

SC/MATH 1090

1- Boolean Formulae

Ref: G. Tourlakis, *Mathematical Logic*, John Wiley & Sons, 2008.

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Overview

- Boolean syntax
 - Boolean Alphabet
 - Strings
 - Formula Calculation; well-formed-formula (WFF)
 - Parsing (top-down and bottom-up)
 - Removing redundant brackets
 - Complexity of formulae

Boolean Alphabet

1. Symbols for Boolean or propositional variables

p, q, r with or without primes or subscripts

Examples: p, p', p_{123}, q''_{45}

2. Symbols for Boolean constants

\top called top, verum, or symbol “true”

\perp called bottom, falsum, or symbol “false”

3. Brackets, (and)

4. Boolean connectives

$\neg, \wedge, \vee, \rightarrow, \equiv$

Strings or Expressions

- Definition:

A **string** (or word, or expression) over a given alphabet is any ordered sequence of the alphabet's symbols, written adjacent to each other without any visible separators (no commas or spaces, etc).

- Examples:

- $(p \vee \perp)$ is a string given Boolean alphabet.
- $(p \sim q)$ is not a string given Boolean alphabet.
- $(p \rightarrow q)$ and $\rightarrow p)q($ are two **different** strings given the Boolean alphabet. Note only the ordering is different.

Strings (cont.)

- String variables
 - Denoted by A, B, C , etc with or without primes or subscripts
- Concatenation
 - Example: if A is abc and B is de (given the English alphabet), then \mathbf{AB} is $abcde$
- Empty string
 - Denoted by ε
 - $A\varepsilon = \varepsilon A = A$
- Substring
 - “ B is a substring of A ” means that for some string C and D we have $A = CBD$
 - If B is a substring of A and $B \neq A$, then B is a **proper substring** of A .

Formula calculation

Procedural definition

- Formula calculation is any finite (ordered) sequence of strings that we may write respecting the following requirements:
 1. At any step, we may write a Boolean variable or a Boolean constant
 2. At any step, we may write $(\neg A)$, provided we have already written string A in a previous step.
 3. At any step, we may write any of the strings $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, $(A \equiv B)$ provided we have already written strings A and B in a previous step.

Well-formed-formula (wff)

- A string A over the Boolean alphabet is called a **Boolean Expression** or a **well-formed-formula** iff it is a string written at some step of some formula-calculation.
 - Examples:
 $(p \equiv q)$
 $((p \vee r) \rightarrow (\neg q))$
- WFF: set of all well-formed-formulae (wffs)
- Bottom- up parsing of a wff is showing the procedural formula calculation steps.

Recursive definition of WFF

- The set of all well-formed-formulae is the smallest set of strings, WFF, that satisfies
 1. All Boolean variables (p, q, r, \dots), and constants (\top, \perp)
 2. If A and B are any strings in WFF, then so are the strings $(\neg A)$, $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, $(A \equiv B)$
- Top-down parsing of a wff is showing the recursive formula calculation steps.
- How do we know recursion terminates?
- The two definitions for WFF are equivalent.

Immediate Predecessors (i.p.)

1. Boolean variables or constants don't have any immediate predecessors
 2. A is an immediate predecessor of $(\neg A)$
 3. A and B are immediate predecessors of $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, $(A \equiv B)$
- We will prove later that the i.p.s are unique for each formula.

Removing brackets

- Redundant brackets
 - Outermost brackets are redundant
 - Any pair of brackets is redundant if its presence can be understood from the priority of the connectives
- Priorities:
 - The order of priorities (decreasing) is agreed to be
 $\neg, \wedge, \vee, \rightarrow, \equiv$
 - For same connectives, the rightmost has the highest priority
- Least parenthesized notation (LPN): writing wff with all redundant brackets removed
 - Note writing wff in LPN is just a short notation and is not a correctly written formula (by formula calculation)

Complexity

- The **complexity** of a formula is the number of connectives occurring in the formula
- The complexity of Boolean variables and constants is zero (they are also called **atomic** formulae)
- Example
 - Complexity of $((p \vee r) \rightarrow (\neg q))$ is 3