Creating a Class Beyond the Basics (pt 2)

Based on slides by Prof. Burton Ma

Arrays as Containers

- Suppose you have an array of unique
 - PhoneNumberS
 - How do you compute whether or not the array

```
public static boolean
       hasPhoneNumber(PhoneNumber p,
                      PhoneNumber[] numbers)
  if (numbers != null) {
    for( PhoneNumber num : numbers ) {
      if (num.equals(p)) {
        return true;
  return false;
```

- Called *linear search* or *sequential search*
 - Doubling the length of the array doubles the amount of searching we need to do
- If there are n PhoneNumbers in the array:
 - Best case: the first PhoneNumber is the one we are searching for → 1 call to equals()
 - → N calls to equals()
 - Average case: the PhoneNumber is somewhere in the middle of the array → approximately (n/2) calls to equals()

hashCode()

- If you override equals() you must override hashCode()
 - Otherwise, the hashed containers won't work

Hash Tables

You can think of a hash table as being an array of buckets where each bucket holds the stored objects

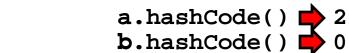
0	1	2	3	•••	N

Insertion into a Hash Table To insert an object a, the hash table calls

a.hashCode() method to compute which

bucket to put the object into

c.hashCode() N d.hashCode() N



0	1	2	3	• • •	N

means the hash table takes the hash code and does something to it to make it fit in the range **0–N**

Search on a Hash Table

To see if a hash table contains an object a, the hash table calls a.hashCode() method to compute which bucket to look for a in a.hashCode() > 1

b	a.equa	als(a) true		z.equa z.equa	ls(c) ls(d) false
0	1	2	3	• • •	N

- Searching a hash table is usually much faster than linear search
 - Doubling the number of elements in the hash table usually does not noticably increase the amount of search needed
- ▶ If there are n PhoneNumbers in the hash table:
 - Best case: the bucket is empty, or the first PhoneNumber in the bucket is the one we are searching for → 0 or 1 call to equals()
 - Worst case: all n of the PhoneNumbers are in the same bucket → N calls to equals()
 - Average case: the PhoneNumber is in a bucket with a small number of other PhoneNumbers → a small number of calls to equals()

Object hashCode()

- If you don't override hashCode(), you get the implementation from Object.hashCode()
 - Object.hashCode() uses the memory address of the object to compute the hash code

```
// client code somewhere
PhoneNumber pizza = new PhoneNumber(416, 967, 1111);
HashSet<PhoneNumber> h = new HashSet<PhoneNumber>();
h.add(pizza);
PhoneNumber pizzapizza = new PhoneNumber(416, 967, 1111);
System.out.println( h.contains(pizzapizza) ); // false
```

- Note that pizza and pizzapizza are distinct objects
 - Therefore, their memory locations must be different
 - Therefore, their hash codes are different (probably)
 - Therefore, the hash table looks in the wrong bucket (probably) and does not find the phone number even though pizzapizza.equals(pizza)

A Bad (but legal) hashCode()

public final class PhoneNumber {
 // attributes, constructors, methods ...

```
@Override public int hashCode()
{
   return 1; // or any other constant int
}
```

}

This will cause a hashed container to put all PhoneNumbers in the same bucket

A Slightly Better hashCode()

```
public final class PhoneNumber {
    // attributes, constructors, methods ...
@Override public int hashCode()
    {
        return (int)(this.getAreaCode() +
            this.getExchangeCode() +
            this.getStationCode());
    }
```

- The basic idea is generate a hash code using the attributes of the object
- It would be nice if two distinct objects had two distinct hash codes
 - But this is not required; two different objects can have the same hash code
- It is required that:
 - 1. If x.equals(y) then x.hashCode() == y.hashCode()
 - 2. x.hashCode() always returns the same value if x does not change its state

Something to Think About

What do you need to be careful of when putting a mutable object into a HashSet?