## Implementing Recursion

## Printing n of Something

- Suppose you want to implement a method that prints out $n$ copies of a string

```
public static void printIt(String s, int n)
{
    for(int i = 0; i < n; i++)
    {
        System.out.print(s);
    }
}
```


## A Different Solution

- Alternatively we can use the following algorithm:

1. if $\mathrm{n}==0$ done, otherwise
I. print the string once
II. print the string ( $\mathrm{n}-1$ ) more times
public static void printItToo(String s, int $n$ )
\{
if ( $\mathrm{n}=\mathbf{0}$ )
\{
return;
\}
else
\{
System.out.print(s);
printItToo(s, n - 1); // method invokes itself
\}

## Recursion

- A method that calls itself is called a recursive method
- A recursive method solves a problem by repeatedly reducing the problem so that a base case can be reached

```
printIt("*", 5)
*printIt("*", 4)
**printIt("*", 3)
***printIt("*", 2)
****printIt("*", 1)
*****printIt("*", 0) base case
* * * * *
```

Notice that the number of times the string is printed decreases
after each recursive call to printlt

Notice that the base case is eventually reached.

## Infinite Recursion

- If the base case(s) is missing, or never reached, a recursive method will run forever (or until the computer runs out of resources)

```
public static void printItForever(String s, int n)
{
    // missing base case; infinite recursion
    System.out.print(s);
    printItForever(s, n - 1);
}
printIt("*", 1)
* printIt("*", 0)
** printIt("*", -1)
*** printIt("*", -2)
```


## Fibonacci Numbers

The sequence of additional pairs
-0, 1, 1, 2, 3, 5, 8, 13, ... are called Fibonacci numbers

- Base cases
- $F(0)=0$
$\circ F(1)=1$
- Recursive definition
$\circ F(n)=F(n-1)+F(n-2)$


## Recursive Methods \& Return

## Values

- A recursive method can return a value
- Example: compute the nth Fibonacci number

```
public static int fibonacci(int n)
{
    if (n == 0)
    {
        return 0;
    }
    else if (n == 1)
    {
        return 1;
    }
    else
    {
        int f = fibonacci(n - 1) + fibonacci(n - 2);
        return f;
    }
```


## Recursive Methods \& Return

## Values

- Example: write a recursive method countZeros that counts the number of zeros in an integer number $\mathbf{n}$
- 10305060700002L has 8 zeros
- Trick: examine the following sequence of numbers

1. 10305060700002
2. 1030506070000
3. 103050607000
4. 10305060700
5. 103050607
6. 1030506

## Recursive Methods \& Return Values

- Not Java:

```
countZeros(n) :
if the last digit in n is a zero
    return 1 + countZeros(n / 10)
else
    return countZeros(n / 10)
```

- Don't forget to establish the base case(s)
- When should the recursion stop? when you reach a single digit (not zero digits; you never reach zero digits!)
- Base case \#1 : n == 0
- return 1
- Base case \#2 : n != 0 \&\& n < 10
- return 0

```
public static int countZeros(long n)
{
    if(n == OL)
    { // base case 1
        return 1;
    }
    else if(n < 10L)
    { // base case 2
        return 0;
    }
    boolean lastDigitlsZero = (n % 10L == 0);
    final long m = n / 10L;
    if(lastDigitlsZero)
    {
        return 1 + countZeros(m);
    }
    else
    {
        return countZeros(m);
}
```


## countZeros Call Stack

## callZeros( 800410L )

last in first out

| callZeros( 8L ) | 0 |
| :---: | :---: |
| callZeros( 80L ) | $1+0$ |
| callZeros( 800L ) | $1+1+0$ |
| callZeros( 8004L ) | $0+1+1+0$ |
| callZeros( 80041L ) | $0+0+1+1+0$ |
| callZeros( 800410L ) | $1+0+0+1+1+0$ |

## Fibonacci Call Tree



## Compute Powers of 10

- Write a recursive method that computes $\mathbf{1 0}^{\text {n }}$ for any integer value $\mathbf{n}$
- Recall:
- $10^{0}=1$
- $10^{n}=10 * 10^{n-1}$
- $10^{-n}=1 / 10^{n}$

```
public static double powerOf10(int n)
    {
    if (n == 0)
    {
        // base case
        return 1.0;
    }
    else if (n > 0)
    {
        // recursive call for positive n
        return 10.0* powerOf10(n-1);
    }
    else
    {
        // recursive call for negative n
        return 1.0 / powerOf10(-n);
    }
```



## Proving Correctness and Termination

- To show that a recursive method accomplishes its goal you must prove:

1. That the base case(s) and the recursive calls are correct
2. That the method terminates

## Proving Correctness

- To prove correctness:

1. Prove that each base case is correct
2. Assume that the recursive invocation is correct and then prove that each recursive case is correct

## printltToo

public static void printltToo(String s, int $n$ )

$$
\text { if }(\mathrm{n}==0)
$$

$$
\{
$$

return;
\}
else \{
System.out.print(s); printltToo(s, n-1);

## Correctness of printltToo

1. (prove the base case) If $\mathbf{n}==\mathbf{0}$ nothing is printed; thus the base case is correct.
2. Assume that printitToo(s, n-1) prints the string s exactly(n-1) times. Then the recursive case prints the string s exactly( $n$ 1)+1 = $n$ times; thus the recursive case is correct.

## Proving Termination

To prove that a recursive method terminates:

1. Define the size of a method invocation; the size must be a non-negative integer number
2. Prove that each recursive invocation has a smaller size than the original invocation

## Termination of printlt

1. printIt( $\mathbf{s}, \mathbf{n})$ prints $\mathbf{n}$ copies of the string $\mathbf{s}$; define the size of printIt $(\mathbf{s}, \mathrm{n})$ to be $\mathbf{n}$
2. The size of the recursive invocation printIt( $s, n-1$ ) is $\mathbf{n - 1}$ (by definition) which is smaller than the original size $n$.

## countZeros

```
public static int countZeros(long n)
{
    if(n == 0L)
    { // base case 1
        return 1;
    }
    else if(n < 10L)
    { // base case 2
        return 0;
    }
    boolean lastDigitlsZero = (n % 10L == 0);
    final long m = n / 10L;
    if(lastDigitlsZero)
    {
        return 1 + countZeros(m);
    }
    else
    {
        return countZeros(m);
}
```


## Correctness of countZeros

1. (Base cases) If the number has only one digit then the method returns 1 if the digit is zero and 0 if the digit is not zero; therefore, the base case is correct.
2. (Recursive cases) Assume that countZeros( $\mathrm{n} / \mathbf{1 0 L}$ ) is correct (it returns the number of zeros in the first ( $\mathbf{d} \mathbf{- 1}$ ) digits of $\mathbf{n}$ ). If the last digit in the number is zero, then the recursive case returns $1+$ the number of zeros in the first ( $\mathbf{d}-\mathbf{1}$ ) digits of $n$, otherwise it returns the number of zeros in the first ( d 1) digits of $n$; therefore, the recursive cases are correct.

## Termination of countZeros

1. Let the size of countZeros( $\mathbf{n}$ ) be $\mathbf{d}$ the number of digits in the number $n$.
2. The size of the recursive invocation countZeros( $\mathrm{n} / 10 \mathrm{~L}$ ) is $\mathbf{d - 1}$, which is smaller than the size of the original invocation.

## Decrease and Conquer

- A common strategy for solving computational problems
- Solves a problem by taking the original problem and converting it to one smaller version of the same problem
- Note the similarity to recursion
- Decrease and conquer, and the closely related divide and conquer method, are widely used in computer science
- Allow you to solve certain complex problems easily
- Help to discover efficient algorithms


## Review of Recursion

- A recursive method calls itself
- To prevent infinite recursion you need to ensure that:

1. The method reaches a base case
2. Each recursive call makes progress towards a base case (i.e. reduces the size of the problem)
To solve a problem with a recursive algorithm:
3. Identify the base cases (the cases corresponding to the smallest version of the problem you are trying to solve)
4. Figure out the recursive call(s)

## Correctness and Termination

- Proving correctness requires that you do two things:

1. Prove that each base case is correct
2. Assume that the recursive invocation is correct and then prove that each recursive case is correct

- Proving termination requires that you do two things:

1. Define the size of each method invocation
2. Prove that each recursive invocation is smaller than the original invocation

## Recursion Examples

The subsequent slides present additional examples of problems that can be solved using recursion

- Depending on time, these examples may or may not be discussed in lecture.


## Palindromes

A palindrome is a sequence of symbols that is the same forwards and backwards: "level"
"yo banana boy"
Write a recursive algorithm that returns true if a string is a palindrome (and false if not); assume that the string has no spaces or punctuation marks.

## Palindromes

- Sketch a small example of the problem
- It will help you find the base cases
- It might help you find the recursive cases


## Palindromes

```
public static boolean isPalindrome(String s)
{
    if (s.length() < 2)
    {
    return true;
    }
    else
    {
        int first = 0;
        int last = s.length() - 1;
        return (s.charAt(first) == s.charAt(last)) &&
        isPalindrome(s.substring(first + 1, last));
    }
```


## Towers of Hanoi

3. $[A J, p$ 685, Q7]


$$
\mathrm{C}
$$

- Move the stack of $n$ disks from A to C
- Can move one disk at a time from the top of one stack onto another stack
- Cannot move a larger disk onto a smaller disk


## Towers of Hanoi

- Legend says that the world will end when a 64 disk version of the puzzle is solved
- Several appearances in pop culture
- Doctor Who (TV series)
- Rise of the Planet of the Apes (Movie)
- Mass Effect (Video game)


## Towers of Hanoi

, $n=1$

, Move disk from A to C

## Towers of Hanoi

- $n=1$



## Towers of Hanoi

- $n=2$

- Move disk from A to B


## Towers of Hanoi

- $n=2$

, Move disk from A to C


## Towers of Hanoi

- $n=2$


Move disk from B to C

## Towers of Hanoi

- $n=2$



## Towers of Hanoi

- $n=3$

- Move disk from A to C


## Towers of Hanoi

- $n=3$

- Move disk from A to B


## Towers of Hanoi

- $n=3$

- Move disk from C to B


## Towers of Hanoi

- $n=3$

- Move disk from A to C


## Towers of Hanoi

- $n=3$

- Move disk from B to A


## Towers of Hanoi

- $n=3$

- Move disk from B to C


## Towers of Hanoi

- $n=3$

- Move disk from A to C


## Towers of Hanoi

, $n=3$


## Towers of Hanoi

- $n=4$

- Move ( $\mathrm{n}-1$ ) disks from A to B using C


## Towers of Hanoi

- $n=4$

, Move disk from A to C


## Towers of Hanoi

- $n=4$

- Move ( $\mathrm{n}-1$ ) disks from B to C using A


## Towers of Hanoi

- $n=4$

- Base case $n=1$

1. Move disk from A to C

- Recursive case

1. Move ( $n-1$ ) disks from A to B
2. Move 1 disk from A to C
3. Move $(n-1)$ disks from B to $C$

## Towers of Hanoi

```
public static void move(int n,
                                    String from,
                                    String to,
                                    String using)
{
    if(n == 1)
    {
        System.out.println("move disk from " + from + " to " + to);
    }
    else
    {
        move(n - 1, from, using, to);
        move(1, from, to, using);
        move(n - 1, using, to, from);
    }
```

