#### Chat

#### A program to make Prolog input English like

A project from Clocksin and Mellish,

page 244 third edition

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#### The main predicate- chat

- The rule repeats itself until the user enters exactly "Stop."
   chat :- repeat
  - > Get a sentence from the user
    - , readLine (Sentence)
  - > Obtain the semantic form, Clause, from the external form, Sentence.
    - , parse ( Clause , Sentence , \_ )
  - > Determine the appropriate response.
    - , respondTo ( Clause )
  - > chat succeeds when the semantic form is stop
    - , Clause = stop .

## readLine (Sentence)

- Read a sentence as a list of words, where each word is the list of characters in ASCII numeric code.
- Split off the periods, question marks and apostrophes
- Create the corresponding list of atoms readLine (Sentence) :- readCharLists (Words) , morphs (Words, Sentence), !.
- Output States User types John is a person.
- Words ==> [ [74, 111, 104, 110 ], [105, 115 ], [97 ], [112, 101, 114, 115, 111, 110, 46 ] ]
- Sentence ==> [ John, is, a, person, .] > John is a constant not a variable

## readCharLists (Words)

Read in a list of words from the keyboard and convert each word to a list of character lists

```
readCharLists ([Word | MoreWords]) :-
```

> Read a word

readWord ( Word , TerminatingChar )

> end of line (ASCII 10 is newLlne) signals the end of the list of words

, ((TerminatingChar = 10), MoreWords = []

; readCharLists (MoreWords)).

MoreWords is a hole

> see parts assembly example

## readWord( Word, CharList )

- Read in a word from the keyboard readWord (Word, TerminatingChar) :- get0 (C)
  - > Check for end of line or space character , ( ( C = 10 ; C = 32 )
  - > Handle eol and space character cases , TerminatingChar = C , Word = []

## Morphs (WordList, AtomList)

Convert list of words (as character lists from readCharLists, for example) to list of atoms, applying morphological rules to split off punctuation and the possessive " 's ".

morphs ([],[]).
morphs ([Word | RestOfWords], Atoms) : morph (Word, Atom)
 , morphs (RestOfWords, RestOfAtoms)
 , append (Atom, RestOfAtoms, Atoms).

# morph (Word, ItsAtoms)

- Convert one word, as a list of characters, to its corresponding atoms.
  - > More than one atom occurs when punctuation is split off, as punctuation is treated as an atom separately from a word.

```
morph ([],[]).
```

- morph ( Word , ItsAtoms ) :-
  - > Use the available rules for morphing a word to a list of component character lists
    - morphrules (Word, WordComponents)
  - > Convert each list of character codes to its corresponding atom
    - , maplist ( name , ItsAtoms , WordComponents ) .

## morphrules (CharList, ComponentLists)

ComponentLists is a sequence of sublists of CharList determined by the splitOff rules

morphrules ( CharList , ComponentLists ) :-

- > Do any split off rules apply?
  - (append (X,Y,CharList)
    - , splitOff (Y)
    - , ComponentLists = [X,Y])
- > Nothing to split off so only one sublist
  - ; ComponentLists = [ CharList ].

# splitOff ( String )

List of strings that are to be split off from words

> Apostrophe s
splitOff ( "'s" ) .

> Question mark
splitOff ( "?" ) .

> Period
splitOff ( "." ) .

# maplist ( P , Arg1 , Arg2 )

- maplist is a predicate that is the equivalent to the Lisp
   mapcar but restricted to one argument
- maplist applies the predicate P to every item in Arg1 and the result is the corresponding item in Arg2.

```
maplist ( _ , [ ] , [ ] ).
```

```
maplist ( P , [ H1 | T1 ] , [ H2 | T2 ] ) :-
```

> Q is the predicate P (H1, H2). The operator =... defines the correspondence of the compound term Q with the list form on the the right.

```
Q =.. [ P, H1, H2 ]
```

```
, call ( Q )
```

```
, maplist (P, T1, T2).
```

#### **Parse rules**

The parse rules analyze the list of atoms in a sentence. The relevant parts are extracted and rearranged for the respondTo rules.

parse ( semantic\_sentence\_representation
 , the\_sentence\_to\_parse
 , remainder\_of\_sentence )

> First rule creates the term stop to terminate the program.

parse ( stop , [ 'Stop' , '.' ] , [ ] ) .

> Last rule matches everything to create the term noparse for the "Can't parse that" response

parse (noparse, \_ , \_ ).

# Parsing "\_ is a \_."

A rule to parse sentences of the form
 John is a person.

- The parsing part of the rule
  parse (Clause) -->
  thing (Name), [is, a], type (T), ['.'].
- ♦ Where

thing (Name) --> [Name]. type (T) --> [T].

- This does not look like Prolog syntax
- What is happening?

#### Parse rule translations

- The previous syntax is in the library of predicates that comes with Edinburgh Prolog
- The predicates define a correspondence with the previous syntax and pure Prolog syntax

Why do we need the predicates?

Writing parsing rules in pure Prolog is tedious

# Parsing "P is a T."

- Syntax as entered in chat Looks fairly straight forward parse (Clause) --> [P], [is, a], [T], ['.'].
- Its equivalent in Prolog compared to the translation parse (Clause, S, Srem) :- det1 (S, S0), det2 (S0, S1), det3 (S1, S2), det4 (S2, Srem).

Query: parse(Clause, [ John, is, a, person, '.' ], \_ ) det1 ([ P | St ], St ). P = John St = [ is , a , person , '.' ] det2 ([ is , a | St ], St ). St = [ person, '.' ] det3 ([ T | St ], St ). T = person St = [ '.' ] det4 ([ '.' | St ], St ). St = [ ] ==> Srem = [ ]

#### Parsing "\_ is a \_." and translation

```
Looks fairly straight forward
     parse (Clause) -->
              thing (Name), [ is , a ], type (T), ['.'].
     thing (Name) --> [Name].
     type(T) --> [T].
                                     compared to the translation
In Prolog is the following
     parse (Clause, S, Srem) :-
           thing (Name, S, S0), det5 (S0, S1)
          , type (T, S1, S2), det6 (S2, Srem).
     thing (Name, S, Srem) :- det7 (S, Srem).
     type (T, S, Srem) :- det8 (S, Strem).
     det5 ([is, a] | St], St). det6 (['.'] | St], St).
     det7 ([Name | St], St). det8 ([T | St], St).
```

#### Semantic representation of a parse

- ♦ We can parse a sentence. So what?
- Need to get a semantic representation for the parse so the respondTo can work.
- That is the role played by the Clause variable in the parse rules

#### Parsing "\_ is a \_." and semantics

- Query: parse (Clause, [John, is, a, person, '.'], \_).
- **The parsing part of the rule**

parse (Clause) --> thing ( Name ) , [ is , a ] , type ( T ) , [ '.' ]

Makes the binding
 Name = John
 T = person

**The semantic part of the rule** 

{...} indicates do not
translate ..., keep as it
is, in the translated rule

# thing (X) & type (X)

 For things we want to check they begin with an upper case letter (capital letter)

```
thing (Name) --> [Name], { capital (Name) }.
```

For types we want to check that it begins with a lower case letter.

```
type (T) --> [T], { not (capital (T)) }.
```

 Rule for determining if a letter is an upper case (capital) letter or not.

> > Character with ASCII code less than 96 means it is an upper case letter.

capital (Name) :- name (Name,  $[FI_]$ ), F < 96.

# Parsing "A \_ is a \_."

The complete rule for parsing sentences like the following
 A woman is a person.

> The parsing part

parse( Clause ) --> [ 'A' ] , type ( T1 ) , [ is , a ] , type ( T2 ) , [ '. ']

> The semantic part

, { Head =.. [T2, X] , Condition =.. [ T1, X ]
 , Clause = (Head :- Condition) , ! } .

The following bindings occur

T1 = womanT2 = personparseHead = person (X)semantics, X is a variableCondition = woman (X)semantics, same XClause = person (X) :- woman (X)semantics

# Parsing "Is \_ a \_?"

- The complete rule for parsing sentences like the following  $\Diamond$ Is Mary a person? > The parsing part parse(Clause) --> ['ls'], thing(Name), [a] , type(T), ['?'] > The semantic part , { Goal =.. [ T, Name ] , Clause = ( '?-' ( Goal ) ) , ! } . Using the example the following bindings occur  $\Diamond$ Name = Mary T = person parse semantics Goal = person (Mary)
  - Clause = ?-(person (Mary )) semantics

?- makes Clause functor unique, correct respondTo is used.

## RespondTo

 The following two clauses are the response to stopping the program and to not finding a parse.

> > The argument is the semantic representation formed in the semantic part of parse rules

respondTo ( stop ) :- write ( 'All done.' ) , nl , ! .

```
respondTo (noparse) :-
write ('Can"t parse that.'), nl, !.
```

#### **RespondTo – enter into database**

 The following matches all clauses, so it would be last on the list

# > It adds the clause to the database – at the beginning

respondTo (Clause) :- asserta (Clause) , write ('Ok'), nl, !.

- ♦ **assertz(Clause)** add at the end of the database
- retract(X) find a clause in the database that matches the argument and remove it from the database

## **RespondTo – Yes/No query**

♦ Match functor **?-** and argument Goal.

> ?- is used to provide a respondTo to correspond to a particular parse rule.

> The operator -> tries to establish the goals to its left. If they succeed, then the goals to its right are attempted

```
respondTo ( '?-' ( Goal ) ) :-
```

```
(Goal -> write ('Yes'); write ('No'))
```

```
, ! , nl , nl .
```

In the case of the "Is Mary a person?" query we only need a yes and no answer.