# **Negation of Quantifiers**

- "There is no student who can ..."
- "Not all professors are bad...."
- "There is no Toronto Raptor that can dunk like Vince ..."
- $\neg \forall x P(x) \equiv \exists x \neg P(x) \text{ why?}$
- $\neg \exists x P(x) \equiv \forall x \neg P(x)$
- Careful: The negation of "Every Canadian loves Hockey" is NOT "No Canadian loves Hockey"! <u>Many, many students make this mistake!</u>

## **Nested Quantifiers**

- Allows simultaneous quantification of many variables.
- E.g. domain integers,
  - $\exists x \exists y \exists z x^2 + y^2 = z^2$  (Pythagorean triples)

 $- \forall n \exists x \exists y \exists z x^n + y^n = z^n$  (Fermat's Last Theorem implies this is false)

- Domain real numbers:
  - $\forall x \forall y \exists z (x < z < y) \lor (y < z < x)$  Is this true?
  - $\forall x \forall y \exists z (x=y) \lor (x < z < y) \lor (y < z < x)$

 $- \forall x \forall y \exists z (x \neq y) \rightarrow (x < z < y) \lor (y < z < x)$ 

# **Nested Quantifiers - 2**

 $\forall x \exists y (x + y = 0)$  is true over the integers

- Assume an arbitrary integer x.
- To show that there exists a y that satisfies the requirement of the predicate, choose y = -x. Clearly y is an integer, and thus is in the domain.
- So x + y = x + (-x) = x x = 0.
- Since we assumed nothing about x (other than it is an integer), the argument holds for any integer x.
- Therefore, the predicate is TRUE.

#### **Nested Quantifiers - 3**

- Analogy: quantifiers are like loops:
  - An inner quantified variable can depend on the outer quantified variable.
  - E.g. in  $\forall x \exists y (x + y = 0)$  we chose y=-x, so for different x we need different y to satisfy the statement.
  - ∀p ∃j Accept (p,j) p,j have different domains does NOT say that there is a j that will accept all p.

#### **Nested Quantifiers - 4**

• Caution: In general, order matters! Consider the following propositions over the integer domain:

 $\forall x \exists y (x < y) \text{ and } \exists y \forall x (x < y)$ 

- ∀x ∃y (x < y) : "there is no maximum integer"</li>
- ∃y ∀x (x < y) : "there is a maximum integer"</li>
- Not the same meaning at all!!!

## **Negation of Nested Quantifiers**

- Use the same rule as before carefully.
- Ex 1:  $\neg \exists x \forall y (x + y = 0)$ 
  - This is equivalent to  $\forall x \neg \forall y (x + y = 0)$
  - This is equivalent to  $\forall x \exists y \neg (x + y = 0)$
  - This is equivalent to  $\forall x \exists y (x + y \neq 0)$
- Ex 2:¬ ∀x ∃y (x < y)
  - This is equivalent to  $\exists x \neg \exists y (x < y)$
  - This is equivalent to  $\exists x \forall y \neg (x < y)$
  - This is equivalent to  $\exists x \forall y \ (x \ge y)$

### Exercises

#### Check that:

- $\forall x \exists y (x + y = 0)$  is not true over the positive integers.
- $\exists x \forall y (x + y = 0)$  is not true over the integers.
- ∀x <>0 ∃y (y = 1/x) is true over the real numbers.

## **Readings and notes**

- Read 1.4-1.5.
- Practice: Q2,8,16,30 (pg 65-67)

• Next: Rules of inference for quantified statements (1.6).

#### Inference rules for quantified statements

- Very intuitive, e.g. Universal instantiation –
  If ∀x P(x) is true, we infer that P(a) is true for any given a
- E.g.: Universal Modus Ponens:

 $\forall x P(x) \rightarrow Q(x) \text{ and } P(a) \text{ imply } Q(a)$ 

If x is odd then x<sup>2</sup> is odd, a is odd. So a<sup>2</sup> is odd.

- Read rules on page 76
- Again, understanding is required, memorization is not.

#### Commonly used technique: Universal generalization

Prove: If x is even, x+2 is even

• Proof:

Prove: If x<sup>2</sup> is even, x is even [Note that the problem is to prove an implication.]

 Proof: if x is not even, x is odd. Therefore x<sup>2</sup> is odd. This is the contrapositive of the original assertion.

## **Aside: Inference and Planning**

- The steps in an inference are useful for planning an action.
- Example: your professor has assigned reading from an out-of-print book. How do you do it?
- Example 2: you are participating in the television show "Amazing race". How do you play?

# Aside 2: Inference and Automatic Theorem-Proving

- The steps in an inference are useful for proving assertions from axioms and facts.
- Why is it important for computers to prove theorems?
  - Proving program-correctness
  - Hardware design
  - Data mining

. . . . .

### Aside 3: Inference and Automatic Theorem-Proving

- Sometimes the steps of an inference (proof) are useful. E.g. on Amazon book recommendations are made.
- You can ask why they recommended a certain book to you (reasoning).

## Next

- Introduction to Proofs (Sec 1.7)
- What is a (valid) proof?
- Why are proofs necessary?

### **Introduction to Proof techniques**

Why are proofs necessary?

What is a (valid) proof?

What details do you include/skip? "Obviously", "clearly"...

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