

# EECS 1028 M: Discrete Mathematics for Engineers

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Course page: <http://www.eecs.yorku.ca/course/1028>

Also on Moodle

# Why Study Logic?

A formal mathematical “language” for precise reasoning.

- Start with propositions.
- Add other constructs like negation, conjunction, disjunction, implication etc.
- All of these are based on ideas we use daily to reason about things.
- Later: A more expressive language – Predicate logic

# Propositions

- Declarative sentence.
- Must be either True or False.
- Examples of propositions:
  - York University is in Toronto
  - York University is in downtown Toronto
  - All students at York are Computer Science majors
- Examples of statements that are not propositions:
  - Do you like this class?
  - There are  $n$  students in this class.

# Propositions - 2

- Truth value: True or False
- Variables:  $p, q, r, s, \dots$
- Negation:  $\neg p$  (In English, “not  $p$ ”)
- Truth tables – enumerative definition of propositions

| $p$ | $\neg p$ |
|-----|----------|
| T   | F        |
| F   | T        |

# Negating Propositions

$\neg p$ : Literally, “it is not the case that  $p$  is true”

- $p$ : “it rained more than 20 inches in Toronto last month”
- $q$ : “John has many iPads”
- Page 12, Q10 (a)  $r$ : “the election is decided”

Practice: Questions 1-7 page 12.

# Combining Propositions

Purpose: express more complex statements

- Conjunction, Disjunction
- Exclusive OR (XOR)
- Conditionals, Biconditionals
- Logical Equivalence

# Conjunctions and Disjunctions

Purpose: combine statements using OR and AND

- Conjunction (AND):  $p \wedge q$  [ $p$  and  $q$ ]
- Disjunction (OR):  $p \vee q$  [ $p$  or  $q$ ]

| $p$ | $q$ | $p \wedge q$ | $p \vee q$ |
|-----|-----|--------------|------------|
| F   | F   | F            | F          |
| F   | T   | F            | T          |
| T   | F   | F            | T          |
| T   | T   | T            | T          |

# Examples

Q11, page 13

$p$ : It is below freezing

$q$ : It is snowing

- It is below freezing and snowing
- It is below freezing but not snowing
- It is either snowing or below freezing (or both)



# Exclusive OR

Notation:  $p \oplus q$

- TRUE if  $p$  and  $q$  have different truth values, FALSE otherwise
- Colloquially, we often use OR ambiguously –
  - “an entree comes with soup or salad” implies XOR, but
  - “students can take MATH XXXX if they have taken MATH 2320 or MATH 1019” usually means the normal OR (so a student who has taken both is still eligible for MATH XXXX).

# Conditionals

Notation:  $p \rightarrow q$  [“if  $p$  then  $q$ ”]

$p$ : hypothesis,  $q$ : conclusion

Examples:

- “If you turn in a homework late, (then) it will not be graded”
- If you get 100% in this course, (then) you will get an A+”

A conditional is a proposition

- Tricky question: Is  $p \rightarrow q$  TRUE if  $p$  is FALSE?
- Think of “If you get 100% in this course, you will get an A+” as a promise – is the promise violated if someone gets 50% and does not receive an A+?

Q: Similarities with `if(...)` then ... statement in programming?

# Conditionals - Truth Table

$p \rightarrow q$ : When is it False?

Q17, pg 14:

- If  $1 + 1 = 3$  then  $2 + 2 = 4$
- If  $1 + 1 = 3$  then  $2 + 2 = 5$
- If  $1 + 1 = 2$  then  $2 + 2 = 4$
- If  $1 + 1 = 2$  then  $2 + 2 = 5$

|       |
|-------|
| FALSE |
|-------|

| $p$ | $q$ | $p \rightarrow q$ | $\neg p \vee q$ |
|-----|-----|-------------------|-----------------|
| F   | F   | T                 | T               |
| F   | T   | T                 | T               |
| T   | F   | F                 | F               |
| T   | T   | T                 | T               |

# English Statements to Conditionals (pg 6)

$p \rightarrow q$  may be expressed as

- A sufficient condition for  $q$  is  $p$
- $q$  whenever  $p$
- $q$  unless  $\neg p$
- Difficult: A necessary condition for  $p$  is  $q$   
if  $p$  happened,  $q$  must have happened, i.e.,  $p$  cannot happen if we do not have  $q$ .
- $p$  only if  $q$ : not the same as  $p$  if  $q$ ! Same as the previous point, if  $p$  happened,  $q$  must have happened

# Logical Equivalence

$p \rightarrow q$  and  $\neg p \vee q$  have the truth table:  
Does that make them equal? equivalent?

- $p \rightarrow q$  and  $\neg p \vee q$  are **logically** equivalent
- Truth tables are the simplest way to prove such facts.
- We will learn other ways later.

# Biconditionals

Notation:  $p \leftrightarrow q$  [“if and only if”]

- True if  $p, q$  have same truth values, false otherwise.
- Can also be defined as  $(p \rightarrow q) \wedge (q \rightarrow p)$
- Example: Q16(c) “ $1+1=3$  if and only if monkeys can fly”.
- Q: How is this related to XOR?

| $p$ | $q$ | $p \leftrightarrow q$ | $p \oplus q$ |
|-----|-----|-----------------------|--------------|
| F   | F   | T                     | F            |
| F   | T   | F                     | T            |
| T   | F   | F                     | T            |
| T   | T   | T                     | F            |

# Contrapositive

Contrapositive of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$

- E.g. The contrapositive of “If you get 100% in this course, you will get an A+” is “If you do not get an A+ in this course, you did not get 100%”.
- Any conditional and its contrapositive are logically equivalent (have the same truth table).

| $p$ | $q$ | $p \rightarrow q$ | $\neg q$ | $\neg p$ | $\neg q \rightarrow \neg p$ |
|-----|-----|-------------------|----------|----------|-----------------------------|
| F   | F   | T                 | T        | T        | T                           |
| F   | T   | T                 | F        | T        | T                           |
| T   | F   | F                 | T        | F        | F                           |
| T   | T   | T                 | F        | F        | T                           |

# Proof using Contrapositive

Prove: If  $x^2$  is even, then  $x$  is even

- Proof 1: Using contradiction, seen before.
- Proof 2:  
 $x^2 = 2a$  for some integer  $a$ . Since 2 is prime, 2 must divide  $x$ .  
(Uses knowledge of primes)
- Proof 3:  
if  $x$  is not even, then  $x$  is odd. Therefore  $x^2$  is odd. This is the contrapositive of the original assertion.  
(Uses only facts about odd and even numbers)



# Converse and Inverse

Converse of  $p \rightarrow q$  is  $q \rightarrow p$

Converse of  $p \rightarrow q$  is  $\neg p \rightarrow \neg q$

- Converse examples:
  - “If you get 100% in this course, you will get an A+”, converse “If you get an A+ in this course, you scored 100%”.
  - “If you won the lottery, you are rich”, converse “If you are rich, you (must have) won the lottery”.
- Neither is logically equivalent to the original conditional

| $p$ | $q$ | $p \rightarrow q$ | $q \rightarrow p$ | $\neg p \rightarrow \neg q$ |
|-----|-----|-------------------|-------------------|-----------------------------|
| F   | F   | T                 | T                 | T                           |
| F   | T   | T                 | F                 | F                           |
| T   | F   | F                 | T                 | T                           |
| T   | T   | T                 | T                 | T                           |

# Tautology and Logical Equivalence

Tautology: A (compound) proposition that is always TRUE,  
e.g.  $q \vee \neg q$

- Logical equivalence redefined:  $p, q$  are logical equivalences (Symbolically  $p \equiv q$ ) if  $p \leftrightarrow q$  is a tautology. .
- Intuition:  $p \leftrightarrow q$  is true precisely when  $p, q$  have the same truth values.

# Compound Propositions: Precedence

Example:  $p \wedge q \vee r$ : Could be interpreted as  $(p \wedge q) \vee r$  or  $p \wedge (q \vee r)$

- precedence order:  $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$   
(Overruled by brackets)
- We use this order to compute truth values of compound propositions.

# Translating English Sentences to Propositional Logic statements

Pages 14-15:

- I will remember to send you the address only if you send me an email message
- The beach erodes whenever there is a storm
- John will go swimming unless the water is too cold
- Getting elected follows from knowing the right people.

# Readings and Notes

- Read pages 1-12.
- Think about the notion of truth tables.
- Master the rationale behind the definition of conditionals.
- Practice translating English sentences to propositional logic statements.

# Manipulating Propositions (Sec 1.3)

- Compound propositions can be simplified by using simple rules. Read page 25 - 28.
- Some are obvious, e.g. Identity, Domination, Idempotence, Negation, Double negation, Commutativity, Associativity
- Less obvious: Distributive, De Morgan's laws, Absorption

TABLE 6 Logical Equivalences.

| <i>Equivalence</i>                                                                                                       | <i>Name</i>         |
|--------------------------------------------------------------------------------------------------------------------------|---------------------|
| $p \wedge \mathbf{T} \equiv p$<br>$p \vee \mathbf{F} \equiv p$                                                           | Identity laws       |
| $p \vee \mathbf{T} \equiv \mathbf{T}$<br>$p \wedge \mathbf{F} \equiv \mathbf{F}$                                         | Domination laws     |
| $p \vee p \equiv p$<br>$p \wedge p \equiv p$                                                                             | Idempotent laws     |
| $\neg(\neg p) \equiv p$                                                                                                  | Double negation law |
| $p \vee q \equiv q \vee p$<br>$p \wedge q \equiv q \wedge p$                                                             | Commutative laws    |
| $(p \vee q) \vee r \equiv p \vee (q \vee r)$<br>$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$                     | Associative laws    |
| $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$<br>$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$ | Distributive laws   |
| $\neg(p \wedge q) \equiv \neg p \vee \neg q$<br>$\neg(p \vee q) \equiv \neg p \wedge \neg q$                             | De Morgan's laws    |
| $p \vee (p \wedge q) \equiv p$<br>$p \wedge (p \vee q) \equiv p$                                                         | Absorption laws     |
| $p \vee \neg p \equiv \mathbf{T}$<br>$p \wedge \neg p \equiv \mathbf{F}$                                                 | Negation laws       |

# Distributive Laws

- $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

Intuition (not a proof!) - For the LHS to be true:  $p$  must be true and  $q$  or  $r$  must be true. This is the same as saying  $p$  and  $q$  must be true or  $p$  and  $r$  must be true.

- $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$

Intuition (less obvious) - For the LHS to be true:  $p$  must be true or both  $q$  and  $r$  must be true. This is the same as saying  $p$  or  $q$  must be true and  $p$  or  $r$  must be true.

Proof: use truth tables.

# De Morgan's Laws

- $\neg(q \vee r) \equiv \neg q \wedge \neg r$

Intuition - For the LHS to be true: neither  $q$  nor  $r$  can be true. This is the same as saying  $q$  and  $r$  must both be false.

- $\neg(q \wedge r) \equiv \neg q \vee \neg r$

Intuition - For the LHS to be true:  $q \wedge r$  must be false. This is the same as saying that  $q$  or  $r$  must be false.

Proof: use truth tables.



# Negating Conditionals

The negation of  $p \rightarrow q$  is NOT  $\neg p \rightarrow \neg q$  or any other conditional

- Easiest to negate the logically equivalent form of  $p \rightarrow q$ , viz.,  $\neg p \vee q$ .

$$\text{So } \neg(p \rightarrow q) \equiv \neg(\neg p \vee q) \equiv p \wedge \neg q$$

- Relate to the truth table of  $p \rightarrow q$

| $p$ | $q$ | $p \rightarrow q$ | $\neg(p \rightarrow q)$ | $p \wedge \neg q$ |
|-----|-----|-------------------|-------------------------|-------------------|
| F   | F   | T                 | F                       | F                 |
| F   | T   | T                 | F                       | F                 |
| T   | F   | F                 | T                       | T                 |
| T   | T   | T                 | F                       | F                 |

# Using the laws

Q: Is  $p \rightarrow (p \rightarrow q)$  a tautology?

- Can use truth tables

| $p$ | $q$ | $p \rightarrow q$ | $p \rightarrow (p \rightarrow q)$ |
|-----|-----|-------------------|-----------------------------------|
| F   | F   | T                 | T                                 |
| F   | T   | T                 | T                                 |
| T   | F   | F                 | F                                 |
| T   | T   | T                 | T                                 |

- Can write a compound proposition and simplify:

$$\begin{aligned}
 p \rightarrow (p \rightarrow q) &\equiv \neg p \vee (\neg p \vee q) \\
 &\equiv \neg p \vee \neg p \vee q \\
 &\equiv \neg p \vee q
 \end{aligned}$$

This is False when  $p$  is True and  $q$  is False