sequence of nodes into an empty BST: (44, 17, 78, 32, 50, 88, 48)
- Is the BST now balanced?
- Insert 46 into the BST.
- Is the BST still *balanced*?
- Perform a *right rotation* on the appropriate node.
- Is the BST again *balanced*?

- Insert the following sequence of nodes into an empty BST: $\langle 44, 17, 78, 32, 50, 88, 82, 95 \rangle$
- Is the BST now *balanced*?
- Insert 85 into the BST.
- Is the BST still *balanced*?
- Perform the *R-L rotations* on the appropriate node.
- Is the BST again *balanced*?

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Trinode Restructuring: Double, L-R Rotation After Deletions: LASSONDE **Continuous Trinode Restructuring** • *Recall* : *Deletion* from a BST results in removing a node with zero or one *internal* child node. • After *deleting* an existing node, say its child is n: Case 1: Nodes on *n*'s *ancestor path* remain *balanced*. ⇒ No rotations Case 2: At least one of n's ancestors becomes unbalanced. **1.** Get the **first/lowest** *unbalanced* node *a* on *n*'s *ancestor path*. 2. Get a's *taller* child node b. [b∉n's **ancestor path**] **3.** Choose *b*'s child node *c* as follows: Perform a Right Rotation on Node c After Left-Right Rotations Perform a *Left Rotation* on Node c • b's two child nodes have **different** heights \Rightarrow c is the **taller** child • b's two child nodes have **same** height $\Rightarrow a, b, c$ slant the **same** way **4.** Perform rotation(s) based on the *alignment* of *a*, *b*, and *c*: $(T_1, b, T_2, c, T_3, a, T_4)$ **BST property** maintained? • Slanted the *same* way \Rightarrow *single rotation* on the **middle** node *b* • Slanted *different* ways \Rightarrow *double* rotations on the lower node c • As n's unbalanced ancestors are found, keep applying Case 2, until Case 1 is satisfied. $O(h) = O(\log n)$ rotations 13 of 22 15 of 22 **L-R Rotations** Single Trinode Restructuring Step LASSONDE LASSONDE

- Insert the following sequence of nodes into an empty BST: (44, 17, 78, 32, 50, 88, 48, 62)
- Is the BST now balanced?
- Insert 54 into the BST.
- Is the BST still *balanced*?
- Perform the *L-R rotations* on the appropriate node.
- Is the BST again balanced?

- *Insert* the following sequence of nodes into an <u>empty</u> BST: (44, 17, 62, 32, 50, 78, 48, 54, 88)
- Is the BST now *balanced*?
- Delete 32 from the BST.
- Is the BST still *balanced*?
- Perform a *left rotation* on the appropriate node.
- Is the BST again *balanced*?

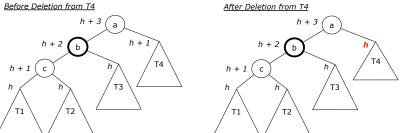
Multiple Trinode Restructuring Steps



- Insert the following sequence of nodes into an empty BST: (50, 25, 10, 30, 5, 15, 27, 1, 75, 60, 80, 55)
- Is the BST now balanced?
- Delete 80 from the BST.
- Is the BST still balanced?
- Perform a *right rotation* on the appropriate node.
- Is the BST now *balanced*?
- Perform another *right rotation* on the appropriate node.
- Is the BST again *balanced*?

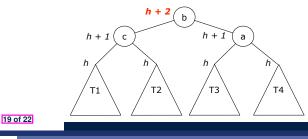
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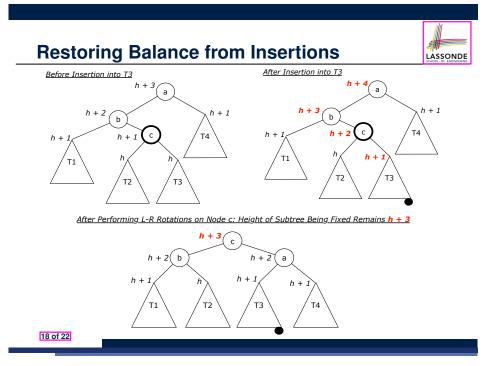




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After Performing Right Rotation on Node b: Height of Subtree Being Fixed Reduces its Height by 1!





Restoring Balance: Insertions vs. Deletions

- Each *rotation* involves only *POs* of setting parent-child references.
 ⇒ *O*(1) running time for each tree *rotation*
- After each *insertion*, a *trinode restructuring* step can *restore the balance* of the subtree rooted at the <u>first</u> *unbalanced* node.
 - \Rightarrow **O(1)** rotations suffices to restore the balance of tree
- After each *deletion*, <u>one or more</u> *trinode restructuring* steps may *restore the balance* of the subtree rooted at the <u>first</u> *unbalanced* node.
 - \Rightarrow May take *O(log n)* rotations to restore the balance of tree

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LASSONDE

Learning Outcomes of this Lecture

Balanced Binary Search Trees: Motivation

Balanced Binary Search Trees: Definition

Fixing Unbalanced BST: Rotations

After Insertions:

Trinode Restructuring via Rotation(s)

Trinode Restructuring: Single, Left Rotation

Left Rotation

Trinode Restructuring: Single, Right Rotation

Right Rotation

Trinode Restructuring: Double, R-L Rotations

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R-L Rotations

Trinode Restructuring: Double, L-R Rotations

L-R Rotations

After Deletions:

Continuous Trinode Restructuring

Single Trinode Restructuring Step

Multiple Trinode Restructuring Steps

Restoring Balance from Insertions

Restoring Balance from Deletions

Restoring Balance: Insertions vs. Deletions