Abstract Data Types Documentation

Documentation

Users are only interested in the properties of the ADT

 Programmers and designers require all the information which a user needs AND all information pertaining to the design and implementation

Useful to think of the documentation as being an annotated definition of an abstract data type

Documentation Table of Contents

- Cover page, table of contents and abstract
- Document introduction
 - » Informal overview of the facilities provided. Help readers determine if this is what they need
- Data type objects
 - » Description of all the objects include diagrams
 - » Split into
 - > Imported which predefined objects are used
 - > Exported for others to use
 - > Hidden used in the implementation

- Operations
 - » Give
 - > Signature
 - > Informal description
 - > pre- and post- conditions
 - » Use natural language, mathematics, diagrams whatever best gets the meaning across.
 - » Be simple, complete, clear, precise, concise as possible

- Example partial axiomatic description of bank accounts
 - » The operation signatures only no pre- post- given

```
new : [] -> account
```

Create an account with a zero balance

```
withdraw: account X amount -> account
```

Remove amount from account

```
deposit : account X amount -> account
```

Add amount to account

```
balance : account -> amount
```

- What is the amount in the account?

- Operation interaction
 - » Previous section describes operations in isolation
 - » Provide better understanding by showing properties when operations are used in combination
 - » Common descriptive method in use is axiomatic
 - > List of axioms or statements which must be true if the ADT is implemented and used correctly

- Axioms about the data type
 - » Axiom 1: New account has a balance of zero dollars
 balance(new) = 0
 - » Axiom 2: Cannot withdraw from a new account
 withdraw(new, amt) = error
 - » Axiom 3: Deposit amt and then withdraw amt with no intervening operations the balance does not change

» Axiom 4: Only withdraw if the balance is ≥ the amount to withdraw. The amount is deducted from the balance

```
withdraw(acct, amt) =
  if balance(acct) < amt then error
  else balance(acct') = balance(acct) - amt
  fi</pre>
```

- How to use the ADT
 - » Tutorial guide on use. Dwell on nuances. Describe various examples
- Dictionary
 - » Define new terminology or domain specific jargon that implementers or users may not know
- Undesired Event Dictionary
 - » Description of possible errors which can occur
 - » Contains warnings
 - » How to recognize error situations
 - » How to recover from error situations
 - » What to do if recovery is impossible

- ADT generation parameters
 - » Describe how instances and variations can be implemented from this generic data type
 - > How to change base types
 - > How to change amount of storage for a customer name
 - » Describe changes that can be made that will not violate assumptions and specifications. Design for a class of similar data types
 - State what programming tools can be used to modify the implementation

- Design issues
 - » What were the design choices and why were the actual choices chosen. Help guide future changes to keep in the spirit of the original
 - > Why was fixed memory allocation used instead of dynamic?
 - > Why were size limits imposed?
 - > Why was a particular data structure chosen?

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- Implementation notes
 - » Designer may have information of use to the implementer. Know properties that can improve implementation
- List of assumptions those assumptions that
 - » Cannot be violated
 - » Not implicit in the context
 - » Global
 - » Note: cannot state all assumptions so state those that
 - > Are most important
 - > Most likely to cause problems if violated
 - > Are not easily detected as causing problems until a long time later

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- Normal use assumptions
 - » Information available from the ADT
 - » Information that must be supplied to the ADT
 - » Events reported by the ADT
 - » Tasks that can be performed by the ADT
 - » Operating states of the ADT and how they affect the Information obtained from and supplied to the ADT
 - » Failure states of the ADT and how they affect the information obtained from and supplied to the ADT

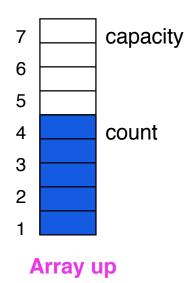
- Incorrect use assumptions
 - » Associated with run time undesired events
 - What may or may not happen if the production version has undesired event handling code removed to speed up the system
- Program source text
 - » If the source test is small may be included with the description of the operations
- Facilities index
 - » A quick look up reference of all programs, modules, operations, objects and terms defined

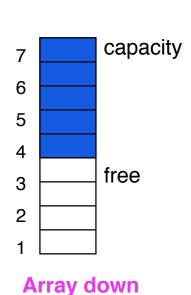
Minimal Documentation

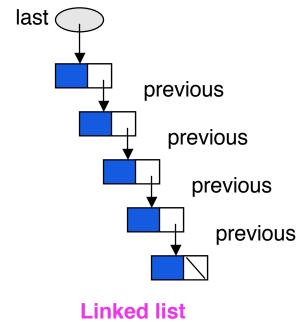
- Objects
 - » Types

Diagrams where possible

- Example stack
 - » Imported none
 - » Exported STACK [G]
 - » Hidden implementation







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Operations – example for a stack

```
>> Signatures, pre & post conditions
>> push : STACK [G] x G → STACK [G]

- require true
    ensure result = x ^ s & count = old count + 1
>> pop : STACK [G] → STACK [G]

- require not empty (s)
    ensure result = s' & count = old count - 1
>> top : STACK [G] → G

- require not empty (s)
    ensure result = s₁
```

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Operations – example for a stack cont'd

- » Note: often "require true" is not written but is assumed
- » It is better to write it as then one can wonder if it was left out by accident
 - > "nothing" is often represented with a special symbol. e.g. nil , λ , ϵ , Δ

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- Operations example for a stack cont'd
 - » axioms

```
> ∀ x : G, s : STACK [ G ] •

top ( push ( s, x ) ) = x

∧ pop ( push ( s, x ) ) = s

∧ empty ( new )

∧ ~ empty ( push ( s, x ) )
```

is read as"it is the case that"

» Alternately can use natural language

```
> forall x : G, s : STACK [ G } ::
          top ( push ( s, x ) ) = x
          and pop ( push ( s, x ) ) = s
          and empty ( new )
          and not empty ( push ( s, x ) )
```

ADT Invariants

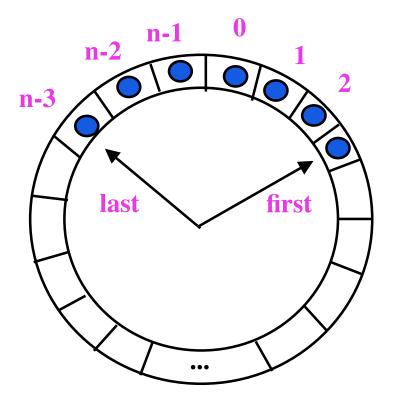
 Conditions that must be true after the execution of any method in the the class

- The conditions that hold, at all times, among the objects in an instance of the ADT
 - » More on this when we discuss design by contract

Example Circular Queue

isEmpty → length = 0 & (last-1) mod Size = first isFull → length = Size - 1 not isFull → length = (Size + first - last + 1) mod Size

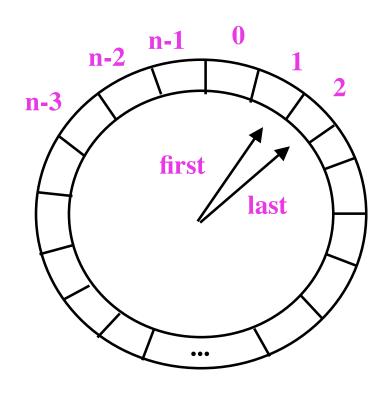
last is the last Item to remove



first is the first Item to remove

Empty Circular Queue

isEmpty → length = 0 & (last-1) mod Size = first isFull → length = Size - 1 not isFull → length = (Size + first - last + 1) mod Size

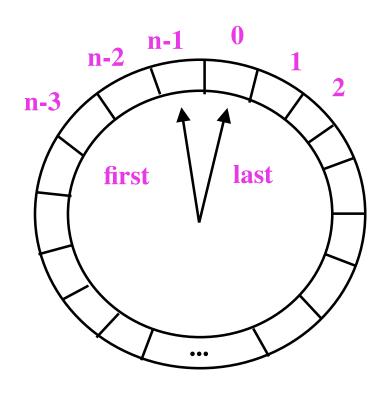


length =
$$(n + 1 - 2 + 1) \mod n$$

= $(n + 0) \mod n$
= 0

Empty Circular Queue – 2

isEmpty → length = 0 & (last-1) mod Size = first isFull → length = Size - 1 not isFull → length = (Size + first - last + 1) mod Size

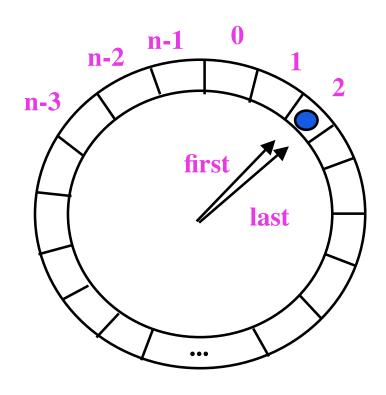


length =
$$(n + (n-1) - 0 + 1) \mod n$$

= $(2n + 0) \mod n$
= 0

Length 1 Circular Queue

```
isEmpty → length = 0 & (last-1) mod Size = first
isFull → length = Size - 1
not isFull → length = (Size + first - last + 1) mod Size
```



length =
$$(n + 2 - 2 + 1) \mod n$$

= $(n + 1) \mod n$
= 1

Longer length Circular Queue

length = (Size + first - last + 1) mod Size

length =
$$(n + 3 - (n-3) + 1) \mod n$$

= $(7) \mod n$
= 7

