Basic Lisp Operations

Function invocation

♦ It is an S-expression – just another list!

(function arg1 arg2 ... argN)

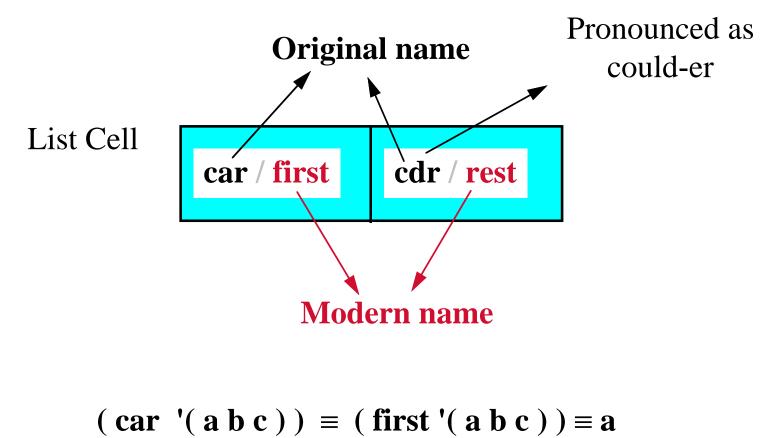
- ♦ First list item is the function **prefix notation**
- ♦ The other list items are the arguments to the function.
- Arguments can themselves be lists
 - » (+ 1 2 3 (+ 4 5 6) 7 8 9) ==> 45
 - » Outer + has 7 arguments, inner + has 3 arguments
 - » Arguments are evaluated before the function

Basic Functions

Can build Lisp out of these functions

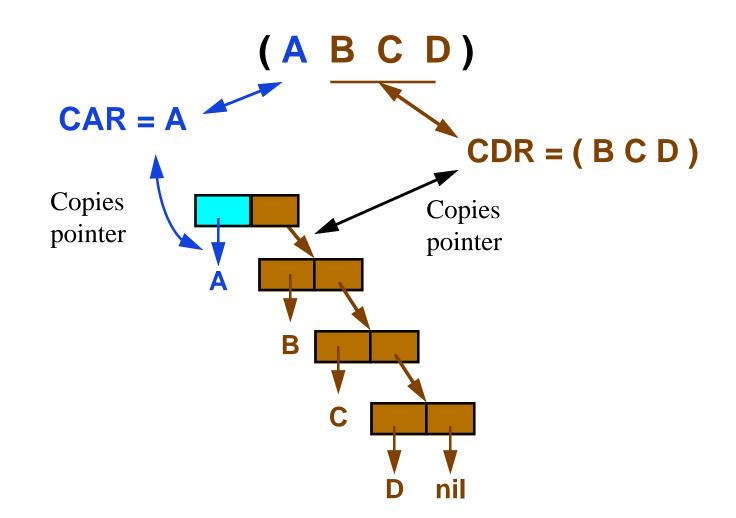
- List access & creation
 - » car or first access first in list
 - » cdr or rest access all but first
 - » cons construct a list cell
- Other
 - » quote or ' take literally, do not interpret
 - » atom true if argument is an atom
 - » eq true if arguments are same object
 - » cond conditional generalized "if ... then ... else"

List access functions

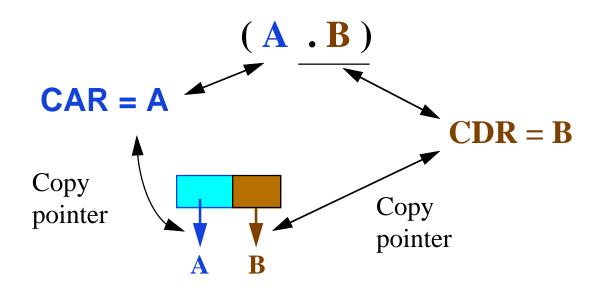


$$(\mathbf{cdr} (\mathbf{a} \mathbf{b} \mathbf{c})) \equiv (\mathbf{rest} (\mathbf{a} \mathbf{b} \mathbf{c})) \equiv (\mathbf{b} \mathbf{c})$$

CAR and CDR – Structural View 1



CAR and CDR – Structural View 2



(car'(abc)) – why the quote?

- Recall that arguments are evaluated before the function
- ♦ If we wrote (car (abc))
 - » argument (a b c) would be evaluated before the car
 - » a would be a function call
 - » but we literally want the list (a b c) not the result of evaluating a on the arguments b and c.
- (...) is syntactic sugar for (quote ...) where Lisp treats the function quote as a special function whose arguments are not evaluated first

 $(car '(a b c)) \equiv (car (quote (a b c)))$

Why the names CAR and CDR?

- Original Lisp developed for an IBM 704 computer which had 18 bit registers
- Pairs of registers could be handled as a single 36 bit 'word'

one word = one lisp cell

address register	decrement register

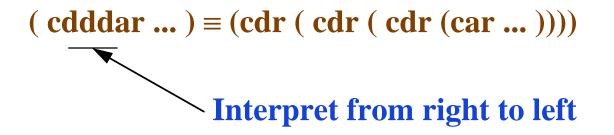
CAR = Contents Address Register

CDR = **Contents Decrement Register**

Short hand for nested car's and cdr's

Accessing deeper into Lisp structures occurs so frequently that additonal functions are introduced into Lisp.

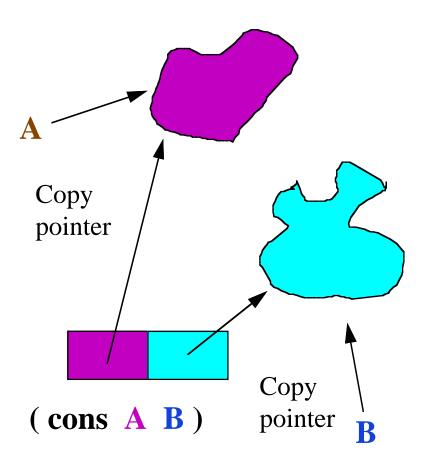
For example



Length depends upon the implementation.

Creating a New Lisp Cell – cons

- Only one constructor function
 cons
- Copies pointers to the arguments
- Laws
 - » (car (cons A B)) = A
 - » (cdr (cons A B)) = B



Destructive List Construction

Ons is expensive as it creates a new cell

- » memory allocation is invoked
 - > But it is non destructive no side effects

Following is dangerous – do not use in the course!

- For efficiency Common Lisp provides a set of destructive operations they change lists
 - » (replca cell newValue) & (replcd cell newValue)

> Replace the car and cdr fields of cell with ponters to newValue

» (nconc x y)

> Replace the cdr field of the last component of x with a pointer to y

SETQ – Define a symbol value

- (setq x value)
 - » If the symbol x does not exist it is created
 - » Symbol x is given the value value
- In this course USE ONLY AT THE GLOBAL LEVEL to create symbols required to test your programs
- ♦ Example
 - » (setq x '(1+ 4)) sets the value of x to the list (1+ 4)
 - » Note the x is not quoted but the second argument is if you do not want to evaluate it.

Compare SET and SETQ

◊ (setq x 'y)

» x has the value y

- ◊ (set x 'z)
 - » x still has the value y
 - » but a new symbol y is created with the value z
 » why?
- See the notes on symbols

DEFUN – define a function

(defun functionName (argumentList)

List of S-expressions to evaluate when the function is invoked – usually only one S-expression

– Example

(defun add (a b) (+ a b))

- Value of the function is the value of the last
 S-expression that is executed
- Functions in Lisp are typically small
 » rarely more than 1/2 a page in length