

Math/CSE 1019:
Discrete Mathematics for Computer Science
Fall 2011

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Course page: <http://www.cs.yorku.ca/course/1019>

Mathematical Induction

- Very simple
- Very powerful proof technique
- “Guess” and verify strategy

Basic steps

- Hypothesis: $P(n)$ is true for all positive integers n
- **Base case**/basis step (starting value)
- **Inductive step**

Formally:

$$[P(1) \wedge \forall k (P(k) \rightarrow P(k+1))] \rightarrow \forall n P(n)$$

Intuition

Iterative modus ponens:

$P(k)$

$P(k) \rightarrow P(k+1)$

$P(k+1)$

Need a starting point (Base case)

Proof is beyond the scope of this course

Example 1

$$P(n): 1 + 2 + \dots + n = n(n+1)/2$$

Follow the steps:

- Base case: $P(1)$.

$$\text{LHS} = 1. \text{ RHS} = 1(1+1)/2 = \text{LHS}$$

- Inductive step:
 - Assume $P(n)$ is true.
 - Show $P(n+1)$ is true.

$$\begin{aligned} \text{Note: } & 1 + 2 + \dots + n + (n+1) \\ & = n(n+1)/2 + (n+1) = (n+1)(n+2)/2 \quad \text{done} \end{aligned}$$

Example 2

- A difficult series (suppose we guess the answer)
- $1^2 + 2^2 + 3^2 + \dots + n^2 = n(n+1)(2n+1)/6$
- Base case: $P(1)$ LHS = 1 = RHS.
- Inductive step:

$$\begin{aligned} 1^2 + 2^2 + 3^2 + \dots + n^2 + (n+1)^2 &= \\ n(n+1)(2n+1)/6 + (n+1)^2 &= \\ (n+1)(n+2)(2n+3)/6 &= \text{RHS}. \end{aligned}$$

Proving Inequalities

- $P(n): n < 4^n$
- Base case: $P(1)$ holds since $1 < 4$.
- Inductive step:
- Assume $n < 4^n$
- Show that $n+1 < 4^{n+1}$

$$n+1 < 4^n + 1 < 4^n + 4^n < 4 \cdot 4^n = 4^{n+1}$$

Points to remember

- Base case does not have to be $n=1$
- Most common mistakes are in not verifying that the base case holds
- Sometimes we need more than $P(n)$ to prove $P(n+1)$ – in these cases **STRONG** induction is used
- Usually guessing the solution is done first.

How can you guess a solution?

- Try simple tricks: e.g. for sums with similar terms – n times the average or n times the maximum; for sums with fast increasing/decreasing terms, some multiple of the maximum term.
- Often proving upper and lower bounds separately helps.

More examples

- Sum of odd integers
- $n^3 - n$ is divisible by 3
- Number of subsets of a finite set

Strong Induction

- Equivalent to induction – use whichever is convenient

- Formally:

$$[P(1) \wedge \forall k (P(1) \wedge \dots \wedge P(k) \rightarrow P(k+1))] \\ \rightarrow \forall n P(n)$$

- Often useful for proving facts about algorithms

Examples

- **Fundamental Theorem of Arithmetic:** every positive integer n , $n > 1$, can be expressed as the product of one or more prime numbers.
- every amount of postage of 12 cents or more can be formed using just 4-cent and 5-cent stamps.

Fallacies/caveats

- “Proof” that all Canadians are of the same age!

<http://www.math.toronto.edu/mathnet/falseProofs/sameAge.html>