

**Homework Assignment #9**  
**Due: March 27, 2020 at 2:30 p.m.**

Note: the answers to the questions on this assignment can be very short.

1. Suppose you have a directed graph with  $n$  nodes and  $\Theta(n)$  edges. Assume the graph is given using an adjacency list representation. Each edge has a non-negative edge weight. There is one special vertex  $w$  in the graph. You wish to find, for every pair of nodes  $u$  and  $v$ , the shortest path from  $u$  to  $v$  that passes through  $w$ .
  - (a) Describe an algorithm to solve this problem. Your algorithm should be as efficient as possible. You can use any algorithm discussed in class or the textbook as a subroutine. Describe how any data structures used in your algorithm are implemented.
  - (b) What is the worst-case running time of your algorithm. Express your answer in terms of  $n$  using  $\Theta$  notation.
  
2. Consider an undirected graph with non-negative edge weights. The graph has  $n$  nodes and  $m$  edges and assume it is given using an adjacency list representation. As usual,  $distance(u, v)$  is the length of the shortest path from  $u$  to  $v$ . Let  $D = \max_{u,v} distance(u, v)$  be the maximum distance between any two nodes in the graph.
  - (a) How can you compute  $D$  in worst-case  $\Theta(n^3)$  time?
  - (b) Suppose we just want to get an approximate value of  $D$ . Let  $s$  be any node in the graph. Let  $D' = \max_v distance(s, v)$  be the maximum distance from  $s$  to any other node in the graph. Show that the value of  $D'$  is within a factor of 2 of  $D$  and can be computed much faster than the exact computation in part (a).