

#### The problem

- How do we use mobile robot(s) to acquire spatial models of unknown physical environments?
- localization and mapping)
- Robot(s) operate in an unknown environment
- Goal: build a map representation of the environment
- One fundamental problem: 'have I been here before?'
  - Disambiguate current place against known locations
- Two basic approaches to solving this problem
  - Metric information in probabilistic approaches
  - Resorting to an 'oracle' in deterministic approaches
- □ Here we are looking at deterministic approaches and the power of specific oracles



### **Basic world and robot model**

- and edges
- Explicit specification of edge ordering exists in each vertex Inaccessible to robot – anonymous graph
- □ Robot can move within the environment by traversing edges
- □ Robot can enumerate edges at a vertex in a consistent manner



- Robot cannot distinguish places using sensing Observing (infinite) sequence of two-door rooms
- □ At a 'new' place, consider <u>all</u> possibilities Maintaining multiple hypotheses (true / false)
- No deterministic algorithm exists!
  - May be infinite valid models
  - But we can use an oracle to make the problem solvable deterministically





# **Exploring Topological Environments**

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Mapping can be efficiently done using DFS, GREEDY, and the like •  $\bigcirc O(m) \le O(n^2)$  exploration cost. Minimum!



## The power of a 'strong' oracle

### $\Box$ *n* super-glue pebbles, *n* footprints, l = c(G) string, and the like

- Marks all the visited place, but not uniquely
- Disambiguation is needed
- Provides additional info that can be exploited to reduce disambiguation cost
- $\bigcirc$   $O(mn) \leq O(n^3)$  with reduced constant than a 'very weak' oracle

### The power of a 'fair' oracle

#### □ Less (≤ *n*) pebbles, shorter (l < c(G)) string, and the like

- Marks multiple visited places, but not uniquely
- Robot may run out of oracles during exploration
- Explores in phases -- pick up and reuse the oracles
- $\bigcirc O(mn) \le O(n^3)$ . Cost varies depending on the amount of the oracles

# **Empirical results with different oracles**



#### Summary

- □ For deterministic mapping of anonymous graphs, an oracle is needed!
- □ Mapping is possible with a `very weak' oracle (e.g., a super-glue pebble) has exploration cost  $O(m^2n) \leq O(n^5)$
- □ The minimum exploration cost  $O(m) \leq O(n^2)$  is possible with a 'super-strong' oracle (e.g. a paint can, a large number of pebbles, a very long string)
- □ Other 'medium' oracles are possible with exploration cost  $O(mn) \leq O(n^3)$

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