## Example predicates

## Showing things to look for

## Infinite loops

$\diamond$ Avoid circular definitions

$$
\begin{aligned}
& \text { parent (A, B ) :- child (B, A ). } \\
& \text { child ( C, D ) :- } \operatorname{parent~(D,~C~).~}
\end{aligned}
$$

$\diamond$ Easy to see here but as database grows you can forget what is in it and circularity can creep in

## Infinite loops - Left Recursion - 1

$\diamond$ Left recursion can cause problems person ( X ) :- person (Y), mother ( Y, X ). person (eve).
» The query person ( P ) loops indefinitely as the first rule is found first on every recursive call.
»Second rule is only tried if first rule fails
$\diamond$ Reordering the rules will correct the problem

Heuristic<br>Put facts before rules

## Infinite loops - Left Recursion - 2

$\diamond$ Left recursion can cause problems - continued person (eve).
person ( X ) :- person (Y), mother ( Y, X ).
» Assuming mother fails, the query person ( P ) loops indefinitely after $P=$ eve
$\diamond$ Left recursion is the problem

Do not assume Prolog will find the facts and rules. Need to know how searching works

## Multiple answers - isList, weakList

$\diamond$ The following gives the following predicate works for specific lists but loops forever on the query isList ( X ).
isList ([AIB]):- isList (B). isList ([]).
$\diamond$ It can be defined by putting the fact first. isList ([]).
isList ([AIB]) :- isList (B).
$\diamond$ But gives more than one answer for the query isList ( X ) but does not loop forever.
$\diamond$ For the latter query, to have only one answer, can assert the following.
> weak_isList ([]). weak_isList ([_I_]).

## Why is weak_isList weak?

$\diamond$ The strong definition says a list must have the correct structure and must end in nil.
$\diamond$ The weak definition simply says the list must have the correct structure for one level and says nothing about nil except for the empty list.
$\diamond$ For example - recall [...] is shorthand for the structure .(...)

```
isList (.(a, [] ) ).
isList (.(a,.(b,[])) ). ==> yes
isList (.(a,.(b,.(c c,[])))). ==> yes
isList (.(a,b)). ==> no
isList (.(a,.(b, c, [] ))). ==> no
```

$\diamond$ But all responses are yes for weak_isList

## Mapping

$\diamond$ Consider the problem of translating a sentence from one form to another
" For example as in the following "dialogue" the second sentence is a translation of the preceding sentence
> you are a computer
I am not a computer
> do you speak french no l speak german
» Assume the following simplistic translations

```
\(>\) you \(==>\) I
    are \(==>\) am not
    do \(==>\) no
    french ==> german
```


## Mapping - 2

$\diamond$ Let us represent sentences as a list of words you are a computer ==> [you, are, a , computer]
$\diamond$ We represent the list of words to change as a set of change rules
change (you, l).
change ( are, [ am , not ]).
change ( french, german ).
change ( do , no ).
change ( $\mathrm{X}, \mathrm{X}$ ). /* catch all to make no changes */

## Mapping - 3

$\diamond$ Then the translation rules can be the following.

```
alter ([],[]).
alter ([ H I T ] , [ X I Y ] ) :- change ( H, X ), alter ( T, Y ).
```

$\diamond$ Then we can translate our example sentences alter ([ you, are, a, computer ], Trans ).
$>$ Trans = [I, am , not , a , computer ]
> Try using ;<return> on the above. Explain why there are multiple answers. Try a trace to see what is happening.
> We need a method to prevent multiple answers

## Mapping - 4

$\diamond$ Try the inverse - with ;<return> alter ( Org , [ I , am , not , a , computer ] ).
$\diamond$ Try a variable - with ; <return> alter ([ you , are , a , X ] , Trans )

## Warning - Caution - Danger

Logic and a finite database can lead to strange<br>and unexpected results.<br>Use with extreme caution.

