Best-First Search Minimizing Space or Time

RBFS Save space, take more time

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 - » A* keeps in memory all of the already generated nodes
 - » RBFS only keeps the current search path and the sibling nodes along the path

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- What is the space complexity?
 Linear the depth of the search
 Same as IDA*

RBFS memory

» When RBFS suspends searching a subtree, what does it remember?

RBFS memory – 2

» When RBFS suspends searching a subtree, what does it remember?

> An updated f-value of the root of the subtree

Updated f-values

» How does RBFS update the f-values?

Updated f-values – 2

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> Backing up the f-values in the same way as A* does

f-value notation

- ♦ Static f-value
 - » f(N)
 - > Value returned by the evaluation function
 - > Always the same

f-value notation – 2

- ♦ Static f-value
 - » f(N)
 - > Value returned by the evaluation function
 - > Always the same
- Backed-up value
 - » F(N)
 - > Changes during the search
 - Depends upon descendants of N

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$$- \mathbf{F}(\mathbf{N}) = \min \left(\mathbf{F} \left(\mathbf{S}_{\mathbf{j}} \right) \right)$$

- Where S_j are the subtrees of N

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RBFS subtree exploration – 4

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- » How does RBFS explore subtrees?
 - > As in A*, within a given f-bound
- **»** How is the bound determined?
 - > From the F-values of the siblings along the current search path
 - > The smallest F-value
 - The closest competitor

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 - > N is expanded
 - > N's children are expanded
 - » Until when?
 - > F(N) > Bound
 - » Then what happens?
 - > Nodes below N are forgotten
 - > N's F-value is updated
 - > RBFS selects which node to expand next

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 - > N_k's F-value can be inherited from N
 - N_k was generated earlier
 - $F(N_k)$ was $\geq F(N)$, otherwise F(N) would be smaller



f(n) in mocha = g(n) in clover + h(n) in magenta

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f(n) in mocha = g(n) in clover + h(n) in magenta

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RBFS-51

Fig 12.2 snapshots – 3





Forget expansion from A

A has backed up F value 10

E is best to expand next

f(n) in mocha = g(n) in clover + h(n) in magenta

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RBFS-52



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Fig 12.2 snapshots – 5





Forget expansion from E

E has backed up F value 11

A is best to expand next

f(n) in mocha = g(n) in clover + h(n) in magenta

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RBFS-54



f(n) in mocha = g(n) in clover + h(n) in magenta

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Fig 12.2 snapshots – 7



F = 12 (A) (E) F = 11

Forget expansion from A

A has backed up F value 12

E is best to expand next

f(n) in mocha = g(n) in clover + h(n) in magenta

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RBFS-56



f(n) in mocha = g(n) in clover + h(n) in magenta

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Algorithm

function NewF (N, F(N), Bound) if F(N) > Bound then NewF := F(N)else if goal(N) then exit search with success else if N has no children then NewF := infinity – dead end else for each child N_k of N do if f(N) < F(N) then $F(N_k) := max(F(N), f(N_k))$ else $F(N_k) := f(N_k)$ sort children N_k in increasing order of F-value while $F(N_1) \leq Bound$ and $F(N_1) < infinity$ do Bound1 := min (Bound, F-value of sibling N_1) $F(N_1) := NewF(N_1, F(N_1), Bound1)$ reorder nodes N_1, N_2, \dots according to new $F(N_1)$ end end NewF := $F(N_1)$

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