

CSE4210

Architecture and Hardware for DSP

Lecture 4

Division and CORDIC

Division

- Division can be implemented as shift and subtract (shift and add for mul.)
- In division, the result of one subtraction determine the next operation (unlike multiplication).
- More difficult, and more time consuming than multiplication.

Subtract and Shift Division

- $X/D=Q \Rightarrow X = Q \cdot D + R, R < D$
- For simplicity, we assume there is no overflow (more bits to represent the quotient than the register size).
- Restoring and non-restoring division

Restoring Division

- Assume that Q contains the dividend (by the end of the calculation, it will contain the quotient).
- A Contains zero, by the end it will contain the remainder
- M contains the divisor

Restoring Division

- Algorithm
- Do n times (number of bits in Q)
 - Left shift A|Q
 - $A \leftarrow A - M$
 - If $A < 0$, $q_0 = 0$, $A \leftarrow A + M$ (restore)
 - Else $q_0 = 1$

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?
A=A-M	1101	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?
A=A-M	1101	
A > 0	0010	0001

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?
A=A-M	1101	
A > 0	0010	0001
Left shift A Q	0100	001?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?
A=A-M	1101	
A > 0	0010	0001
Left shift A Q	0100	001?
A=A-M	1101	

Restoring Division

	[M]=0011	[Q]=1010
	[A]=0000	
Left shift A Q	0001	010?
A=A-M	1101	
A < 0	1110	0100
A=A+M	0011	
	0001	0100
Left shift A Q	0010	100?
A=A-M	1101	
A < 0	1111	1000
A=A+M	0011	
	0010	1000
Left shift A Q	0101	000?
A=A-M	1101	
A > 0	0010	0001
Left shift A Q	0100	001?
A=A-M	1101	
A > 0	0001	0011

R=1 Q=3

Non Restoring Division

- Do n times {
 - Left shift A|Q
 - If previous $A \geq 0$ then $A = A \leftarrow A - M$
 - Else $A \leftarrow A + M$
 - If current $A \geq 0$ then $q_0 \leftarrow 1$
 - Else $q_0 \leftarrow 0$
 - }
- If $A < 0$ then $A \leftarrow A + M$

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev</i> +	1101	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev</i> +	1101	
A < 0	1110	0100

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M <i>prev -</i>	0011	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M <i>prev -</i>	0011	
A < 0	1111	1000

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M <i>prev -</i>	0011	
A < 0	1111	1000
Left shift A Q	1111	000?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M prev +	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M prev -	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M prev -	0011	

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M prev +	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M prev -	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M prev -	0011	
A > 0	0010	0001

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M prev +	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M prev -	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M prev -	0011	
A > 0	0010	0001
Left shift A Q	0100	001?

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M prev +	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M prev -	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M prev -	0011	
A > 0	0010	0001
Left shift A Q	0100	001?
A = A-M prev +	1101	

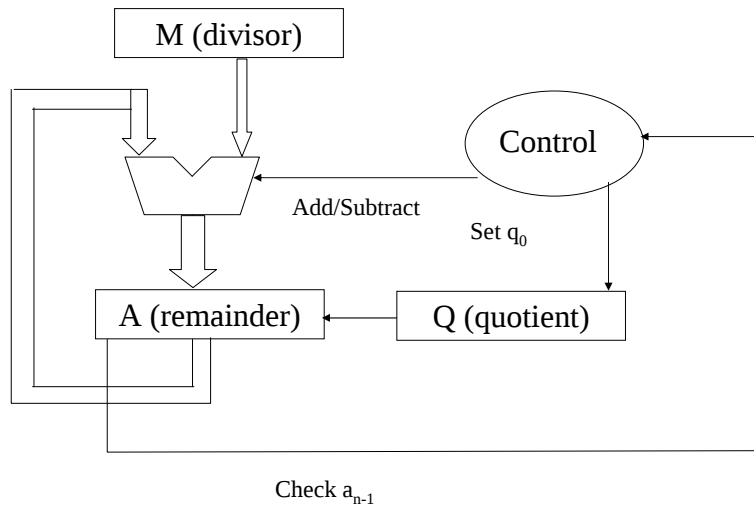
Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M <i>prev -</i>	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M <i>prev -</i>	0011	
A > 0	0010	0001
Left shift A Q	0100	001?
A = A-M <i>prev +</i>	1101	
A > 0	0001	0011

Restoring Division

	[M]=0011	
	[A]=0000	[Q]=1010
Left shift A Q	0001	010?
A=A-M <i>prev +</i>	1101	
A < 0	1110	0100
Left shift A Q	1100	100?
A=A+M	0011	
A < 0	1111	1000
Left shift A Q	1111	000?
A=A+M	0011	
A > 0	0010	0001
Left shift A Q	0100	001?
A = A-M	1101	
A > 0	0001	0011
A > 0 NO CORRECTION		

R
Q



CORDIC

- CORDIC (COordinate Rotation DIgital Computer)
- Used to calculate trigonometric functions (and many other)
- Reduces calculation to a series of shift and add operations
- Delay and hardware cost comparable to division.

CORDIC

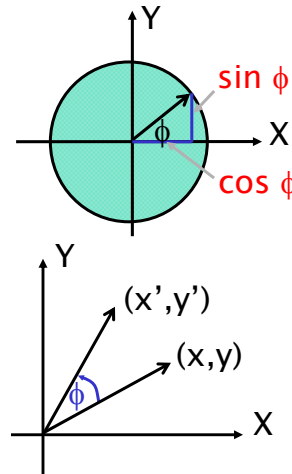
- The vector $(1,0)$ is rotated by an angle ϕ

$$x' = x \cdot \cos(\phi) - y \cdot \sin(\phi)$$

$$y' = y \cdot \cos(\phi) + x \cdot \sin(\phi)$$

$$x' = \cos(\phi)[x - y \tan(\phi)]$$

$$y' = \cos(\phi)[x + y \tan(\phi)]$$



CORDIC

$$x' = \cos(\phi)[x - y \tan(\phi)]$$

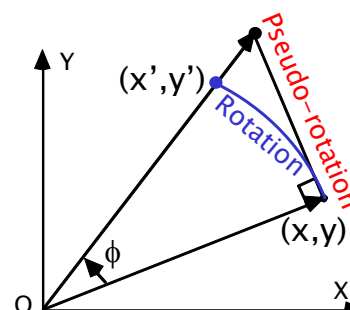
$$y' = \cos(\phi)[x + y \tan(\phi)]$$

If we do not multiply by $\cos(\phi) \rightarrow$
pseudo rotation

If we start by $(1,0)$ then x' is $\cos(\phi)$

If we choose $\tan(\phi)$ to be 2^{-1} , then
the multiplication is a shift.

Any angle can be the sum of many ϕ 's
with $\tan = 2^{-i}$



CORDIC

i	α_i	$\tan(\alpha_i)$	
0	45.0°	1	$= 2^{-0}$
1	26.6°	0.5	$= 2^{-1}$
2	14.0°	0.25	$= 2^{-2}$
3	7.1°	0.125	$= 2^{-3}$
4	3.6°	0.0625	$= 2^{-4}$
5	1.8°	0.03125	$= 2^{-5}$
6	0.9°	0.015625	$= 2^{-6}$
7	0.4°	0.0078125	$= 2^{-7}$
8	0.2°	0.00390625	$= 2^{-8}$
9	0.1°	0.001953125	$= 2^{-9}$
10	0.055	0.000976562	$= 2^{-10}$

$$\bullet \quad 37.2^\circ = 45 - 26.6 + 14.0 + 7.1 - 3.6 + 1.8 - 0.9 + 0.4 \\ = 37.2^\circ$$

CORDIC

$$x_{i+1} = \cos(\phi_i) \cdot [x_i - y_i \cdot \sigma_i \cdot \tan(\phi_i)] \quad \sigma_i \in \{-1, 1\}$$

$$y_{i+1} = \cos(\phi_i) \cdot [y_i + x_i \cdot \sigma_i \cdot \tan(\phi_i)]$$

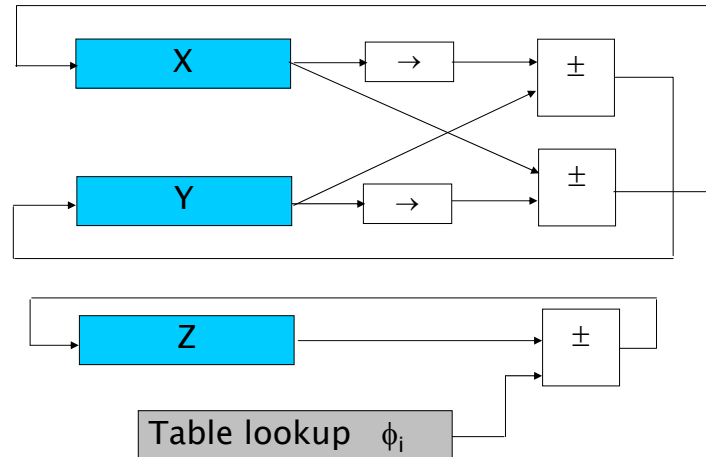
$$x_{i+1} = \cos(\phi_i) \cdot [x_i - \sigma_i \cdot 2^{-i} \cdot y_i]$$

$$y_{i+1} = \cos(\phi_i) \cdot [y_i + \sigma_i \cdot 2^{-i} \cdot x_i]$$

$$z_{i+1} = z_i - \sigma_i \tan^{-1}(2^{-i})$$

$\cos(\phi_i)$ is a constant K_i

CORDIC -- Hardware



CORDIC

$$x_{i+1} = K_i \cdot [x_i - \sigma_i \cdot 2^{-i} \cdot y_i]$$

$$y_{i+1} = K_i \cdot [y_i + \sigma_i \cdot 2^{-i} \cdot x_i]$$

We can multiply by a constant K once at the end

$$K = \prod_{i=1}^n K_i \quad n \rightarrow \infty, \quad K = 0.607252935\dots$$

Or, we can start with (0.60725, 0) instead of (1, 0) and save one multiplication