

**Test 1****First Name:** \_\_\_\_\_**Last Name:** \_\_\_\_\_**Student Number:** \_\_\_\_\_

*This test lasts 80 minutes. No aids allowed.*

*You may use any result that was proved in class or in the textbook without reproving it.*

*Make sure your test has 5 pages, including this cover page.*

*Answer in the space provided. (If you need more space, use the reverse side of the page and indicate **clearly** which part of your work should be marked.)*

*Write legibly.*

Question 1	/4
Question 2	/4
Question 3	/3
Question 4	/3
Question 5	/3
Question 6	/3
Total	/20

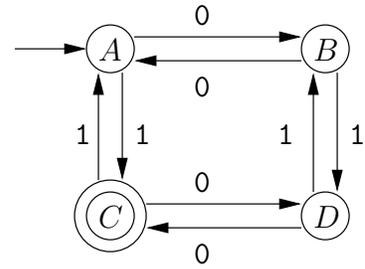
1. [4 marks] Let  $L_1 = \{x \in \{0, 1\}^* : x \text{ contains a } 1 \text{ and the length of } x \text{ is a multiple of } 3\}$ . Draw the transition diagram of a deterministic finite automaton for  $L_1$ . For each state of your machine, describe the set of strings that take the machine to that state.

2. [4 marks] Consider the DFA  $M$  shown at right. Recall that  $\delta^*(A, x)$  denotes the state that  $M$  is in after processing the string  $x$ , starting from the initial state  $A$ . Give a careful proof of the following claim.

**Claim:** For any string  $x \in \{0, 1\}^*$ ,

$\delta^*(A, x) \in \{C, D\}$  if  $x$  contains an odd number of 1's, and

$\delta^*(A, x) \in \{A, B\}$  if  $x$  contains an even number of 1's.



3. [3 marks] Let  $L_3 = \{x \in \{0, 1\}^* : x \text{ starts with } 0 \text{ and has even length}\}$ . Write a regular expression for  $L_3$ .
4. [3 marks] Let  $L_4 = \{x \in \{0, 1\}^* : x \text{ contains exactly twice as many } 0\text{'s as } 1\text{'s}\}$ . For example, 001010 is in  $L_4$  but 11000 is not in  $L_4$ . Prove that  $L_4$  is not regular.

5. [3 marks] Give a high-level description of how a non-deterministic finite automaton  $N$  can be transformed into a deterministic finite automaton  $M$  that accepts the same language. Your answer must fit inside the box below. Anything written outside the box will be ignored.

6. [3 marks] If  $L \subseteq \Sigma^*$  is a language, define  $PREFIX(L)$  to be the set of all strings that are prefixes of strings in  $L$ . More formally,

$$PREFIX(L) = \{x \in \Sigma^* : \text{there exists } y \in \Sigma^* \text{ such that } xy \in L\}.$$

Show that for every regular language  $L$ , the language  $PREFIX(L)$  is also regular.