

## Void Safety



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## Java Program: Example 2



```
1 class Point {
2   double x;
3   double y;
4   Point(double x, double y) {
5     this.x = x;
6     this.y = y;
7   }
8 }

1 class PointCollector {
2   ArrayList<Point> points;
3   PointCollector() {
4     points = new ArrayList<>();
5   }
6   void addPoint(Point p) {
7     points.add(p);
8   }
9   Point getPointAt(int i) {
10    return points.get(i);
11  }
```

```
1 @Test
2 public void test2() {
3   PointCollector pc = new PointCollector();
4   Point p = null;
5   pc.addPoint(p);
6   p = pc.getPointAt(0);
7   assertTrue(p.x == 3 && p.y == 4); }
```

The above Java code **compiles**. But anything wrong?

**L5** adds `p` (which stores `null`).

$\therefore$  **NullPointerException** when **L7** calls `p.x`.

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## Java Program: Example 1



```
1 class Point {
2   double x;
3   double y;
4   Point(double x, double y) {
5     this.x = x;
6     this.y = y;
7   }
8 }

1 class PointCollector {
2   ArrayList<Point> points;
3   PointCollector() { }
4   void addPoint(Point p) {
5     points.add(p);
6   }
7   Point getPointAt(int i) {
8     return points.get(i);
9   }
10 }
```

The above Java code **compiles**. But anything wrong?

```
1 @Test
2 public void test1() {
3   PointCollector pc = new PointCollector();
4   pc.addPoint(new Point(3, 4));
5   Point p = pc.getPointAt(0);
6   assertTrue(p.x == 3 && p.y == 4); }
```

**L3** calls `PointCollector` constructor not initializing `points`.

$\therefore$  **NullPointerException** when **L4** calls **L5** of `PointCollector`.

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## Java Program: Example 3



```
1 class Point {
2   double x;
3   double y;
4   Point(double x, double y) {
5     this.x = x;
6     this.y = y;
7   }
8 }

1 class PointCollector {
2   ArrayList<Point> points;
3   PointCollector() {
4     points = new ArrayList<>();
5   }
6   void addPoint(Point p) {
7     points.add(p);
8   }
9   Point getPointAt(int i) {
10    return points.get(i);
11  }
```

```
1 public void test3() {
2   PointCollector pc = new PointCollector();
3   Scanner input = new Scanner(System.in);
4   System.out.println("Enter an integer:");
5   int i = input.nextInt();
6   if(i < 0) { pc = null; }
7   pc.addPoint(new Point(3, 4));
8   assertTrue(pc.getPointAt(0).x == 3 && pc.getPointAt(0).y == 4);
9 }
```

The above Java code **compiles**. But anything wrong?

**NullPointerException** when user's input at **L5** is non-positive.

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## Limitation of Java's Type System



- A program that compiles does not guarantee that it is free from **NullPointerExceptions**:
  - Uninitialized attributes (in constructors).
  - Passing **nullable** variable as a method argument.
  - Calling methods on **nullable** local variables.
- The notion of Null references was back in 1965 in ALGO W.
- Tony Hoare (inventor of Quick Sort), introduced this notion of Null references “simply because *it was so easy to implement*”.
- But he later considers it as his “**billion-dollar mistake**”.
  - When your program manipulates reference/object variables whose types include the legitimate value of Null or Void, then there is always a possibility of having a **NullPointerExceptions**.
  - For undisciplined programmers, this means the final software product **crashes** often!

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## Eiffel Program: Example 1



```

1 class
2   POINT
3 create
4   make
5 feature
6   x: REAL
7   y: REAL
8 feature
9   make (nx: REAL; ny: REAL)
10  do x := nx
11    y := ny
12  end
13 end

1 class
2   POINT_COLLECTOR_1
3 create
4   make
5 feature
6   points: LINKED_LIST[POINT]
7 feature
8   make do end
9   add_point (p: POINT)
10  do points.extend (p) end
11  get_point_at (i: INTEGER): POINT
12  do Result := points [i] end
13 end
    
```

- Above code is semantically equivalent to Example 1 Java code.
- Eiffel compiler won't allow you to run it.
  - ∴ L8 of POINT\_COLLECTOR\_1 does **not compile**
  - ∴ It is **void safe** [Possibility of **NullPointerException** ruled out]

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## Eiffel's Type System for Void Safety



- By default, a reference variable is **non-detachable**.
  - e.g., `acc: ACCOUNT` means that `acc` is always **attached** to some valid ACCOUNT point.
- **VOID** is an illegal value for **non-detachable** variables.
  - ⇒ Scenarios that might make a reference variable **detached** are considered as **compile-time errors**:
    - **Non-detachable** variables can only be re-assigned to **non-detachable** variables.
      - e.g., `acc2: ACCOUNT` ⇒ `acc := acc2` **compilable**
      - e.g., `acc3: detachable ACCOUNT` ⇒ `acc := acc3` **non-compilable**
    - Creating variables (e.g., `create acc.make`) **compilable**
    - **Non-detachable** attribute not explicitly initialized (via creation or assignment) in all constructors is **non-compilable**.

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## Eiffel Program: Example 2



```

1 class
2   POINT
3 create
4   make
5 feature
6   x: REAL
7   y: REAL
8 feature
9   make (nx: REAL; ny: REAL)
10  do x := nx
11    y := ny
12  end
13 end

1 class
2   POINT_COLLECTOR_2
3 create
4   make
5 feature
6   points: LINKED_LIST[POINT]
7 feature
8   make do create points.make end
9   add_point (p: POINT)
10  do points.extend (p) end
11  get_point_at (i: INTEGER): POINT
12  do Result := points [i] end
13 end

1 test_2: BOOLEAN
2 local
3   pc: POINT_COLLECTOR_2 ; p: POINT
4 do
5   create pc.make
6   pc := Void
7   pc.add_point (p)
8   p := pc.get_point_at (0)
9   Result := p.x = 3 and p.y = 4
10 end
    
```

- Above code is semantically equivalent to Example 2 Java code.
  - L7 does **not compile** ∴ `pc` might be void. **[void safe]**

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## Eiffel Program: Example 3



```
1 class
2   POINT
3 create
4   make
5 feature
6   x: REAL
7   y: REAL
8 feature
9   make (nx: REAL; ny: REAL)
10  do x := nx
11    y := ny
12  end
13 end

1 class
2   POINT_COLLECTOR_2
3 create
4   make
5 feature
6   points: LINKED_LIST[POINT]
7 feature
8   make do create points.make end
9   add_point (p: POINT)
10  do points.extend (p) end
11  get_point_at (i: INTEGER): POINT
12  do Result := points [i] end
13 end

1 test_3: BOOLEAN
2 local pc: POINT_COLLECTOR_2 ; p: POINT ; i: INTEGER
3 do create pc.make
4   io.print ("Enter an integer:\N")
5   io.read_integer
6   if io.last_integer < 0 then pc := Void end
7   pc.add_point (create {POINT}.make (3, 4))
8   p := pc.get_point_at (0)
9   Result := p.x = 3 and p.y = 4
10 end
```

- Above code is semantically equivalent to Example 3 Java code. L7 and L8 do **not compile** ∴ pc might be void. [void safe]

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## Beyond this lecture...



- Tutorial Series on Void Safety by Bertrand Meyer (inventor of Eiffel):
  - The End of Null Pointer Dereferencing
  - The Object Test
  - The Type Rules
  - Final Rules
- Null Pointer as a Billion-Dollar Mistake by Tony Hoare
- More notes on void safety

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## Lessons from Void Safety



- It is much more costly to recover from **crashing** programs (due to **NullPointerException**) than to fix **uncompilable** programs.  
e.g., You'd rather have a **void-safe design** for an airplane, rather than hoping that the plane won't crash after taking off.
- If you are used to the standard by which Eiffel compiler checks your code for **void safety**, then you are most likely to write Java/C/C++/C#/Python code that is **void-safe** (i.e., free from **NullPointerExceptions**).

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