
COSC4201

Loop Level Parallelism

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Loop Level Parallelism LLP

- Loop-Level Parallelism (LLP) analysis focuses on whether data accesses in later iterations of a loop are data dependent on data values produced in earlier iterations and possibly making loop iterations independent.

- e.g. in **for (i=1; i<=1000; i++)**
x[i] = x[i] + s;

the computation in each iteration is independent of the previous iterations and the loop is thus parallel. The use of **x[i]** twice is within a single iteration.

⇒ Thus loop iterations are parallel (or independent from each other).

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Loop Level Parallelism LLP

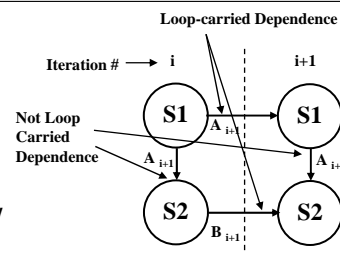
- **Loop-carried Dependence:** A data dependence between different loop iterations (data produced in earlier iteration used in a later one).
- LLP analysis is important in software optimizations such as loop unrolling since it usually requires loop iterations to be independent.
- LLP analysis is normally done at the source code level or close to it since assembly language and target machine code generation introduces loop-carried name dependence in the registers used for addressing and incrementing.
- Instruction level parallelism (ILP) analysis, on the other hand, is usually done when instructions are generated by the compiler

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Loop Level Parallelism LLP

```
for (i=1; i<=100; i=i+1) {
    A[i+1] = A[i] + C[i]; /* S1 */
    B[i+1] = B[i] + A[i+1]; /* S2 */
}
```



Dependency Graph

- S2 uses the value $A[i+1]$, computed by S1 in the same iteration. This data dependence is within the same iteration (not a loop-carried dependence).
 \Rightarrow does not prevent loop iteration parallelism.
- S1 uses a value computed by S1 in an earlier iteration, since iteration i computes $A[i+1]$ read in iteration $i+1$ (loop-carried dependence, prevents parallelism). The same applies for S2 for $B[i]$ and $B[i+1]$
 \Rightarrow These two dependencies are loop-carried spanning more than one iteration preventing loop parallelism.

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```
for (i=1; i<=100; i=i+1) {
```

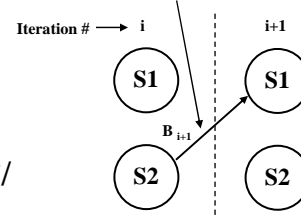
```
    A[i] = A[i] + B[i];      /* S1 */
```

```
    B[i+1] = C[i] + D[i];   /* S2 */
```

```
}
```

- S1 uses the value B[i] computed by S2 in the previous iteration (loop-carried dependence)
- This dependence is not circular:
 - S1 depends on S2 but S2 does not depend on S1.

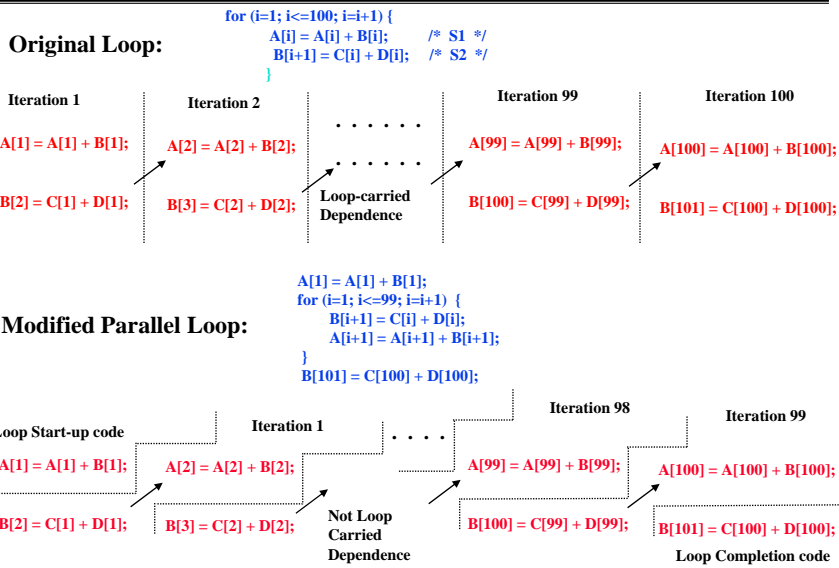
Dependency Graph Loop-carried Dependence



LLP Analysis Example 2

```
A[1] = A[1] + B[1];
for (i=1; i<=99; i=i+1) {
    B[i+1] = C[i] + D[i];
    A[i+1] = A[i+1] + B[i+1];
}
B[101] = C[100] + D[100];
```

LLP Analysis Example 2



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LLP

```

for(i=2; i<=100; i++) {
    y[i]=y[i-1]+y[i]
}

```

```

for(i=2; i<=100; i++) {
    y[i]=y[i-5]+y[i]
}

```

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Finding Dependences

- Finding dependences in the program is very important for renaming and executing instructions in parallel.
- Arrays and pointers makes finding dependences very difficult.
- Assume array indices are *affine*, which means on the form $a \times i + b$ where a and b are constant.
- GCD test can be used to detect dependences.

GCD Test

- Assume we stored an array with index value of $a \times i + b$ and loaded an array with an index value of $c \times i + d$
- Are they pointing to the same location?
- Assume the loop limit is m, n
- Are there

$$j, k \quad m \leq j, k \leq n \text{ such that } a \times j + b = c \times k + d$$

GCD Test

- ° A simple and **sufficient** test for absence can be found.
- ° If a loop dependence exists, then

$GCD(c, a)$ must divides $(d - b)$

- ° If that test fails, there is no guarantee there is dependence (loop bound)

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GCD Test

```
for(i=1; i<=100; i=i+1) {  
    x[2*i+3] = x[2*i] * 5.0;  
}
```

$a = 2$ $b = 3$ $c = 2$ $d = 0$

$GCD(a, c) = 2$

$d - b = -3$

2 does not divide -3 \Rightarrow No
dependence is not possible.

5,7,9,11,13,15,17,19,21,23,....

4,6,8,10,12,14,16,18,20,22,....

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Dependence Analysis -- Difficulties

- Dependence analysis is a very important tool for exploiting LLP, it can not be used in these situations
- Objects are referenced using pointers
- Array indexing using another array $A[b[I]]$
- Dependence may exist for some values of input, but in reality the input never takes these values.
- When we want to more than the possibility of dependence (which write causes it?)
- Dependence analysis across procedure boundaries

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Dependence Analysis -- Difficulties

- Sometimes, *points-to* analysis might help.
- We might be able to answer *simpler* questions, or get some hints.
- Do 2 pointers point to the same list?
- Type information
- Information derived when the object was allocated
- Pointer assignments

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