

Test 1**First Name:** _____**Last Name:** _____**Student Number:** _____

This test lasts 80 minutes. No aids allowed.

You may use any result that was proved in class or in the textbook without reproving it.

Make sure your test has 5 pages, including this cover page.

*Answer in the space provided. (If you need more space, use the reverse side of the page and indicate **clearly** which part of your work should be marked.)*

Write legibly.

Question 1	/4
Question 2	/4
Question 3	/3
Question 4	/3
Question 5	/3
Question 6	/3
Total	/20

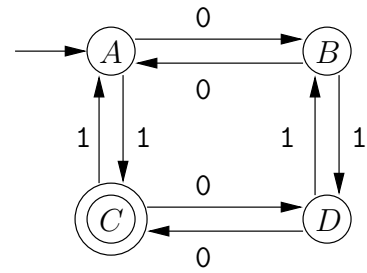
1. [4 marks] Let $L_1 = \{x \in \{0, 1\}^* : x \text{ contains a } 1 \text{ and the length of } x \text{ is a multiple of } 3\}$. Draw the transition diagram of a deterministic finite automaton for L_1 . For each state of your machine, describe the set of strings that take the machine to that state.

2. [4 marks] Consider the DFA M shown at right. Recall that $\delta^*(A, x)$ denotes the state that M is in after processing the string x , starting from the initial state A . Give a careful proof of the following claim.

Claim: For any string $x \in \{0, 1\}^*$,

$\delta^*(A, x) \in \{C, D\}$ if x contains an odd number of 1's, and

$\delta^*(A, x) \in \{A, B\}$ if x contains an even number of 1's.



3. [3 marks] Let $L_3 = \{x \in \{0, 1\}^* : x \text{ starts with } 0 \text{ and has even length}\}$. Write a regular expression for L_3 .
4. [3 marks] Let $L_4 = \{x \in \{0, 1\}^* : x \text{ contains exactly twice as many } 0\text{'s as } 1\text{'s}\}$. For example, 001010 is in L_4 but 11000 is not in L_4 . Prove that L_4 is not regular.

5. [3 marks] Give a high-level description of how a non-deterministic finite automaton N can be transformed into a deterministic finite automaton M that accepts the same language. Your answer must fit inside the box below. Anything written outside the box will be ignored.

6. [3 marks] If $L \subseteq \Sigma^*$ is a language, define $PREFIX(L)$ to be the set of all strings that are prefixes of strings in L . More formally,

$$PREFIX(L) = \{x \in \Sigma^* : \text{there exists } y \in \Sigma^* \text{ such that } xy \in L\}.$$

Show that for every regular language L , the language $PREFIX(L)$ is also regular.