# Digital Logic Design 

## Chapter 1

Introduction

## Introduction

- This course is about Design Techniques for Digital System, a more exact name will be Synchronous Digital Hardware System.
- Synchronous means clocked i.e. all changes in the system are controlled by a global clock and happen at the same time
- Digital means all values (input, output, and internal) can take on discrete values.
- A/D if the input is analog (voice or music).
- Text: Digital Design, Mano and Ciletti $4^{\text {th }}$ ED Prentice Hall
- References:
- Digital Design: Principles and Practices, Wakerly, Prentice Hall
- Advanced Digital Design with the Verilog HDL, M. Ciletti, Prentice hall
- Contemporary Logic Design, Katz and Borriello, Prentice Hall


## Course contents

- Number system and how to represent things digitally.
- Boolean algebra and logic circuits.
- Combinational design
- Sequential design
- This is not a course on transistor physics or circuits, but we need to know these to better understand the building blocks of the system.
- Not a course on computer organization, but we will look at these as example of what we can do


## Good Design

- Digital systems are very complex and large, in order to do a good design, consider these issues:
- Modularize your design
- Top-down Design
- Bottom-up Design
- Design issues: Speed, Cost, Power
- Usually these are contradictory (a fast system is not cheap.
- Design is more of an art than science, but luckily we have measures for the design (cost, speed, .)


## Specifications

- Like in any other design, we start with the specification.
- The specification is basically what do we want to achieve.
- High level specification
- Binary level specification
- Algorithmic level specification



## CAD TOOLS

- Nowadays, design is usually automated, we use design tools for our design.
- HDL is used to describe the system (or the proposed system) in a high-level C-like language.
- Synthesis tools are used to map this design to FPGA
- Simulation tools are used to check the design (timing or functional simulation).
- Finally, testing


## Implementation

- Integrated Circuits (IC's)
- Crystalline silicon
- 1-100's of Millions of transistors
- Feature size 0.13 um or 0.1390 n (0.09micron) 60n and dropping ...
- CMOS (mostly)
- Standard microprocessors
- ASIC (application Specific IC's)
- FPGA's (Field Programmable Gate Arrays)
- Printed Ciruits Board (PCB)
- Fiberglass or ceramics
- Many conductive layers (1-20)
- Multiple Chips Modules (MCM's) multiple chips directly connected to a substarte
- Chassis


## The Design Process



## Number System

- Decimal numbers (9735) $=9 \times 10^{3}+7 \times$ $10^{2}+3 \times 10^{1}+510^{0}$
- Binary numbers (101) = 5 decimal

| Number | Quotient | Remainder |
| :--- | :--- | :--- |
| $29 / 2$ | 14 | 1 |
| $14 / 2$ | 7 | 0 |
| $7 / 2$ | 3 | 1 |
| $3 / 2$ | 1 | 1 |
| $1 / 2$ | 0 | 1 |

- How to convert 29 to binary (successive division by 2 the answer is 11101


## Number System

- Octal and Hex
- Conversion is the same idea


## Complements—Diminished radix

- Given a number $N$, in base $r$ having $n$ digits is defined as $\left(r^{n}-1\right)-N$
- For example $9^{\text {th }}$ complement of 456325 is 543674
- The 1's complement of any binary number is obtained by changing every 1 to 0 and every 0 to 1
- The 1's complement of 101100010 is 010011101


## Radix Complement

- Given a number $N$, in base $r$ having $n$ digits is defined as $r^{n}-N=\left(r^{n}-1\right)-N+1$
- For the $10^{\text {th }}$ complement, the rule is
- Leave the least significant 0's unchanged, the first nonzero digit is subtracted from 10, the rest of the digit are subtracted from $910^{\text {th }}$ complement of 3451600 is

6548400

- For 2's complement, the LSB zeros are left unchanged, the first 1 un changed, the remaining bits are complemented 010010


## Signed Binary Numbers

|  | Decimal | 2's complement |
| :---: | :---: | :---: |
|  | +7 | 0111 |
|  | +6 | 0110 |
|  | +5 | 0101 |
|  | +4 | 0100 |
|  | +3 | 0011 |
|  | +2 | 0010 |
|  | +1 | 0001 |
|  | +0 | 0000 |
|  | -0 |  |
|  | -1 | 1111 |
|  | -2 | 1110 |
|  | -3 | 1101 |
|  | -4 | 1100 |
|  | -5 | 1011 |
|  | -6 | 1010 |
| Fall 2008 | -7 | 1001 |
|  | -8 | 1000 |


| 1's complement | signed magnitude |  |
| :--- | :--- | :--- |
| 0111 | 0111 |  |
| 0110 | 0110 |  |
| 0101 | 0101 |  |
| 0100 | 0100 |  |
| 0011 | 0011 |  |
| 0010 | 0010 |  |
| 0001 | 0001 |  |
| 0000 | 0000 |  |
| 1111 | 1000 |  |
| 1110 | 1001 |  |
| 1101 | 1010 |  |
| 1100 | 1011 |  |
| 1011 | 1100 |  |
| 1010 | 1101 | 1110 |
| 1001 | 1111 |  |
| EeSe3201 |  |  |

## Subtraction with complement

- To subtract two n-digit unsigned numbers $M-N$ in base $r$ is done as follows

1. Add the minuend $M$ to the $r$ 's complement of the subtrahend $N$ yielding $M+\left(r^{n}-N\right)=M-N+r^{n}$
2. If $M \geq N$ The sum will produce an end carry, that is basically the $r^{n}$
3. If $M \leq N$ The sum does not produce a carry and is equal to $r^{n}-(N-M)$ which is the $r$ 's complement of $(N-M)$. To obtain the answer take the $r$ 's complement of the sum and place a -ve sign next to it.

## Addition (examples)

## Addition (Example)

## Binary Codes

- BCD Note that in adding 2 BCD numbers, the digits are added as if they are 2 binary numbers, if the result is greater than or equal 1010, we add 0110 to obtain the correct BCD digit sum and a carry
- Gray code (why do we care?)
- ASCII and ASCII with parity


## BCD Addition

## Gray Codes

- Only I digit change when we go from any number to number+1
- To form a sequence, put the sequence from left to right, followed by the sequence reversed. Add 0 as a MSB to the left sequence and 1 to the right sequence
$0 \quad 1$

| 00 | 01 | 11 | 10 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 000 | 001 | 011 | 010 | 110 | 111 | 101 | 100 |
| 0 | 1 | 3 | 2 | 6 | 7 | 5 | 4 |

## Register Transfer Logic

- Example R2 $\rightarrow$ R1
- What is an $R$ ?


## Binary Logic

- Low 0-1 volt
- High 1.2-4 volts
- In-between not defined
- AND, OR NOT (EXOR, NAND, NOR)

