
Prolog Core Concepts and Notation

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Adapted from Peter Roosen-Runge

Readings: C & M Ch 1, 2, 3.1-3.3, 8

declarative/logic programming

- ◆ idea: write a program that is a logical theory about some domain and then query it
- ◆ most well known instance is Prolog
- ◆ core constructs, terms and statements, are inherited from first order logic

terms

- ◆ Prolog statements express relationships among *terms*
- ◆ terms are (a generalization) of the same notion in first order logic, i.e. a constant, a variable, or a function applied to some argument terms
- ◆ E.g. john, john_smith, X, Node, _person, fatherOf(paul), date(25,10,2005)
- ◆ fatherOf and date are functors; date has arity 3; it takes 3 arguments

terms

- ◆ variables begin with upper-case letter or _
- ◆ constants and functors (symbols) begin with lower-case
- ◆ terms denote objects
- ◆ compound terms are called structures
- ◆ E.g. course(complexity,time(Monday,9,11),lecturer(patrick,dymond),location(CSE,3311))
- ◆ used to represent complex data
- ◆ terms (usually) have a tree structure

facts

- ◆ *facts* are like atomic formulas in first order logic.
- ◆ syntax is same as terms, but ending with a period.
- ◆ e.g. fatherOf(paul,henry).
mortal(ulyssus).
likes(X,iceCream).
likes(mary,brotherOf(helen)).
- ◆ variables are implicitly universally quantified.

rules

- ◆ *rules* are conditional statements.
- ◆ e.g. mortal(X) :- human(X).
i.e. $\forall x \text{ Human}(x) \rightarrow \text{Mortal}(x)$,
all humans are mortal.
- ◆ daughter(X,Y) :- father(Y,X), female(X).
- ◆ , represents conjunction.
- ◆ likes(mary,X) :- isSweet(X).

rules

- ◆ `ancestor(X,Y) :- father(X,Z), ancestor(Z,Y).`
- ◆ variables are universally quantified from outside; can think of variables that appear only in rule body as existentially quantified.

queries

- ◆ A query asks whether a given statement is true, i.e. whether it follows from the program.
- ◆ e.g. `?- mortal(ulyssus).` given `mortal(X) :- human(X).`
`human(ulyssus).` `human(penelope).`
`god(zeus).`
Prolog answers Yes

queries

- ◆ `?- mortal(X).`
X = ulyssus ;
X = penelope
Yes
- ◆ variables in queries are existentially quantified; can be used to retrieve information.
- ◆ can have conjunctive queries, e.g.
`?- mortal(X), mortal(Y), not(X = Y).`

lists

- ◆ lists are a special kind of term that allows arbitrary number of components
- ◆ `[]` is the empty list
- ◆ `.(a,b)` is a dotted pair
- ◆ `[a, b, c] = .(a,.(b,.(c,[])))` is a list of 3 components.
- ◆ the functor `.` builds binary trees
- ◆ can use `display(X)` to print internal representation of X

lists

- ◆ can refer to the first and rest of a list using the notation: [First | Rest]
- ◆ e.g. ?- X = [a,b,c], X = [F|R].
X = [a,b,c]
F = a
R = [b,c]
- ◆ E.g. X = [b], Y = a, Z = [Y|X].
X = [b]
Y = a
Z = [a,b]

unification

- ◆ this was an instance of the kind of pattern matching called unification that Prolog performs
- ◆ Prolog tries to find a way to instantiate the variables (substitute terms for them) that satisfies the query
- ◆ more on this later

terms can represent graphs

- ◆ ?- $X = [a|X]$.
 $X = [a, a, a, a, a, a, a, a, a|...]$
Yes
- ◆ here X denotes an infinite or circular list
- ◆ this is not allowed in first-order logic; a variable cannot denote a term and one of its subterms; but Prolog omits the “occurs check”

building a knowledge base

- ◆ to be used in a computation, facts and rules must be stored in the (dynamic) database
- ◆ facts and rules get into the database through *assertion* and *consultation*
- ◆ consultation loads facts and rules from a file

assertion

- ◆ ?- assert(human(ulyssus)).
- ◆ ?- human(X).
X = ulyssus
Yes
- ◆ assertion can be done dynamically
- ◆ also retract to remove facts and rules from the DB
- ◆ like assignment, change state; avoid when possible

consultation

- ◆ ?- consult('family.pl').
loads facts and rules from file family.pl
- ◆ ?- [family].
does the same thing
- ◆ ?- [user].
lets you enter facts and rules from the keyboard

denotation/meaning of Prolog programs

- ◆ a Prolog program defines a set of relations, i.e. specifies which tuples of objects/terms belong to a particular relation
- ◆ in logic, this is called a model
- ◆ declarative programming is very different from usual procedural programming where programs perform many state changing operations

denotation of Prolog program e.g.

- ◆ fatherOf(john,paul).
fatherOf(mary,paul).
motherOf(john,lisa).
parentOf(X,Y) :- fatherOf(X,Y).
parentOf(X,Y) :- motherOf(X,Y).
- ◆ fatherOf is the relation {<john,paul>, <mary,paul>}
- ◆ what is the relation associated with motherOf and parentOf?

rules as procedures

- ◆ rule has form goal :- body
- ◆ goal or head is like name of procedure
- ◆ terms on the RHS are like the body of the procedure, the sub-goals that have to be achieved to show that the goal holds
- ◆ the sub-goals will be attempted left-to-right
- ◆ rule succeeds if all sub-goals succeed

passing values

- ◆ calling/querying a goal can instantiate its variables
- ◆ a sub-goal's success can bind a variable within it, also binding the same variable in the goal
- ◆ binding or instantiating a variable is giving it a value
- ◆ compare to passing values into or out of a procedure

passing values e.g.

- ◆ Assume program:
motherOf(john,lisa).
parentOf(X,Y) :- motherOf(X,Y).
- ◆ Queries:
?- parentOf(john,X).
X = lisa Yes
?- parentOf(X,lisa).
X = john Yes
?- parentOf(X,Y).
X = john, Y = lisa Yes
- ◆ No fixed input and output parameters

relational thinking

- ◆ in Prolog, formulate statements about function values as relational facts, e.g.
factorial(0,1).
factorial(N,M):- K is N -1, factorial(K,L),
M is N * L.
- ◆ to compose functions, e.g. $Y = f(g(X))$,
you must name intermediate results
fg(X,Y):- g(X,Z), f(Z,Y).

almost everything is syntactically a term

- ◆ lists are terms; what is the functor?
- ◆ rules are terms:
grandfather(X,Y):- father(X,Z), father(Z,Y).
What are the functors?
- ◆ queries are terms

arithmetic functions

- ◆ Prolog retains arithmetic functions as functions (more intuitive):
?- X **is** exp(1). % exp(1) = e¹
X = 2.71828
Yes
?- X is (4 + 2) * 5.
X = 30
Yes
- ◆ How does is compare with =, assignment?

operators

- ◆ some functors are represented by *infix* or *prefix* or *postfix* operators
- ◆ Some infix operators: `is`, `=`, `+`, `*`, `/`, `mod`, `>`, `>=`, `:-`, `,`, etc.
- ◆ `+` and `-` are both prefix and infix
- ◆ `:-` as prefix is a command, used for declarations
- ◆ operators have precedence
- ◆ can define our own operators

help is sometimes helpful

?- help(reverse).

reverse(+List1, -List2)

Reverse the order of the elements in List1 and unify the result with the elements of List2.

+arg: arg is input and should be instantiated.

-arg: arg is output and can be initially uninstantiated; if the query succeeds, the arg is instantiated with the "output" of the query.

?arg: arg can be either input or output

online help

```
?- help(lists).  
No help available for lists  
Yes  
?- apropos(lists).  
merge/3           Merge two sorted lists  
append/3          Concatenate lists  
Section 11-1      "lists: List Manipulation"  
Section 15-2-1   "lists"  
Yes  
?- help(append/3).  
append(?List1, ?List2, ?List3)  
  Succeeds when List3 unifies with the concatenation of List1 and  
  List2. The predicate can be used with any instantiation pattern  
  (even three variables).
```

examples

```
?- append([a,b],[c],X).  
X = [a, b, c]  
  
Yes  
?- append(X,[c],[a,b,c]).  
X = [a, b]  
  
Yes  
?- append([a,b],[c],[a,b,d]).  
  
No
```

more examples

```
?- append([a,b],X,Y).
X = _G187
Y = [a, b|_G187]
Yes
?- append(X,Y,Z).
X = []
Y = _G181
Z = _G181 ;

X = [_G262]
Y = _G181
Z = [_G262|_G181] ;

X = [_G262, _G268]
Y = _G181
Z = [_G262, _G268|_G181]
```

append is an example of a reversible or **steadfast** predicate (Richard O'Keefe)

reversible programming

- ◆ good predicates are steadfast
- ◆ they give correct answers even if unusual values are supplied
 - e. g. variables for inputs, constants for outputs
- ◆ non-steadfast predicates require specific arguments to be instantiated (input) or variables (output)

unification

- ◆ Prolog matches terms by *unifying* them, i.e. applying a most general unifier to them
- ◆ it instantiates variables as little as possible to make them match, e.g.

?- $X = f(Y,b,Z)$, $X = f(a,V,W)$.

$X = f(a, b, _G182)$

$Y = a$

$Z = _G182$

$V = b$

$W = _G182$

family relations example

family relations

- ◆ the database:

rules

parent(Parent, Child) :- mother(Parent, Child).

parent(Parent, Child) :- father(Parent, Child).

facts

father('George', 'Elizabeth'). father('George', 'Margaret').

mother('Mary', 'Elizabeth'). mother('Mary', 'Margaret').

- ◆ Note encoding of disjunction

finding all solutions

| ?- parent(Parent, Child).

Parent = 'Mary',

Child = 'Elizabeth' ;

Parent = 'Mary',

Child = 'Margaret' ;

Parent = 'George',

Child = 'Elizabeth' ;

Parent = 'George',

Child = 'Margaret' ;

no

how prolog finds solutions

```
trace] ?-
  parent(Parent, Child1),
  parent(Parent, Child2), not
  (Child1 = Child2).
  Call: (8) parent(_G313,
    _G314) ? creep
  Call: (9) mother(_G313,
    _G314) ? creep
  Exit: (9) mother('Mary',
    'Elizabeth') ? creep
  Exit: (8) parent('Mary',
    'Elizabeth') ? creep
  Exit: (8) parent('Mary',
    'Elizabeth') ? creep
  Call: (8) parent('Mary',
    _G317) ? creep
  Call: (9) mother('Mary',
    _G317) ? creep

  Exit: (9) mother('Mary',
    'Elizabeth') ? creep
  Exit: (8) parent('Mary',
    'Elizabeth') ? creep
  Redo: (9) mother('Mary',
    _G317) ? creep
  Exit: (9) mother('Mary',
    'Margaret') ? creep
  Exit: (8) parent('Mary',
    'Margaret') ? creep

  Parent = 'Mary'
  Child1 = 'Elizabeth'
  Child2 = 'Margaret'
```

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Prolog's query answering process

- ◆ a query is a conjunction of terms
- ◆ answer to the query is *yes* if all terms *succeed*
- ◆ A term in a query *succeeds* if
 - ❖ it matches a fact in the database *or*
 - ❖ it matches the head of a rule whose body succeeds
- ◆ the substitution used to unify the term and the fact/head is applied to the rest of the query
- ◆ works on query terms in **left to right order**;
databases facts/rules that match are tried in **top to bottom** order

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recursion examples

generating permutations

- ◆ A permutation P of a list L is a list whose first is some element E of L and whose rest is a permutation of L with E removed.
- ◆ [] is a permutation of []
- ◆ In Prolog:
permutation([], []).
permutation(L, [E|PR]) :- select(E, L, R),
permutation(R, PR).

selecting an element from a list

- ◆ To select an element from a list, can either select the first leaving the rest, or select some element from the rest and leaving the first plus the unselected elements from the rest.
- ◆ In Prolog:
select(X,[X|R],R).
select(X,[Y|R],[Y|RS]):- select(X,R,RS).

sorting by the definition

- ◆ Find a permutation that is ordered
sort(L,P):- permutation(L,P),
ordered(P).
ordered([]).
ordered([E]).
ordered([E1,E2|R]) :- E1 <= E2,
ordered([E2|R]).
- ◆ an example of “generate and test”

reverse

- ◆ reverse(L,RL) holds if RL is a list with the components of L reversed
- ◆ ordinary recursive definition

```
reverse([],[]).
reverse([F|R],RL):- reverse(R,RR),
    append(RR, [F], RL).
append([],L,L).
append([F|R],L,[F|RL]):-
    append(R,L,RL).
```

reverse

- ◆ Tail recursive definition:

```
reverse(L,RL):- reverse(L,[],RL).
reverse([],Acc,Acc).
reverse([F|R],Acc,RL):-
    reverse(R,[F|Acc],RL).
```
- ◆ recursive call is last thing done
- ◆ can avoid saving calls on stack

solving a logic puzzle with Prolog

the zebra puzzle

1. There are 5 houses, occupied by politically-incorrect gentlemen of 5 different nationalities, who all have different coloured houses, keep different pets, drink different drinks, and smoke different (now-extinct) brands of cigarettes.
2. The Englishman lives in a red house.
3. The Spaniard keeps a dog.
4. The owner of the green house drinks coffee.
- ...
6. The ivory house is just to the left of the green house.
- ...
11. The Chesterfields smoker lives next to a house with a fox.

Who owns the zebra and who drinks water?

Prolog implementation

- ◆ represent the 5 houses by a structure of 5 terms
house(Colour, Nationality, Pet, Drink, Cigarettes)
- ◆ create a partial structure using variables, to be filled by the solution process
- ◆ specify constraints to instantiate variables

house building

```
makehouses(0,[]).
```

```
makehouses(N,[house(Col, Nat, Pet, Drk, Cig)|List])  
:- N>0, N1 is N - 1, makehouses(N1,List).
```

or more cleanly with anonymous variables:

```
makehouses(N,[house(_,_,_,_)|List])  
:- N>0, N1 is N - 1, makehouses(N1,List).
```

Why is this equivalent? (See p. 159.)

the empty houses

?- makehouses(5, List).

```
List = [house(_G233, _G234, _G235, _G236, _G237),
        house(_G245, _G246, _G247, _G248, _G249), house
        (_G257, _G258, _G259, _G260, _G261), house(_G269,
        _G270, _G271, _G272, _G273), house(_G281, _G282,
        _G283, _G284, _G285)]
```

constraints

- ◆ The Englishman lives in a red house.
house(red, englishman, _, _, _) on List,
- ◆ The Spaniard keeps a dog.
house(_, spaniard, dog, _, _) on List,
- ◆ The owner of the green house drinks coffee.
house(green, _, _, coffee, _) on List
- ◆ The ivory house is just to the left of the green house
sublist2([house(ivy, _, _, _, _), house(green, _, _, _, _)], List),
- ◆ The Chesterfields smoker lives next to a house with a fox.
nextto(house(_, _, _, chesterfields), house(_, _, fox, _, _), List),

defining the on operator

- ◆ on is a user-defined infix operator that is a version of member/2
- ◆ :- op(100,zfy,on).
X on List :- member(X,List).
amounts to
 X on [X|_].
 X on [_|R] :- X on R.

predicates for defining constraints

- ◆ “just to the left of”? “lives next to”?
- ◆ define sublist(S,L)
 sublist2([S1, S2], [S1, S2 | _]) .
 sublist2(S, [_ | T]) :- sublist2(S, T).
- ◆ define nextto predicate
 nextto(H1, H2, L) :- sublist2([H1, H2], L).
 nextto(H1, H2 ,L) :- sublist2([H2, H1], L).

translating the constraints

- ◆ The ivory house is just to the left of the green house

```
sublist2( [house(ivy, _, _, _, _),  
          house(green, _, _, _, _)], List),
```
- ◆ The Chesterfields smoker lives next to a house with a fox.

```
nextto(house( _, _, _, _, chesterfields),  
       house( _, _, fox, _, _), List),
```

looking for the zebra

- ◆ Who owns the zebra and who drinks water?

```
find(ZebraOwner, WaterDrinker) :-  
  makehouses(5, List),  
  house(red, englishman, _, _, _) on List,  
  ... % all other constraints  
  house( _, WaterDrinker, _, water, _) on List,  
  house( _, ZebraOwner, zebra, _, _) on List.
```
- ◆ solution is generated and queried in the same clause
- ◆ neither water or zebra are mentioned in the constraints

solving the puzzle

```
?- [zebra].  
% zebra compiled 0.00 sec, 5,360 bytes
```

```
Yes  
?- find(ZebraOwner, WaterDrinker).
```

```
ZebraOwner = japanese  
WaterDrinker = norwegian ;
```

```
No
```

how Prolog finds solution

After first 8 constraints:

```
List = [  
house(red, englishman, snail, _G251, old_gold),  
house(green, spaniard, dog, coffee, _G264),  
house(ivory, ukrainian, _G274, tea, _G276),  
house(green, _G285, _G286, _G287, _G288),  
house(yellow, _G297, _G298, _G299, kools)]
```

how Prolog solves the puzzle

Then need to satisfy “the owner of the third house drinks milk”, i.e.

```
List = [_ , _ , house( _ , _ , _ , milk, _), _ , _],
```

Can't be done with current instantiation of List. So Prolog will **backtrack** and find another.

how Prolog solves the puzzle

The unique complete solution is

```
L = [  
house(yellow, norwegian, fox, water, kools),  
house(blue, ukrainian, horse, tea, chesterfields),  
house(red, englishman, snail, milk, old_gold),  
house(ivory, spaniard, dog, orange,  
lucky_strike),  
house(green, japanese, zebra, coffee,  
parliaments)]
```

See course web page for code of the example.