

# Mathematical Induction

- Very simple
- Very powerful proof technique
- “Guess” and verify strategy

# Basic steps

- Hypothesis:  $P(n)$  is true for all positive integers  $n$
- **Base case**/basis step (starting value)
- **Inductive step**

Formally:

$$[P(1) \wedge \forall k (P(k) \rightarrow P(k+1))] \rightarrow \forall n P(n)$$

# Intuition

Iterative modus ponens:

$P(k)$

$P(k) \rightarrow P(k+1)$

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$P(k+1)$

Need a starting point (Base case)

Proof is beyond the scope of this course

# Example 1

$$P(n): 1 + 2 + \dots + n = n(n+1)/2$$

Follow the steps:

- Base case:  $P(1)$ .

$$\text{LHS} = 1. \text{ RHS} = 1(1+1)/2 = \text{LHS}$$

- Inductive step:
  - Assume  $P(n)$  is true.
  - Show  $P(n+1)$  is true.

$$\begin{aligned} \text{Note: } & 1 + 2 + \dots + n + (n+1) \\ & = n(n+1)/2 + (n+1) = (n+1)(n+2)/2 \quad \text{done} \end{aligned}$$

## Example 2

- A difficult series (suppose we guess the answer)
- $1^2 + 2^2 + 3^2 + \dots + n^2 = n(n+1)(2n+1)/6$
- Base case:  $P(1)$  LHS = 1 = RHS.
- Inductive step:

$$\begin{aligned} 1^2 + 2^2 + 3^2 + \dots + n^2 + (n+1)^2 &= \\ n(n+1)(2n+1)/6 + (n+1)^2 &= \\ (n+1)(n+2)(2n+3)/6 &= \text{RHS}. \end{aligned}$$

# Proving Inequalities

- $P(n): n < 4^n$
- Base case:  $P(1)$  holds since  $1 < 4$ .
- Inductive step:
- Assume  $n < 4^n$
- Show that  $n+1 < 4^{n+1}$

$$n+1 < 4^n + 1 < 4^n + 4^n < 4 \cdot 4^n = 4^{n+1}$$

# More examples

- Sum of odd integers
- $n^3 - n$  is divisible by 3
- Number of subsets of a finite set

# Points to remember

- Base case does not have to be  $n=1$
- Most common mistakes are in not verifying that the base case holds
- Sometimes we need more than  $P(n)$  to prove  $P(n+1)$  – in these cases STRONG induction is used
- Usually guessing the solution is done first.



# How can you guess a solution?

- Try simple tricks: e.g. for sums with similar terms –  $n$  times the average or  $n$  times the maximum; for sums with fast increasing/decreasing terms, some multiple of the maximum term.
- Often proving upper and lower bounds separately helps.

# Strong Induction

- Equivalent to induction – use whichever is convenient

- Formally:

$$[ P(1) \wedge \forall k (P(1) \wedge \dots \wedge P(k) \rightarrow P(k+1)) ] \\ \rightarrow \forall n P(n)$$

- Often useful for proving facts about algorithms

# Examples

- **Fundamental Theorem of Arithmetic:** every positive integer  $n$ ,  $n > 1$ , can be expressed as the product of one or more prime numbers.
- every amount of postage of 12 cents or more can be formed using just 4-cent and 5-cent stamps.

# Fallacies/caveats

- “Proof” that all Canadians are of the same age!

<http://www.math.toronto.edu/mathnet/falseProofs/sameAge.html>