

**Types: Reference vs. Expanded
Copies: Reference vs. Shallow vs. Deep
Writing Complete Postconditions**



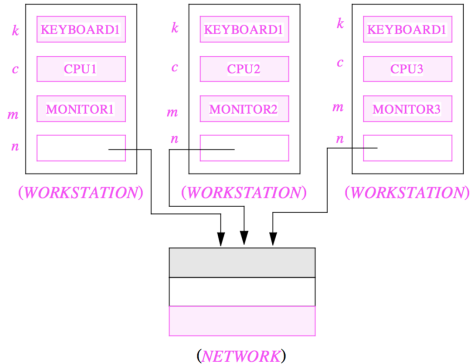
EECS3311 A: Software Design
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Expanded Class: Modelling

- We may want to have objects which are:
 - Integral parts of some other objects
 - Not** shared among objects

e.g., Each workstation has its own CPU, monitor, and keyboard.
All workstations share the same network.



Expanded Class: Programming (2)

```
class KEYBOARD ... end class CPU ... end  
class MONITOR ... end class NETWORK ... end  
class WORKSTATION  
  k: expanded KEYBOARD  
  c: expanded CPU  
  m: expanded MONITOR  
  n: NETWORK  
end
```

Alternatively:

```
expanded class KEYBOARD ... end  
expanded class CPU ... end  
expanded class MONITOR ... end  
class NETWORK ... end  
class WORKSTATION  
  k: KEYBOARD  
  c: CPU  
  m: MONITOR  
  n: NETWORK  
end
```

Expanded Class: Programming (3)

```
expanded class
  B
  feature
    change_i (ni: INTEGER)
      do
        i := ni
      end
  feature
    i: INTEGER
  end
```

```
1  test_expanded: BOOLEAN
2  local
3    eb1, eb2: B
4  do
5    Result := eb1.i = 0 and eb2.i = 0
6    check Result end
7    Result := eb1 = eb2
8    check Result end
9    eb2.change_i (15)
10   Result := eb1.i = 0 and eb2.i = 15
11   check Result end
12   Result := eb1 /= eb2
13   check Result end
14  end
```

- **L5:** object of expanded type is automatically initialized.
- **L9 & L10:** no sharing among objects of expanded type.
- **L7 & L12:** = between expanded objects compare their contents.

Reference vs. Expanded (1)

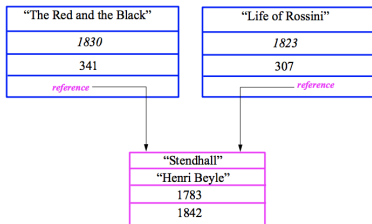
- Every entity must be declared to be of a certain type (based on a class).
- Every type is either *referenced* or *expanded*.
- In *reference* types:
 - y denotes *a reference* to some object
 - $x := y$ attaches x to same object as does y
 - $x = y$ compares references
- In *expanded* types:
 - y denotes *some object* (of expanded type)
 - $x := y$ copies contents of y into x
 - $x = y$ compares contents

$[x \sim y]$

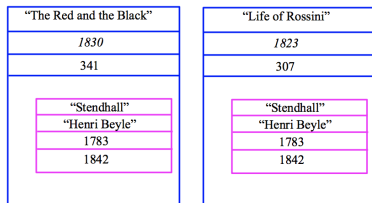
Reference vs. Expanded (2)

Problem: Every published book has an author. Every author may publish more than one books. Should the author field of a book *reference*-typed or *expanded*-typed?

reference-typed author



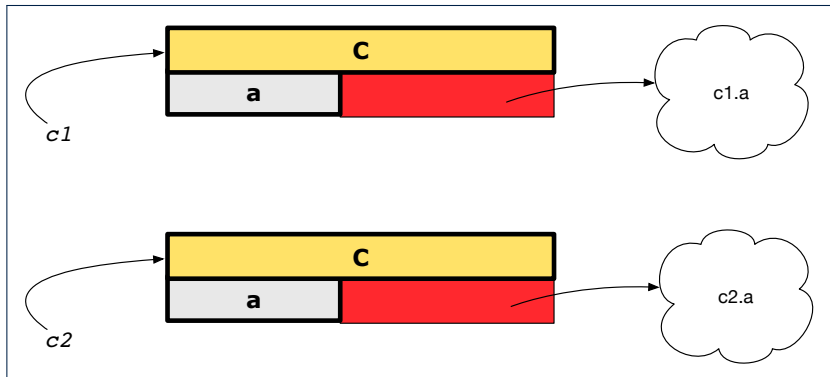
expanded-typed author



Copying Objects

Say variables `c1` and `c2` are both declared of type `C`. [`c1, c2: C`]

- There is only one attribute `a` declared in class `C`.
- `c1.a` and `c2.a` may be of either:
 - **expanded** type or
 - **reference** type



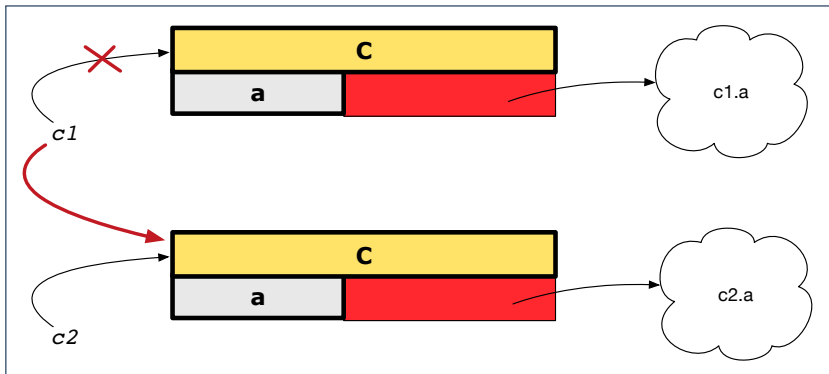
Copying Objects: Reference Copy

Reference Copy

```
c1 := c2
```

- Copy the address stored in variable `c2` and store it in `c1`.
 - ⇒ Both `c1` and `c2` point to the same object.
 - ⇒ Updates performed via `c1` also visible to `c2`.

[*aliasing*]

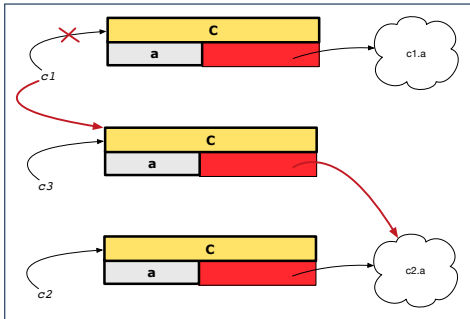


Copying Objects: Shallow Copy

Shallow Copy

```
c1 := c2.twin
```

- Create a temporary, behind-the-scene object $c3$ of type C .
- Initialize each attribute a of $c3$ via **reference copy**: $c3.a := c2.a$
- Make a **reference copy** of $c3$: $c1 := c3$
 $\Rightarrow c1$ and $c2$ **are not** pointing to the same object. $[c1 \neq c2]$
 $\Rightarrow c1.a$ and $c2.a$ **are** pointing to the same object.
 \Rightarrow **Aliasing** still occurs: at 1st level (i.e., attributes of $c1$ and $c2$)

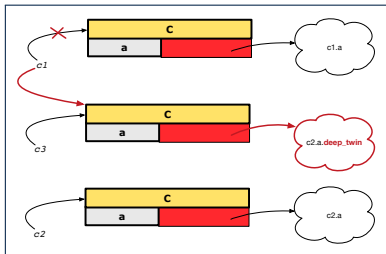


Copying Objects: Deep Copy

Deep Copy

```
c1 := c2.deep_twin
```

- Create a temporary, behind-the-scene object $c3$ of type C .
- **Recursively** initialize each attribute a of $c3$ as follows:
 - Base Case:** a is expanded (e.g., INTEGER). $\Rightarrow c3.a := c2.a$.
 - Recursive Case:** a is referenced. $\Rightarrow c3.a := c2.a.deep_twin$
- Make a **reference copy** of $c3$: $c1 := c3$
 - $\Rightarrow c1$ and $c2$ **are not** pointing to the same object.
 - $\Rightarrow c1.a$ and $c2.a$ **are not** pointing to the same object.
 - \Rightarrow **No aliasing** occurs at any levels.



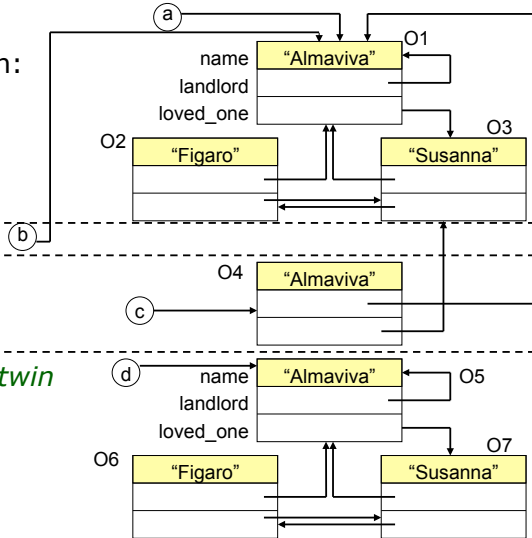
Copying Objects

- Initial situation:
- Result of:

$b := a$

$c := a.twin$

$d := a.deep_twin$



Example: Collection Objects (1)

- In any OOP, when a variable is declared of a **type** that corresponds to a **known class** (e.g., STRING, ARRAY, LINKED_LIST, etc.):

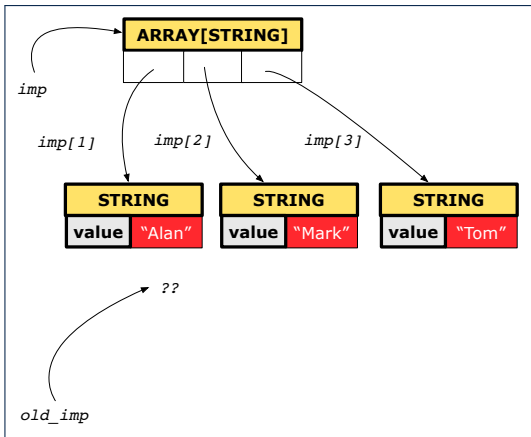
At **runtime**, that variable stores the **address** of an object of that type (as opposed to storing the object in its entirety).

- Assume the following variables of the same type:

```
...
local
  imp : ARRAY[STRING]
  old_imp: ARRAY[STRING]
do
  create {ARRAY[STRING]} imp.make_empty
  imp.force("Alan", 1)
  imp.force("Mark", 2)
  imp.force("Tom", 3)
  ...
```

Example: Collection Objects (2)

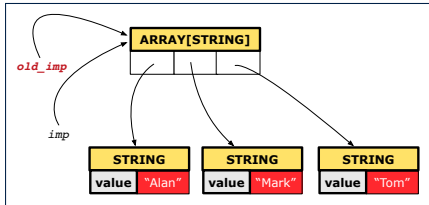
- Variables `imp` and `old_imp` store address(es) of some array(s).
- Each “slot” of these arrays stores a `STRING` object’s address.



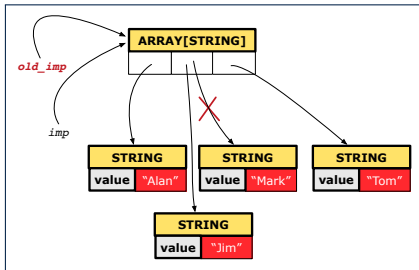
Reference Copy of Collection Object

```
1  old_imp := imp
2  Result := old_imp = imp -- Result = true
3  imp[2] := "Jim"
4  Result :=
5  across 1 |..| imp.count as j
6  all imp [j.item] ~ old_imp [j.item]
7  end -- Result = true
```

Before Executing L3



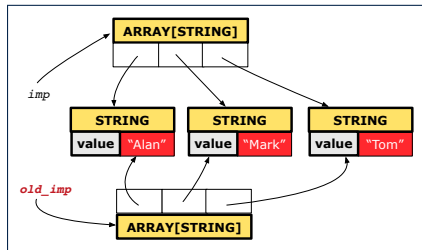
After Executing L3



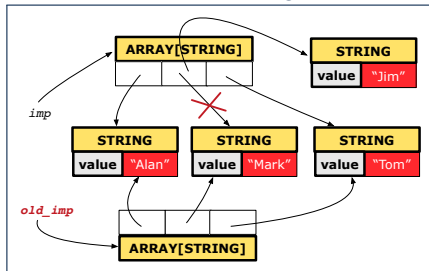
Shallow Copy of Collection Object (1)

```
1  old_imp := imp.twin
2  Result := old_imp = imp  -- Result = false
3  imp[2] := "Jim"
4  Result :=
5  across 1 |..| imp.count as j
6  all imp [j.item] ~ old_imp [j.item]
7  end  -- Result = false
```

Before Executing L3



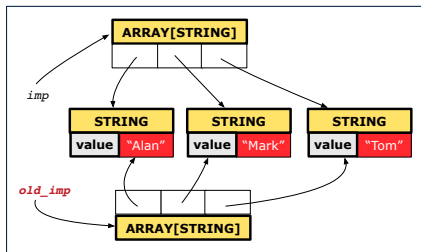
After Executing L3



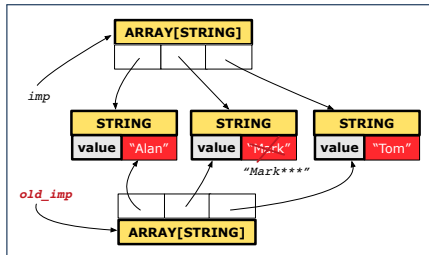
Shallow Copy of Collection Object (2)

```
1 old_imp := imp.twin
2 Result := old_imp = imp -- Result = false
3 imp[2].append ("****")
4 Result :=
5   across 1 |..| imp.count as j
6   all imp [j.item] ~ old_imp [j.item]
7   end -- Result = true
```

Before Executing L3



After Executing L3

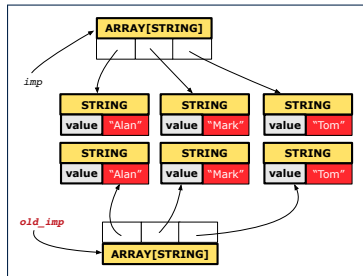


Deep Copy of Collection Object (1)

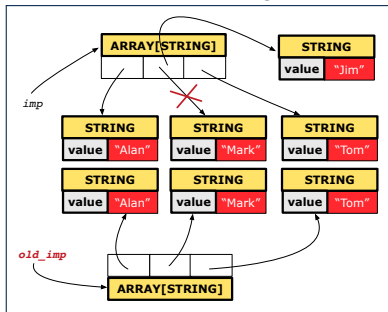
```

1  old_imp := imp.deep_twin
2  Result := old_imp = imp  -- Result = false
3  imp[2] := "Jim"
4  Result :=
5  across 1 |..| imp.count as j
6  all imp [j.item] ~ old_imp [j.item] end  -- Result = false
  
```

Before Executing L3



After Executing L3

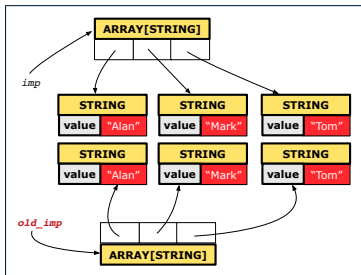


Deep Copy of Collection Object (2)

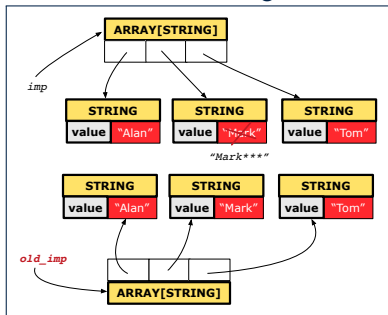
```

1  old_imp := imp.deep_twin
2  Result := old_imp = imp -- Result = false
3  imp[2].append ("****")
4  Result :=
5  across 1 |..| imp.count as j
6  all imp [j.item] ~ old_imp [j.item] end -- Result = false
  
```

Before Executing L3



After Executing L3



How are contracts checked at runtime?

- All contracts are specified as Boolean expressions.
- Right **before** a feature call (e.g., `acc.withdraw(10)`):
 - The current state of `acc` is called its **pre-state**.
 - Evaluate **pre-condition** using **current values** of attributes/queries.
 - Cache values, via `:=`, of **old expressions** in the **post-condition**.
 - e.g., `old balance = balance - a` [*old_balance* := *balance*]
 - e.g., `old accounts[i].id` [*old_accounts_i_id* := *accounts[i].id*]
 - e.g., `(old accounts)[i].id` [*old_accounts_i* := *accounts[i]*]
 - e.g., `(old accounts)[i].id` [*old_accounts* := *accounts*]
 - e.g., `(old Current).accounts[i].id` [*old_current* := *Current*]
- Right **after** the feature call:
 - The current state of `acc` is called its **post-state**.
 - Evaluate **invariant** using **current values** of attributes and queries.
 - Evaluate **post-condition** using both **current values** and **“cached” values** of attributes and queries.

When are contracts complete?

- In *post-condition*, for *each attribute*, specify the relationship between its *pre-state* value and its *post-state* value.
 - Eiffel supports this purpose using the **old** keyword.
- This is tricky for attributes whose structures are **composite** rather than **simple**:
 - e.g., *ARRAY*, *LINKED_LIST* are composite-structured.
 - e.g., *INTEGER*, *BOOLEAN* are simple-structured.
- **Rule of thumb:** For an attribute whose structure is composite, we should specify that after the update:
 1. The intended change is present; **and**
 2. *The rest of the structure is unchanged*.
- The second contract is much harder to specify:
 - Reference aliasing [ref copy vs. shallow copy vs. deep copy]
 - Iterable structure [use **across**]

Account

```
class
  ACCOUNT

inherit
  ANY
  redefine is_equal end

create
  make

feature -- Attributes
  owner: STRING
  balance: INTEGER

feature -- Commands
  make (n: STRING)
  do
    owner := n
    balance := 0
  end
```

```
deposit(a: INTEGER)
  do
    balance := balance + a
  ensure
    balance = old balance + a
  end

is_equal(other: ACCOUNT): BOOLEAN
  do
    Result :=
      owner ~ other.owner
      and balance = other.balance
  end
end
```

Bank

```
class BANK
create make
feature
  accounts: ARRAY[ACCOUNT]
  make do create accounts.make_empty end
  account_of (n: STRING): ACCOUNT
    require -- the input name exists
      existing: across accounts as acc some acc.item.owner ~ n end
      -- not (across accounts as acc all acc.item.owner /~ n end)
    do ...
    ensure Result.owner ~ n
    end
  add (n: STRING)
    require -- the input name does not exist
      non_existing: across accounts as acc all acc.item.owner /~ n end
      -- not (across accounts as acc some acc.item.owner ~ n end)
    local new_account: ACCOUNT
    do
      create new_account.make (n)
      accounts.force (new_account, accounts.upper + 1)
    end
  end
end
```

Roadmap of Illustrations

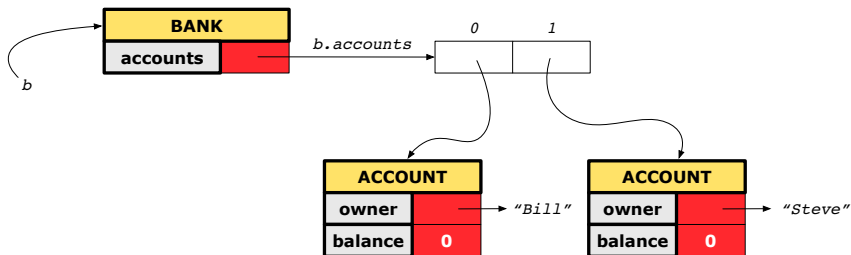
We examine 5 different versions of a command

deposit_on (*n* : *STRING*; *a* : *INTEGER*)

VERSION	IMPLEMENTATION	CONTRACTS	SATISFACTORY?
1	<i>Correct</i>	<i>Incomplete</i>	<i>No</i>
2	<i>Wrong</i>	<i>Incomplete</i>	<i>No</i>
3	<i>Wrong</i>	<i>Complete</i> (reference copy)	<i>No</i>
4	<i>Wrong</i>	<i>Complete</i> (shallow copy)	<i>No</i>
5	<i>Wrong</i>	<i>Complete</i> (deep copy)	<i>Yes</i>

Object Structure for Illustration

We will test each version by starting with the same runtime object structure:



Version 1: Incomplete Contracts, Correct Implementation

```
class BANK
  deposit_on_v1 (n: STRING; a: INTEGER)
    require across accounts as acc some acc.item.owner ~ n end
    local i: INTEGER
    do
      from i := accounts.lower
      until i > accounts.upper
      loop
        if accounts[i].owner ~ n then accounts[i].deposit(a) end
        i := i + 1
      end
    ensure
      num_of_accounts_unchanged:
        accounts.count = old accounts.count
      balance_of_n_increased:
        account_of (n).balance = old account_of (n).balance + a
    end
end
```

Test of Version 1

```
class TEST_BANK
  test_bank_deposit_correct_imp_incomplete_contract: BOOLEAN
  local
    b: BANK
  do
    comment("t1: correct imp and incomplete contract")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v1 ("Steve", 100)
    Result :=
      b.account_of ("Bill").balance = 0
      and b.account_of ("Steve").balance = 100
    check Result end
  end
end
```

Test of Version 1: Result

APPLICATION

Note: * indicates a violation test case

PASSED (1 out of 1)		
Case Type	Passed	Total
Violation	0	0
Boolean	1	1
All Cases	1	1
State	Contract Violation	Test Name
Test1	TEST_BANK	
PASSED	NONE	t1: test deposit_on with correct imp and incomplete contract

Version 2: Incomplete Contracts, Wrong Implementation

```
class BANK
  deposit_on_v2 (n: STRING; a: INTEGER)
    require across accounts as acc some acc.item.owner ~ n end
    local i: INTEGER
    do
      -- same loop as in version 1

      -- wrong implementation: also deposit in the first account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged:
        accounts.count = old accounts.count
      balance_of_n_increased:
        account_of (n).balance = old account_of (n).balance + a
    end
  end
end
```

Current postconditions lack a check that accounts other than n are unchanged.

Test of Version 2

```
class TEST_BANK
test_bank_deposit_wrong_imp_incomplete_contract: BOOLEAN
  local
    b: BANK
  do
    comment ("t2: wrong imp and incomplete contract")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v2 ("Steve", 100)
    Result :=
      b.account_of ("Bill").balance = 0
      and b.account_of ("Steve").balance = 100
    check Result end
  end
end
```

Test of Version 2: Result

APPLICATION

Note: * indicates a violation test case

FAILED (1 failed & 1 passed out of 2)		
Case Type	Passed	Total
Violation	0	0
Boolean	1	2
All Cases	1	2
State	Contract Violation	Test Name
Test1	TEST_BANK	
PASSED	NONE	t1: test deposit_on with correct imp and incomplete contract
FAILED	Check assertion violated.	t2: test deposit_on with wrong imp but incomplete contract

Version 3: Complete Contracts with Reference Copy

```
class BANK
  deposit_on_v3 (n: STRING; a: INTEGER)
    require across accounts as acc some acc.item.owner ~ n end
    local i: INTEGER
    do
      -- same loop as in version 1
      -- wrong implementation: also deposit in the first account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged: accounts.count = old accounts.count
      balance_of_n_increased:
        account_of(n).balance = old account_of(n).balance + a
      others_unchanged :
        across old accounts as cursor
          all cursor.item.owner /~ n implies
            cursor.item ~ account_of (cursor.item.owner)
        end
    end
end
end
```

Test of Version 3

```
class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_ref_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t3: wrong imp and complete contract with ref copy")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v3 ("Steve", 100)
    Result :=
      b.account_of ("Bill").balance = 0
      and b.account_of ("Steve").balance = 100
    check Result end
  end
end
```


Test of Version 3: Result

APPLICATION

Note: * indicates a violation test case

FAILED (2 failed & 1 passed out of 3)		
Case Type	Passed	Total
Violation	0	0
Boolean	1	3
All Cases	1	3
State	Contract Violation	Test Name
Test1	TEST_BANK	
PASSED	NONE	t1: test deposit_on with correct imp and incomplete contract
FAILED	Check assertion violated.	t2: test deposit_on with wrong imp but incomplete contract
FAILED	Check assertion violated.	t3: test deposit_on with wrong imp, complete contract with reference copy

Version 4:

Complete Contracts with Shallow Object Copy

```
class BANK
  deposit_on_v4 (n: STRING; a: INTEGER)
    require across accounts as acc some acc.item.owner ~ n end
    local i: INTEGER
    do
      -- same loop as in version 1
      -- wrong implementation: also deposit in the first account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged: accounts.count = old accounts.count
      balance_of_n_increased:
        account_of (n).balance = old account_of (n).balance + a
        others_unchanged :
          across old accounts.twin as cursor
            all cursor.item.owner /~ n implies
              cursor.item ~ account_of (cursor.item.owner)
    end
  end
end
```

Test of Version 4

```
class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_shallow_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t4: wrong imp and complete contract with shallow copy")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v4 ("Steve", 100)
    Result :=
      b.account_of ("Bill").balance = 0
      and b.account_of ("Steve").balance = 100
    check Result end
  end
end
```

Test of Version 4: Result

APPLICATION

Note: * indicates a violation test case

FAILED (3 failed & 1 passed out of 4)		
Case Type	Passed	Total
Violation	0	0
Boolean	1	4
All Cases	1	4
State	Contract Violation	Test Name
Test1	TEST_BANK	
PASSED	NONE	t1: test deposit_on with correct imp and incomplete contract
FAILED	Check assertion violated.	t2: test deposit_on with wrong imp but incomplete contract
FAILED	Check assertion violated.	t3: test deposit_on with wrong imp, complete contract with reference copy
FAILED	Check assertion violated.	t4: test deposit_on with wrong imp, complete contract with shallow object copy

Version 5:

Complete Contracts with Deep Object Copy

```
class BANK
  deposit_on_v5 (n: STRING; a: INTEGER)
    require across accounts as acc some acc.item.owner ~ n end
    local i: INTEGER
    do
      -- same loop as in version 1
      -- wrong implementation: also deposit in the first account
      accounts[accounts.lower].deposit(a)
    ensure
      num_of_accounts_unchanged: accounts.count = old accounts.count
      balance_of_n_increased:
        account_of (n).balance = old account_of (n).balance + a
        others_unchanged :
          across old accounts.deep_twin as cursor
            all cursor.item.owner /~ n implies
              cursor.item ~ account_of (cursor.item.owner)
          end
    end
  end
end
```

Test of Version 5

```
class TEST_BANK
  test_bank_deposit_wrong_imp_complete_contract_deep_copy: BOOLEAN
  local
    b: BANK
  do
    comment("t5: wrong imp and complete contract with deep copy")
    create b.make
    b.add ("Bill")
    b.add ("Steve")

    -- deposit 100 dollars to Steve's account
    b.deposit_on_v5 ("Steve", 100)
    Result :=
      b.account_of ("Bill").balance = 0
      and b.account_of ("Steve").balance = 100
    check Result end
  end
end
```

Test of Version 5: Result

APPLICATION

Note: * indicates a violation test case

FAILED (4 failed & 1 passed out of 5)		
Case Type	Passed	Total
Violation	0	0
Boolean	1	5
All Cases	1	5
State	Contract Violation	Test Name
Test1	TEST_BANK	
PASSED	NONE	t1: test deposit_on with correct imp and incomplete contract
FAILED	Check assertion violated.	t2: test deposit_on with wrong imp but incomplete contract
FAILED	Check assertion violated.	t3: test deposit_on with wrong imp, complete contract with reference copy
FAILED	Check assertion violated.	t4: test deposit_on with wrong imp, complete contract with shallow object copy
FAILED	Postcondition violated.	t5: test deposit_on with wrong imp, complete contract with deep object copy

Exercise

- Consider the query *account_of* (*n*: *STRING*) of *BANK*.
- How do we specify (part of) its postcondition to assert that the state of the bank remains unchanged:

- `accounts = old accounts` [×]
- `accounts = old accounts.twin` [×]
- `accounts = old accounts.deep_twin` [×]
- `accounts ~ old accounts` [×]
- `accounts ~ old accounts.twin` [×]
- `accounts ~ old accounts.deep_twin` [✓]

- Which equality of the above is appropriate for the postcondition?
- Why is each one of the other equalities not appropriate?

Index (1)

Expanded Class: Modelling

Expanded Class: Programming (2)

Expanded Class: Programming (3)

Reference vs. Expanded (1)

Reference vs. Expanded (2)

Copying Objects

Copying Objects: Reference Copy

Copying Objects: Shallow Copy

Copying Objects: Deep Copy

Example: Copying Objects

Example: Collection Objects (1)

Example: Collection Objects (2)

Reference Copy of Collection Object

Shallow Copy of Collection Object (1)

Index (2)

Shallow Copy of Collection Object (2)

Deep Copy of Collection Object (1)

Deep Copy of Collection Object (2)

How are contracts checked at runtime?

When are contracts complete?

Account

Bank

Roadmap of Illustrations

Object Structure for Illustration

Version 1:

Incomplete Contracts, Correct Implementation

Test of Version 1

Test of Version 1: Result

Version 2:

Incomplete Contracts, Wrong Implementation

Index (3)

Test of Version 2

Test of Version 2: Result

Version 3:

Complete Contracts with Reference Copy

Test of Version 3

Test of Version 3: Result

Version 4:

Complete Contracts with Shallow Object Copy

Test of Version 4

Test of Version 4: Result

Version 5:

Complete Contracts with Deep Object Copy

Test of Version 5

Test of Version 5: Result

Exercise