## Link Analysis: PageRank and HITS

Thanks to Jure Leskovec, Stanford and Panayiotis Tsaparas, Univ. of loannina for slides

## Agenda

- Web Search: How to Organize the Web?
- Ranking Nodes on Graphs
- Hubs and Authorities
- PageRank
- How to Solve PageRank
- Personalized PageRank
- Link Prediction in Networks
- Estimating Scores for Missing Edges
- Case studies:
- Facebook: Supervised Random Walks for Link Prediction
- Twitter: The who to follow service at Twitter


## How to Organize the Web?

- How to organize the Web?
- First try: Human curated Web directories
- Yahoo, DMOZ, LookSmart
- Second try: Web Search
- Information Retrieval attempts to find relevant docs in a small and trusted set
- Newspaper articles, Patents, etc.

- But: Web is huge, full of untrusted documents, random things, web spam, etc.
- So we need a good way to rank webpages!


## Web Search: 2 Challenges

2 challenges of web search:

- (1) Web contains many sources of information Who to "trust"?
" Insight: Trustworthy pages may point to each other!
- (2) What is the "best" answer to query "newspaper"?
- No single right answer
- Insight: Pages that actually know about newspapers might all be pointing to many newspapers


## Ranking Nodes on the Graph

- All web pages are not equally "important" www.joe-schmoe.com vs. www.stanford.edu
- We already know:

There is large diversity in the web-graph node connectivity.

- So, let's rank the pages using the web graph link structure!



## Link Analysis Algorithms

- We will cover the following Link Analysis approaches to computing importance of nodes in a graph:
- Hubs and Authorities (HITS)
- Page Rank
- Topic-Specific (Personalized) Page Rank

Sidenote: Various notions of node centrality: Node $\boldsymbol{u}$

- Degree centrality = degree of $u$
- Betweenness centrality = \#shortest paths passing through $u$
- Closeness centrality = avg. length of shortest paths from $u$ to all other nodes of the network
- Eigenvector centrality = like PageRank


## Hubs and Authorities

## Link Analysis

- Goal (back to the newspaper example):
- Don't just find newspapers. Find "experts" - pages that link in a coordinated way to good newspapers
- Idea: Links as votes
- Page is more important if it has more links
- In-coming links? Out-going links?
- Hubs and Authorities

Each page has 2 scores:

- Quality as an expert (hub):
- Total sum of votes of pages pointed to
- Quality as a content (authority):
- Total sum of votes of experts

- Principle of repeated improvement


## Hubs and Authorities

Interesting pages fall into two classes:

1. Authorities are pages containing useful information

- Newspaper home pages
- Course home pages
- Home pages of auto manufacturers


2. Hubs are pages that link to authorities

- List of newspapers
- Course bulletin
- List of U.S. auto manufacturers



## Counting in-links: Authority



Each page starts with hub score 1 Authorities collect their votes
(Note this is idealized example. In reality graph is not bipartite and each page has both the hub and authority score)

## Expert Quality: Hub


(Note this is idealized example. In reality graph is not bipartite and each page has both the hub and authority score)

## Reweighting



## Mutually Recursive Definition

- A good hub links to many good authorities
- A good authority is linked from many good hubs
- Note a self-reinforcing recursive definition
- Model using two scores for each node:
- Hub score and Authority score
- Represented as vectors $\boldsymbol{h}$ and $\boldsymbol{a}$, where the $\boldsymbol{i}$-th element is the hub/authority score of the $i$-th node


## Hubs and Authorities

- Each page $i$ has 2 scores:
- Authority score: $\boldsymbol{a}_{\boldsymbol{i}}$
- Hub score: $\boldsymbol{h}_{\boldsymbol{i}}$

Convergence criteria:

$$
\begin{aligned}
& \sum_{i}\left(h_{i}^{(t)}-h_{i}^{(t+1)}\right)^{2}<\varepsilon \\
& \sum_{i}\left(a_{i}^{(t)}-a_{i}^{(t+1)}\right)^{2}<\varepsilon
\end{aligned}
$$

HITS algorithm:

- Initialize: $a_{j}^{(0)}=1 / \sqrt{\mathrm{n}}, \mathrm{h}_{\mathrm{j}}^{(0)}=1 / \sqrt{\mathrm{n}}$
- Then keep iterating until convergence:
- $\forall i$ : Authority: $a_{i}^{(t+1)}=\sum_{j \rightarrow i} h_{j}^{(t)}$
- $\forall \boldsymbol{i}$ : Hub: $h_{i}^{(t+1)}=\sum_{i \rightarrow j} a_{j}^{(t)}$
- $\forall \boldsymbol{i}$ : Normalize:

$$
\sum_{i}\left(a_{i}^{(t+1)}\right)^{2}=1, \sum_{j}\left(h_{j}^{(t+1)}\right)^{2}=1
$$

## Hubs and Authorities

- Definition: Eigenvectors \& Eigenvalues
- Let $\boldsymbol{R} \cdot \boldsymbol{x}=\boldsymbol{\lambda} \cdot \boldsymbol{x}$
for some scalar $\boldsymbol{\lambda}$, vector $\boldsymbol{x}$, matrix $\boldsymbol{R}$
- Then $x$ is an eigenvector, and $\lambda$ is its eigenvalue
- The steady state (HITS has converged) is:
${ }^{-} \boldsymbol{A}^{\boldsymbol{T}} \cdot \boldsymbol{A} \cdot \boldsymbol{a}=\boldsymbol{c}^{\prime} \cdot \boldsymbol{a}$
- $\boldsymbol{A} \cdot \boldsymbol{A}^{\boldsymbol{T}} \cdot \boldsymbol{h}=\boldsymbol{c}^{\prime \prime} \cdot \boldsymbol{h}$

Note constants c', c' don't matter as we normalize them out every step of HITS

- So, authority $a$ is eigenvector of $\boldsymbol{A}^{T} A$ (associated with the largest eigenvalue) Similarly: hub $\boldsymbol{h}$ is eigenvector of $\boldsymbol{A} \boldsymbol{A}^{\boldsymbol{T}}$


## PageRank

## Links as Votes

- Still the same idea: Links as votes
- Page is more important if it has more links
- In-coming links? Out-going links?
- Think of in-links as votes:
- www.stanford.edu (many in-links)
- www.edessacity.gr (few in-link)
- Are all in-links equal?
- Links from important pages count more
- Recursive question!


## PageRank: The "Flow" Model

- A "vote" from an important page is worth more:
- Each link's vote is proportional to the importance of its source page
- If page $i$ with importance $r_{i}$ has $d_{i}$ out-links, each link gets $r_{i} / d_{i}$ votes

- Page j's own importance $r_{j}$ is the sum of the votes on its inlinks


## PageRank: The "Flow" Model

- A page is important if it is pointed to by other important pages
- Define a "rank" $r_{j}$ for node $j$

$$
r_{j}=\sum_{i \rightarrow j} \frac{r_{i}}{\mathrm{~d}_{\mathrm{i}}}
$$

$d_{i} \ldots$ out-degree of node $i$

You might wonder: Let's just use Gaussian elimination to solve this system of linear equations. Bad idea!

The web in 1839

"Flow" equations:

$$
\begin{aligned}
& \mathbf{r}_{\mathrm{y}}=\mathbf{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{a}} / 2 \\
& \mathbf{r}_{\mathrm{a}}=\mathbf{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{m}} \\
& \mathbf{r}_{\mathrm{m}}=\mathbf{r}_{\mathrm{a}} / 2
\end{aligned}
$$

## PageRank: Matrix Formulation

- Stochastic adjacency matrix M
- Let page $\boldsymbol{j}$ have $\boldsymbol{d}_{\boldsymbol{j}}$ out-links
- If $\boldsymbol{j} \rightarrow \boldsymbol{i}$, then $\boldsymbol{M}_{\boldsymbol{i} \boldsymbol{j}}=\frac{\mathbf{1}}{\boldsymbol{d}_{\boldsymbol{i}}}$
- $\boldsymbol{M}$ is a column stochastic matrix
- Columns sum to 1
- Rank vector $r$ : An entry per page

- $\boldsymbol{r}_{\boldsymbol{i}}$ is the importance score of page $\boldsymbol{i}$
$-\sum_{i} r_{i}=1$
- The flow equations can be written

$$
\boldsymbol{r}=\boldsymbol{M} \cdot \boldsymbol{r} \quad r_{j}=\sum_{i \rightarrow j} \frac{r_{i}}{\mathrm{~d}_{i}}
$$

## Random Walk Interpretation

- Imagine a random web surfer:
- At any time $\boldsymbol{t}$, surfer is on some page $i$
- At time $\boldsymbol{t}+\mathbf{1}$, the surfer follows an out-link from $i$ uniformly at random
- Ends up on some page $\boldsymbol{j}$ linked from $\boldsymbol{i}$

- Process repeats indefinitely
- Let:
- $\boldsymbol{p}(\boldsymbol{t})$... vector whose $i^{\text {th }}$ coordinate is the prob. that the surfer is at page $i$ at time $t$
- So, $\boldsymbol{p}(\boldsymbol{t})$ is a probability distribution over pages


## The Stationary Distribution

- Where is the surfer at time $t+1$ ?
- Follows a link uniformly at random

$$
p(t+1)=M \cdot p(t)
$$

- Suppose the random walk reaches a state

$$
p(t+1)=M \cdot p(t)=p(t)
$$

then $\boldsymbol{p}(t)$ is stationary distribution of a random walk

- Our original rank vector $\boldsymbol{r}$ satisfies $\boldsymbol{r}=\boldsymbol{M} \cdot \boldsymbol{r}$
- So, $r$ is a stationary distribution for the random walk


# PageRank: How to solve? 

## PageRank: How to solve?

Given a web graph with $n$ nodes, where the nodes are pages and edges are hyperlinks

- Assign each node an initial page rank
- Repeat until convergence $\left(\Sigma_{i}\left|r_{i}^{(t+1)}-r_{i}^{(t)}\right|<\varepsilon\right)$
- Calculate the page rank of each node

$$
r_{j}^{(t+1)}=\sum_{i \rightarrow j} \frac{r_{i}^{(t)}}{\mathrm{d}_{\mathrm{i}}}
$$

$d_{i} \ldots$ out-degree of node $\boldsymbol{i}$

## PageRank: How to solve?

- Power Iteration:
- Set $r_{j} \leftarrow 1 / N$
- 1: $r^{\prime}{ }_{j} \leftarrow \sum_{i \rightarrow j} \frac{r_{i}}{d_{i}}$
- 2: $r \leftarrow r^{\prime}$
- If $\left|r-r^{\prime}\right|>\varepsilon$ : goto 1


|  | y | a | m |
| :---: | :---: | :---: | :---: |
| y | $1 / 2$ | $1 / 2$ | 0 |
| a | $1 / 2$ | 0 | 1 |
| m | 0 | $1 / 2$ | 0 |
|  |  |  |  |

$$
\begin{aligned}
& \mathbf{r}_{\mathrm{y}}=\mathrm{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{a}} / 2 \\
& \mathbf{r}_{\mathrm{a}}=\mathbf{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{m}} \\
& \mathbf{r}_{\mathrm{m}}=\mathbf{r}_{\mathrm{a}} / 2
\end{aligned}
$$

- Example:
\(\left(\begin{array}{l}r_{y} <br>
r_{\mathrm{a}} <br>

\mathrm{r}_{\mathrm{m}}\end{array}\right)=\)| $1 / 3$ |
| :--- |
| $1 / 3$ |
| $1 / 3$ |

Iteration 0, 1, 2, ...

## PageRank: How to solve?

- Power Iteration:
- Set $r_{j} \leftarrow 1 / N$
- 1: $r^{\prime}{ }_{j} \leftarrow \sum_{i \rightarrow j} \frac{r_{i}}{d_{i}}$
- 2: $r \leftarrow r^{\prime}$
- If $\left|r-r^{\prime}\right|>\varepsilon$ : goto 1
- Example:
\(\left(\begin{array}{l}r_{y} <br>
r_{\mathrm{a}} <br>

\mathrm{r}_{\mathrm{m}}\end{array}\right)=\)| $1 / 3$ | $1 / 3$ | $5 / 12$ | $9 / 24$ |  | $6 / 15$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $1 / 3$ | $3 / 6$ | $1 / 3$ | $11 / 24$ | $\ldots$ | $6 / 15$ |
| $1 / 3$ | $1 / 6$ | $3 / 12$ | $1 / 6$ |  | $3 / 15$ |

Iteration 0, 1, 2, ...


6/15
6/15
3/15

|  | y |  | a |
| ---: | :---: | :---: | :---: |
| m |  |  |  |
| y | $1 / 2$ | $1 / 2$ | 0 |
| a | $1 / 2$ | 0 | 1 |
| m | 0 | $1 / 2$ | 0 |
|  |  |  |  |

$$
\begin{aligned}
& \mathbf{r}_{\mathrm{y}}=\mathbf{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{a}} / 2 \\
& \mathbf{r}_{\mathrm{a}}=\mathbf{r}_{\mathrm{y}} / 2+\mathbf{r}_{\mathrm{m}} \\
& \mathbf{r}_{\mathrm{m}}=\mathbf{r}_{\mathrm{a}} / 2
\end{aligned}
$$

## PageRank: Three Questions

$$
r_{j}^{(t+1)}=\sum_{i \rightarrow j} \frac{r_{i}^{(t)}}{\mathrm{d}_{\mathrm{i}}} \underset{\text { equivalently }}{\text { or }} \quad r=M \boldsymbol{r}
$$

- Does this converge?
- Does it converge to what we want?
- Are results reasonable?


## Does this converge?

The "Spider trap" problem:


$$
r_{j}^{(t+1)}=\sum_{i \rightarrow j} \frac{r_{i}^{(t)}}{\mathrm{d}_{\mathrm{i}}}
$$

- Example:

Iteration: 0, 1, 2, 3...

| $\mathrm{r}_{\mathrm{a}}$ |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{r}_{\mathrm{b}}$ |$=$| 1 | 0 | 1 |
| :--- | :--- | :--- |
| 0 | 1 | 0 |
|  | 1 |  |

## Does it converge to what we want?

- The "Dead end" problem:

$$
\text { (a) } \longrightarrow r_{j}^{(t+1)}=\sum_{i \rightarrow j} \frac{r_{i}^{(t)}}{\mathrm{d}_{\mathrm{i}}}
$$

- Example:

Iteration: 0, 1, 2, 3...

| $\mathrm{r}_{\mathrm{a}}$ |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{r}_{\mathrm{b}}$ |$=$| 1 | 0 | 0 |
| :--- | :--- | :--- |
| 0 | 1 | 0 |
|  | 0 |  |

## RageRank: Problems

2 problems:

- (1) Some pages are dead ends (have no out-links)
- Such pages cause importance to "leak out"

- (2) Spider traps
(all out-links are within the group)
- Eventually spider traps absorb all importance


## Problem: Spider Traps

- Power Iteration:
- Set $r_{j}=\frac{1}{N}$
- $r_{j}=\sum_{i \rightarrow j} \frac{r_{i}}{d_{i}}$
- And iterate


$$
\begin{aligned}
& r_{y}=r_{y} / 2+r_{a} / 2 \\
& \mathbf{r}_{a}=r_{y} / 2 \\
& r_{m}=r_{a} / 2+r_{m}
\end{aligned}
$$

- Example:

| $\mathrm{r}_{\mathrm{y}}$ ) |  | 1/3 | 2/6 | 3/12 | 5/24 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{r}_{\mathrm{a}}$ | $=$ | 1/3 | 1/6 | 2/12 | 3/24 | 0 |
| $\mathrm{r}_{\mathrm{m}}$ |  | 1/3 | 3/6 | 7/12 | 16/24 | 1 |

Iteration 0, 1, 2, $\ldots$

## Solution: Random Teleports

- The Google solution for spider traps: At each time step, the random surfer has two options
- With prob. $\beta$, follow a link at random
- With prob. 1- $\beta$, jump to a random page
- Common values for $\beta$ are in the range 0.8 to 0.9
- Surfer will teleport out of spider trap within a few time steps



## Problem: Dead Ends

- Power Iteration:
- Set $r_{j}=\frac{1}{N}$
- $r_{j}=\sum_{i \rightarrow j} \frac{r_{i}}{d_{i}}$
- And iterate


|  | y | a | m |
| :---: | :---: | :---: | :---: |
| y | 1/2 | 1/2 | 0 |
| a | 1/2 | 0 | 0 |
| m | 0 | 1/2 | 0 |

$$
\begin{aligned}
& r_{y}=r_{y} / 2+r_{a} / 2 \\
& r_{a}=r_{y} / 2 \\
& r_{m}=r_{a} / 2
\end{aligned}
$$

- Example:

| $\mathrm{r}_{\mathrm{y}}$ ) |  | 1/3 | 2/6 | 3/12 | 5/24 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{r}_{\mathrm{a}}$ | $=$ | 1/3 | 1/6 | 2/12 | 3/24 | 0 |
| $\mathrm{r}_{\mathrm{m}}$ |  | 1/3 | 1/6 | 1/12 | 2/24 | 0 |

Iteration 0, 1, 2, ...

## Solution: Always Teleport

- Teleports: Follow random teleport links with probability 1.0 from dead-ends
- Adjust matrix accordingly


|  | y | a | m |
| ---: | :---: | :---: | :---: |
| y | $1 / 2$ | $1 / 2$ | 0 |
| a | $1 / 2$ | 0 | 0 |
| m | 0 | $1 / 2$ | 0 |
|  |  |  |  |


|  | y | a | m |
| :---: | :---: | :---: | :---: |
| y | $1 / 2$ | $1 / 2$ | $1 / 3$ |
| y | $1 / 2$ | 0 | $1 / 3$ |
| m | 0 | $1 / 2$ | $1 / 3$ |
|  |  |  |  |

## Final PageRank Equation

- Google's solution: At each step, random surfer has two options:
- With probability $\boldsymbol{\beta}$, follow a link at random
- With probability $\mathbf{1 - \beta}$, jump to some random page
- PageRank equation [Brin-Page, '98]

$$
r_{j}=\sum_{i \rightarrow j} \beta \frac{r_{i}}{d_{i}}+(1-\beta) \frac{1}{n}
$$

$d_{i} \ldots$ out-degree of node i

The above formulation assumes that $\boldsymbol{M}$ has no dead ends. We can either preprocess matrix $\boldsymbol{M}$ (bad!) or explicitly follow random teleport links with probability 1.0 from dead-ends. See P. Berkhin, A Survey on PageRank Computing, Internet Mathematics, 2005.

## PageRank \& Eigenvectors

- PageRank as a principal eigenvector $\boldsymbol{r}=\boldsymbol{M} \cdot \boldsymbol{r}$ or equivalently $\boldsymbol{r}_{\boldsymbol{j}}=\sum_{i} \frac{\boldsymbol{r}_{i}}{\boldsymbol{d}_{\boldsymbol{i}}}$
- But we really want ( ${ }^{* *}$ ):

$$
r_{j}=\beta \sum_{i} \frac{r_{i}}{d_{i}}+(1-\beta) \frac{1}{n}
$$

$d_{i}$... out-degree
of node i

- Let's define:

$$
M_{i j}^{\prime}=\beta M_{i j}+(1-\beta) \frac{1}{n}
$$

- Now we get what we want:

$$
\boldsymbol{r}=\boldsymbol{M}^{\prime} \cdot \boldsymbol{r}
$$

Note: $M$ is a sparse matrix but $\boldsymbol{M}^{\prime}$ is dense (all entries $\neq 0$ ). In practice we never "materialize" $M$ but rather we use the "sum" formulation (**)

- What is $\mathbf{1}-\boldsymbol{\beta}$ ?
- In practice 0.15 (Jump approx. every 5-6 links)


## The PageRank Algorithm

- Input: Graph $\boldsymbol{G}$ and parameter $\boldsymbol{\beta}$
- Directed graph $\boldsymbol{G}$ with spider traps and dead ends
- Parameter $\beta$
- Output: PageRank vector $r$
- Set: $r_{j}^{(0)}=\frac{1}{N}, \quad t=1$
- do:
$\forall j: \boldsymbol{r}_{j}^{\prime(t)}=\sum_{i \rightarrow j} \beta \frac{r_{i}^{(t-1)}}{d_{i}}$

$$
\boldsymbol{r}_{\boldsymbol{j}}^{\prime(t)}=\mathbf{0} \text { if in-deg. of } \boldsymbol{j} \text { is } \mathbf{0}
$$

- Now re-insert the leaked PageRank:

$$
\forall j: r_{j}^{(t)}=r_{j}^{\prime(t)}+\frac{1-S}{N} \quad \text { where: } S=\sum_{j} r_{j}^{(t)}
$$

- $t=t+1$
- while $\sum_{j}\left|r_{j}^{(t)}-r_{j}^{(t-1)}\right|>\varepsilon$


## Example



## PageRank and HITS

- PageRank and HITS are two solutions to the same problem:
- What is the value of an in-link from $u$ to $v$ ?
- In the PageRank model, the value of the link depends on the links into $u$
- In the HITS model, it depends on the value of the other links out of $u$
- The destinies of PageRank and HITS post-1998 were very different


# Personalized PageRank, Random Walk with Restarts 


a.k.a.: Relevance, Closeness, 'Similarity'...

## Example Application: Graph Search

- Given:

Conferences-to-authors graph

- Goal:

Proximity on graphs

- Q: What is most related conference to ICDM?
IJCAI
[Tong et al. ‘08]
Automatic Image Captioning

\{Sea Sun Sky Wave\}

\{Cat Forest Grass Tiger\}

\{?, ?, ?,\}
[Tong et al. ‘08]

[Tong et al. ‘08]



## Good proximity measure?

- Shortest path is not good:

- No influence for degree-1 nodes (E, F, G)!
- Multi-faceted relationships


## Good proximity measure?

- Network Flow is not good:

- Does not punish long paths


## What is good notion of proximity?



- Multiple Connections
- Quality of connection
-Direct \& In-direct connections
-Length, Degree,
Weight...


## Random Walk with Restarts



## Personalized PageRank

- Goal: Evaluate pages not just by popularity but by how close they are to the topic
- Teleporting can go to:
- Any page with equal probability
- (we used this so far)
- A topic-specific set of "relevant" pages
" Topic-specific (personalized) PageRank ( $S$...teleport set)

$$
\begin{aligned}
M_{i j}^{\prime} & =\beta M_{i j}+(1-\beta) /|S| \quad \text { if } i \in S \\
& =\beta M_{i j} \quad \text { otherwise }
\end{aligned}
$$

- Random Walk with Restart: $\mathbf{S}$ is a single element


## PageRank: Applications

- Graphs and web search:
- Ranks nodes by "importance"
- Personalized PageRank:
- Ranks proximity of nodes to the teleport nodes $\boldsymbol{S}$
- Proximity on graphs:
- Q: What is most related conference to ICDM?



## Random Walk with Restarts



Nearby nodes, higher scores
More red, more relevant

|  | Node 4 |
| :--- | :--- |
| Node 1 | 0.13 |
| Node 2 | 0.10 |
| Node 3 | 0.13 |
| Node 4 | 0.22 |
| Node 5 | 0.13 |
| Node 6 | 0.05 |
| Node 7 | 0.05 |
| Node 8 | 0.08 |
| Node 9 | 0.04 |
| Node 10 | 0.03 |
| Node 11 | 0.04 |
| Node 12 | 0.02 |

Ranking vector
$\vec{r}_{4}$

## Most related conferences to ICDM



DMKD

## Personalized PageRank

Q: Which conferences are closest to KDD \& ICDM?

A: Personalized PageRank with teleport set $S=\{K D D$, ICDM

## Link Prediction in Networks

Thanks to Jure Leskovec, Stanford and Panayiotis Tsaparas, Univ. of loannina for slides

## Agenda

- Link Prediction in Networks
- Estimating Scores for Missing Edges
- Classification Approach (Omitted)
- Case studies:
- Facebook: Supervised Random Walks for Link Prediction
- Twitter: The who to follow service at Twitter


## Link Prediction Motivation

- Recommending new friends in online social networks
- Predicting the participation of actors in events
- Suggesting interactions between the members of a company/organization that are external to the hierarchical structure of the organization itself
- Predicting connections between members of communities/organizations who have not been directly observed together
- Suggesting collaborations between researchers based on coauthorship
- Overcoming the data-sparsity problem in recommender systems using collaborative filtering
- The link prediction task:
- Given $G\left[t_{0}, t_{0}^{\prime}\right]$ a graph on edges up to time $t_{0}^{\prime}$, output a ranked list $L$ of links (not in $G\left[t_{0}, t_{0}^{\prime}\right]$ ) that are predicted to appear in $G\left[t_{1}, t_{1}^{\prime}\right]$
- Evaluation:

$G\left[t_{0}, t_{0}^{\prime}\right]$
$G\left[t_{1}, t_{1}^{\prime}\right]$
- $n=\left|E_{\text {new }}\right|$ : \# new edges that appear during the test period $\left[t_{1}, t_{1}^{\prime}\right]$
- Take top $n$ elements of $L$ and count correct edges


## Link Prediction

- Predict links in a evolving collaboration network

|  | training period $^{c \mid}$ |  |  | Core |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
|  | authors | papers | collaborations ${ }^{1}$ | authors | $\left\|E_{\text {old }}\right\|$ | $\left\|E_{\text {new }}\right\|$ |
| astro-ph | 5343 | 5816 | 41852 | 1561 | 6178 | 5751 |
| cond-mat | 5469 | 6700 | 19881 | 1253 | 1899 | 1150 |
| gr-qc | 2122 | 3287 | 5724 | 486 | 519 | 400 |
| hep-ph | 5414 | 10254 | 47806 | 1790 | 6654 | 3294 |
| hep-th | 5241 | 9498 | 15842 | 1438 | 2311 | 1576 |

- Core: Because network data is very sparse
- Consider only nodes with degree of at least 3
- Because we don't know enough about these nodes to make good inferences


## Link Prediction

- Methodology:
- For each pair of nodes ( $x, y$ ) compute score $c(x, y)$
- For example, $c(x, y)$ could be the \# of common neighbors of $x$ and $y$
- Sort pairs $(x, y)$ by the decreasing score $c(x, y)$
- Note: Only consider/predict edges where both endpoints are in the core (deg. $\geq 3$ )
- Predict top $n$ pairs as new links
- See which of these links actually appear in $G\left[t_{1}, t_{1}^{\prime}\right]$



## Link Prediction

- Different scoring functions $c(x, y)=$
- Graph distance: (negated) Shortest path length
- Common neighbors: $|\Gamma(x) \cap \Gamma(y)|$
- Jaccard's coefficient: $|\Gamma(x) \cap \Gamma(y)| /|\Gamma(x) \cup \Gamma(y)|$
- Adamic/Adar: $\sum_{z \in \Gamma(x) \cap \Gamma(y)} 1 / \log |\Gamma(z)|$
- Preferential attachment: $|\Gamma(x)| \cdot|\Gamma(y)| \quad \Gamma(x) \ldots$ neighbors
- PageRank: $r_{x}(y)+r_{y}(x)$ of node $x$
- $r_{x}(y)$... stationary distribution score of $y$ under the random walk:
- with prob. 0.15 , jump to $x$
- with prob. 0.85 , go to random neighbor of current node
- Then, for a particular choice of $c(\cdot)$
- For every pair of nodes $(x, y)$ compute $c(x, y)$
- Sort pairs $(x, y)$ by the decreasing score $c(x, y)$
- Predict top $n$ pairs as new links


## Link Prediction Methods

Thanks to Jure Leskovec, Stanford and Panayiotis Tsaparas, Univ. of loannina for slides

## Link Prediction Methods

- How to assign the score $c(x, y)$ for each pair $(x, y)$ ?
- Some form of similarity between $\mathbf{x}$ and $\mathbf{y}$
- Some form of node proximity between $\mathbf{x}$ and $\mathbf{y}$
- Methods
- Neighborhood-based (shared neighbors)
- Network proximity based (paths between $\mathbf{x}$ and $\mathbf{y}$ )
- Other


## Methods for Link Prediction

Neighborhood-based

## Neighborhood-based Methods

Let $\Gamma(\mathbf{x})$ be the set of neighbors of $\mathbf{x}$ in $\mathbf{G}_{\text {old }}$

- Methods
- Common Neighbors Overlap
- Jaccard
- Adamic/Adar
- Preferential Attachment


## Neighborhood-based Methods

Intuition: The larger the overlap of the neighbors of two nodes, the more likely the nodes to be linked in the future

- Common neighbors
- A: adjacency matrix, $\mathrm{A}_{\mathrm{x}, \mathrm{y}}{ }^{2}$ : \#paths of length 2

$$
\operatorname{score}(x, y)=\| \Gamma(x) \cap \Gamma(y) \mid
$$

- Jaccard coefficient
- The probability that both $\mathbf{x}$ and $\mathbf{y}$ have a feature for a randomly selected feature that either $\mathbf{x}$ or $\mathbf{y}$ has

$$
\operatorname{score}(x, y)=\frac{\|\Gamma(x) \cap \Gamma(y)\|}{\| \Gamma(x) \cup \Gamma(y) \mid}
$$

## Neighborhood-based Methods

- Adamic/Adar
- Assigns large weights to common neighbors $\mathbf{z}$ of $\mathbf{x}$ and $\mathbf{y}$ which themselves have few neighbors (weight rare features more heavily)

$$
\operatorname{score}(x, y)=\sum_{z \in \Gamma(x) \cap \Gamma(y)} \frac{1}{\log |\Gamma(z)|}
$$

- Preferential attachment
- Based on the premise that the probability that a new edge has node $\mathbf{x}$ as its endpoint is proportional to $|\Gamma(\mathbf{x})|$, i.e., nodes like to form ties with 'popular' nodes

$$
\operatorname{score}(x, y)=|\Gamma(x)||\Gamma(y)|
$$

## Methods for Link Prediction

Network proximity based

## Network Proximity Methods

Intuition: The "closer" two nodes are in the network, the more likely are to be linked in the future

- Methods
- based on shortest path length between $\mathbf{x}$ and $\mathbf{y}$
- based on all paths between $\mathbf{x}$ and $\mathbf{y}$
- Katz ${ }_{\beta}$ measure (unweighted, weighted)
- Random walk-based
- hitting time
- commute time
- Rooted PageRank
- SimRank


## Shortest Path Based

For $\mathbf{x}, \mathbf{y} \in \mathbf{V} \times \mathbf{V}-\mathbf{E}_{\text {old }}$,

## score $(x, y)=($ negated $)$ length of shortest path between $x$ and $y$

If there are more than $\mathbf{n}$ pairs of nodes tied for the shortest path length, order them at random

## Ensemble of All Paths

- Katz $_{\beta}$ measure

$$
\operatorname{score}(x, y):=\sum_{\ell=1}^{\infty} \beta^{\ell} \cdot\left|\operatorname{paths}_{x, y}^{\langle\ell\rangle}\right|
$$

- Sum over all paths of length $\ell$
- $0<6<1$ : a parameter of the predictor, exponentially damped to count short paths more heavily


## Ensemble of All Paths

- Katz $_{\beta}$ measure

$$
\begin{aligned}
\operatorname{score}(x, y) & :=\sum_{\ell=1}^{\infty} \beta^{\ell} \cdot\left|\operatorname{paths}_{x, y}^{\langle\ell,}\right| \\
\sum_{l=1}^{\infty} \beta^{l} \cdot \mid \text { paths }_{x y}^{(l)} \mid & =\beta A_{x y}+\beta^{2}\left(A^{2}\right)_{x y}+\beta^{3}\left(A^{3}\right)_{x y}+\cdots
\end{aligned}
$$

- Unweighted version: path $_{x, y}(\mathbf{1})=1$, if $x$ and $y$ have collaborated, $\mathbf{0}$ otherwise
- Weighted version: path ${ }_{x, y}(\mathbf{1})=$ \#times $x$ and $y$ have collaborated


## Random Walk Based

Consider a random walk on $\mathbf{G}_{\text {old }}$ that starts at $\mathbf{x}$ and iteratively moves to a neighbor of $\mathbf{x}$ chosen uniformly at random from $\Gamma(x)$

- Hitting $\mathbf{H}_{\mathbf{x}, \mathbf{y}}$ (from $\mathbf{x}$ to $\mathbf{y}$ ): the expected number of steps it takes for the random walk starting at $\mathbf{x}$ to reach $\mathbf{y}$

$$
\operatorname{score}(x, y)=-H_{x, y}
$$

- Commute Time $C_{x, y}$ (from $x$ to $y$ ): the expected number of steps to travel from $\mathbf{x}$ to $\mathbf{y}$ and from $\mathbf{y}$ to $\mathbf{x}$

$$
\operatorname{score}(x, y)=-\left(H_{x, y}+H_{y, x}\right)
$$

Not symmetric, can be shown

$$
\begin{aligned}
& h_{v u}=\Theta\left(n^{2}\right) \\
& h_{u v}=\Theta\left(n^{3}\right)
\end{aligned}
$$



## Random Walk Based

- The hitting time and commute time measures are sensitive to parts of the graph far away from $x$ and $y \rightarrow$ periodically reset the walk
- Random walk on $\mathbf{G}_{\text {old }}$ that starts at $\boldsymbol{x}$ and has a probability $\boldsymbol{\alpha}$ of returning to $\mathbf{x}$ at each step
- Rooted PageRank
- Starts from x
- with probability $(1-a)$ moves to a random neighbor
- with probability $\boldsymbol{a}$ returns to $\mathbf{x}$
$\operatorname{score}(x, y)=$ stationary probability of $y$ in a
rooted PageRank


## SimRank

Intuition: Two objects are similar, if they are related to similar objects

- Two objects $\boldsymbol{x}$ and $\boldsymbol{y}$ are similar, if they are related to objects $a$ and $b$ respectively and $a$ and $b$ are themselves similar

$$
\operatorname{similarity}(x, y):=\gamma \cdot \frac{\sum_{a \in \Gamma(x)} \sum_{b \in \Gamma(y)} \text { similarity }(a, b)}{|\Gamma(x)| \cdot|\Gamma(y)|}
$$

- Expresses the average similarity between neighbors of $\mathbf{x}$ and neighbors of $\mathbf{y}$ $\operatorname{score}(x, y)=\operatorname{similarity}(x, y)$


## Methods for Link Prediction

Other Methods

## Other Methods

Low-rank Approximations

- Unseen bigrams
- High-level Clustering


## Low Rank Approximations

Intuition: represent the adjacency matrix $\mathbf{M}$ with a
lower rank matrix $\mathbf{M}_{\mathbf{k}}$

- Method
- Apply SVD (singular value decomposition)
- Obtain the rank-k matrix that best approximates M


## Singular Value Decomposition

$$
\underset{\substack{[n \times r][r \times r][r \times n]}}{\mathrm{A}}=\mathrm{U} \quad \Sigma \quad \mathrm{~V}^{\top}=\left[\begin{array}{llll}
\overrightarrow{\mathrm{u}}_{1} & \overrightarrow{\mathrm{u}}_{2} & \cdots & \overrightarrow{\mathrm{u}}_{\mathrm{r}}
\end{array}\right]\left[\begin{array}{llll}
\sigma_{1} & & & \\
& \sigma_{2} & & \\
& & \ddots & \\
& & & \sigma_{\mathrm{r}}
\end{array}\right]\left[\begin{array}{c}
\overrightarrow{\mathrm{v}}_{1} \\
\overrightarrow{\mathrm{v}}_{2} \\
\vdots \\
\overrightarrow{\mathrm{v}}_{\mathrm{r}}
\end{array}\right]
$$

- r: rank of matrix A
- $\sigma_{1} \geq \sigma_{2} \geq \ldots \geq \sigma_{r}$ : singular values (square roots of eig-vals $A A^{\top}, A^{\top} A$ )
- $\overrightarrow{\mathrm{u}}_{1}, \overrightarrow{\mathrm{u}}_{2}, \cdots, \overrightarrow{\mathrm{u}}_{r} \quad$ : left singular vectors (eig-vectors of $A A^{\top}$ )
- $\vec{V}_{1}, \overrightarrow{\mathrm{~V}}_{2}, \ldots, \overrightarrow{\mathrm{~V}}_{\mathrm{r}}$ : right singular vectors (eig-vectors of $\mathrm{A}^{\top} A$ )
- $A=\sigma_{1} \overrightarrow{\mathrm{u}}_{1} \overrightarrow{\mathrm{v}}_{1}^{\mathrm{T}}+\sigma_{2} \overrightarrow{\mathrm{u}}_{2} \overrightarrow{\mathrm{v}}_{2}^{\mathrm{T}}+\cdots+\sigma_{\mathrm{r}} \overrightarrow{\mathrm{u}}_{\mathrm{r}} \overrightarrow{\mathrm{v}}_{\mathrm{r}}^{\mathrm{T}}$
- $\mathrm{A}_{\mathrm{k}}=\sigma_{1} \overrightarrow{\mathrm{u}}_{1} \overrightarrow{\mathrm{v}}_{1}^{\mathrm{T}}+\sigma_{2} \overrightarrow{\mathrm{u}}_{2} \overrightarrow{\mathrm{v}}_{2}^{\mathrm{T}}+\cdots+\sigma_{\mathrm{k}} \overrightarrow{\mathrm{u}}_{\mathrm{k}} \overrightarrow{\mathrm{v}}_{\mathrm{k}}^{\mathrm{T}}, k \in\{1,2, \ldots, r\}$


## Unseen bigrams

- Unseen bigrams: Predict pairs of words that co-occur in a test corpus, but not in the corresponding training corpus
- Not just score( $\mathbf{x}, \mathbf{y}$ ) but score(z, $\mathbf{y})$ for nodes $\mathbf{z}$ that are similar to $\mathbf{x}--\mathbf{S}_{\mathbf{x}}{ }^{(\delta)}$ : the $\delta$ nodes most related to $x$

$$
\begin{aligned}
\text { score }_{\text {unweighted }}^{*}(x, y) & :=\left|\left\{z: z \in \Gamma(y) \cap S_{x}^{\langle\delta\rangle}\right\}\right| \\
\text { score }_{\text {weighted }}^{*}(x, y) & :=\sum_{z \in \Gamma(y) \cap S_{x}^{(\delta\rangle}} \operatorname{score}(x, z)
\end{aligned}
$$



## Clustering

- Compute score( $\mathbf{x}, \mathbf{y}$ ) for all edges in $\mathbf{E}_{\text {old }}$
- Delete the (1-p) fraction of the edges whose score is the lowest, for some parameter $\mathbf{p}$
- Re-compute score( $\mathbf{x}, \mathrm{y}$ ) for all pairs in the subgraph


## Evaluation \& Results

## How to Evaluate the Prediction

- Each link predictor $\boldsymbol{p}$ outputs a ranked list $L_{p}$ of pairs in $\mathrm{V} \times \mathrm{V}-\mathrm{E}_{\text {old }}$ in decreasing order of confidence
- focus on Core network, (d>3)

$$
E *_{\text {new }}=E_{\text {new }} \cap(\text { Core } \times \text { Core })=\left|E *_{\text {new }}\right|
$$

- Evaluation method: Size of intersection of
- the first $\mathbf{n}$ edge predictions from $\mathbf{L}_{\mathrm{p}}$ that are in Core $\times$ Core, and
- the actual set $E *_{\text {new }}$


## Evaluation: Baseline Predictor

- Random Predictor: Randomly select pairs of authors who did not collaborate in the training interval
- Probability that a random prediction is correct:

$$
\frac{\left|E_{\text {new }}\right|}{\binom{\text { Corel }}{2}-\left|E_{\text {old }}\right|}
$$

In the datasets, from $0.15 \%$ (cond-mat) to $0.48 \%$ (astro-ph)

## Average Relevance Performance

- Improvement over random predictor

- average ratio over the five datasets of the given predictor's performance versus a baseline predictor's performance.
- the error bars indicate the minimum and maximum of this ratio over the five datasets.
- the parameters for the starred predictors are: (1) for weighted Katz, $\beta=0.005$; (2) for Katz clustering, $\beta 1=0.001 ; \rho=0.15 ; \beta 2$ = 0.1; (3) for low-rank inner product, rank = 256; (4) for rooted Pagerank, $\alpha=0.15$; (5) for unseen bigrams, unweighted, common neighbors with $\delta=8$; and (6) for SimRank, C ( $\gamma$ ) = 0.8.


## Results: Common Neighbors

- Improvement over \#common neighbors



## Results: Common Neighbors

- Improvement over graph distance predictor



## Results: Improvement




## Factor Improvement Over Random

| predictor | astro-ph | cond-mat | gr-qc | hep-ph | hep-th |
| :---: | :---: | :---: | :---: | :---: | :---: |
| probability that a random prediction is correct | 0.475\% | 0.147\% | 0.341\% | 0.207\% | 0.153\% |
| graph distance (all distance-two pairs) | 9.4 | 25.1 | 21.3 | 12.0 | 29.0 |
| common neighbors | 18.0 | 40.8 | 27.1 | 26.9 | 46.9 |
| preferential attachment | 4.7 | 6.0 | 7.5 | 15.2 | 7.4 |
| Adamic/Adar | 16.8 | 54.4 | 30.1 | 33.2 | 50.2 |
| Jaccard | 16.4 | 42.0 | 19.8 | 27.6 | 41.5 |
| SimRank $\gamma=0.8$ | 14.5 | 39.0 | 22.7 | 26.0 | 41.5 |
| hitting time | 6.4 | 23.7 | 24.9 | 3.8 | 13.3 |
| hitting time - normed by stationary distribution | 5.3 | 23.7 | 11.0 | 11.3 | 21.2 |
| commute time | 5.2 | 15.4 | 33.0 | 17.0 | 23.2 |
| commute time - normed by stationary distribution | 5.3 | 16.0 | 11.0 | 11.3 | 16.2 |
| rooted PageRank $\quad \alpha=0.01$ | 10.8 | 27.8 | 33.0 | 18.7 | 29.1 |
| $\alpha=0.05$ | 13.8 | 39.6 | 35.2 | 24.5 | 41.1 |
| $\alpha=0.15$ | 16.6 | 40.8 | 27.1 | 27.5 | 42.3 |
| $\alpha=0.30$ | 17.1 | 42.0 | 24.9 | 29.8 | 46.5 |
| $\alpha=0.50$ | 16.8 | 40.8 | 24.2 | 30.6 | 46.5 |
| Katz (weighted) $\beta=0.05$ | 3.0 | 21.3 | 19.8 | 2.4 | 12.9 |
| $\beta=0.005$ | 13.4 | 54.4 | 30.1 | 24.0 | 51.9 |
| $\beta=0.0005$ | 14.5 | 53.8 | 30.1 | 32.5 | 51.5 |
| Katz (unweighted) $\beta=0.05$ | 10.9 | 41.4 | 37.4 | 18.7 | 47.7 |
| $\beta=0.005$ | 16.8 | 41.4 | 37.4 | 24.1 | 49.4 |
| $\beta=0.0005$ | 16.7 | 41.4 | 37.4 | 24.8 | 49.4 |

## Factor Improvement Over Random

| predictor | astro-ph | cond-mat | gr-qc | hep-ph | hep-th |
| :---: | :---: | :---: | :---: | :---: | :---: |
| probability that a random prediction is correct | 0.475\% | 0.147\% | 0.341\% | 0.207\% | 0.153\% |
| graph distance (all distance-two pairs) | 9.4 | 25.1 | 21.3 | 12.0 | 29.0 |
| common neighbors | 18.0 | 40.8 | 27.1 | 26.9 | 46.9 |
| Low-rank approximation: rank $=1024$ | 15.2 | 53.8 | 29.3 | 34.8 | 49.8 |
| Inner product rank $=256$ | 14.6 | 46.7 | 29.3 | 32.3 | 46.9 |
| rank $=64$ | 13.0 | 44.4 | 27.1 | 30.7 | 47.3 |
| rank $=16$ | 10.0 | 21.3 | 31.5 | 27.8 | 35.3 |
| rank $=4$ | 8.8 | 15.4 | 42.5 | 19.5 | 22.8 |
| rank $=1$ | 6.9 | 5.9 | 44.7 | 17.6 | 14.5 |
| Low-rank approximation: $\quad$ rank $=1024$ | 8.2 | 16.6 | 6.6 | 18.5 | 21.6 |
| Matrix entry rank $=256$ | 15.4 | 36.1 | 8.1 | 26.2 | 37.4 |
| rank $=64$ | 13.7 | 46.1 | 16.9 | 28.1 | 40.7 |
| rank $=16$ | 9.1 | 21.3 | 26.4 | 23.1 | 34.0 |
| rank $=4$ | 8.8 | 15.4 | 39.6 | 20.0 | 22.4 |
| rank $=1$ | 6.9 | 5.9 | 44.7 | 17.6 | 14.5 |
| Low-rank approximation: rank $=1024$ | 11.4 | 27.2 | 30.1 | 27.0 | 32.0 |
| Katz ( $\beta=0.005$ ) rank $=256$ | 15.4 | 42.0 | 11.0 | 34.2 | 38.6 |
| rank $=64$ | 13.1 | 45.0 | 19.1 | 32.2 | 41.1 |
| rank $=16$ | 9.2 | 21.3 | 27.1 | 24.8 | 34.9 |
| rank $=4$ | 7.0 | 15.4 | 41.1 | 19.7 | 22.8 |
| rank $=1$ | 0.4 | 5.9 | 44.7 | 17.6 | 14.5 |
| unseen bigrams common neighbors, $\delta=8$ | 13.5 | 36.7 | 30.1 | 15.6 | 46.9 |
| (weighted) common neighbors, $\delta=16$ | 13.4 | 39.6 | 38.9 | 18.5 | 48.6 |
| Katz $(\beta=0.005), \delta=8$ | 16.8 | 37.9 | 24.9 | 24.1 | 51.1 |
| Katz $(\beta=0.005), \delta=16$ | 16.5 | 39.6 | 35.2 | 24.7 | 50.6 |
| unseen bigrams common neighbors, $\delta=8$ | 14.1 | 40.2 | 27.9 | 22.2 | 39.4 |
| (unweighted) common neighbors, $\delta=16$ | 15.3 | 39.0 | 42.5 | 22.0 | 42.3 |
| Katz $(\beta=0.005), \delta=8$ | 13.1 | 36.7 | 32.3 | 21.6 | 37.8 |
| Katz $(\beta=0.005), \delta=16$ | 10.3 | 29.6 | 41.8 | 12.2 | 37.8 |
| clustering: $\quad \rho=0.10$ | 7.4 | 37.3 | 46.9 | 32.9 | 37.8 |
| Katz ( $\left.\beta_{1}=0.001, \beta_{2}=0.1\right) \quad \rho=0.15$ | 12.0 | 46.1 | 46.9 | 21.0 | 44.0 |
| $\rho=0.20$ | 4.6 | 34.3 | 19.8 | 21.2 | 35.7 |
| $\rho=0.25$ | 3.3 | 27.2 | 20.5 | 19.4 | 17.4 |

## Evaluation：Prediction Overlap



How similar are the predictions made by the different methods？ \＃common predictions
\＃correct predictions

|  | $\begin{aligned} & \text { 券 } \\ & \text { 品 } \\ & \text { ? } \end{aligned}$ |  |  |  |  | $\begin{aligned} & \text { N } \\ & \text { 要 } \\ & \text { 曹 } \\ & \text { है } \end{aligned}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Adamic／Adar | 92 | 65 | 53 | 22 | 43 | 87 | 72 | 44 | 36 | 49 |
| Katz clustering |  | 78 | 41 | 20 | 29 | 66 | 60 | 31 | 22 | 37 |
| common neighbors |  |  | 69 | 13 | 43 | 52 | 43 | 27 | 26 | 40 |
| hitting time |  |  |  | 40 | 8 | 22 | 19 | 17 | 9 | 15 |
| Jaccard＇s coefficient |  |  |  |  | 71 | 41 | 32 | 39 | 51 | 43 |
| weighted Katz |  |  |  |  |  | 92 | 75 | 44 | 32 | 51 |
| low－rank inner product |  |  |  |  |  |  | 79 | 39 | 26 | 46 |
| rooted Pagerank |  |  |  |  |  |  |  | 69 | 48 | 39 |
| SimRank |  |  |  |  |  |  |  |  | 66 | 34 |
| unseen bigrams |  |  |  |  |  |  |  |  |  | 68 |

## Extensions

- Improve performance. Even the best (Katz clustering on gr-qc) correct on only about 16\% of its prediction
- Improve efficiency on very large networks (approximation of distances)
- Treat more recent links (e.g., collaborations) as more important
- Additional information (paper titles, author institutions, etc) latently present in the graph

Facebook: Supervised Random Walks for Link Prediction

## Supervised Link Prediction

- Can we learn to predict new friends?
- Facebook's People You May Know

Friendship by Nunber of Hops

- Let's look at the FB data:
- 92\% of new friendships on FB are friend-of-a-friend
- More mutual friends helps




## Supervised Link Prediction

- Goal: Recommend a list of possible friends
- Supervised machine learning setting:
- Labeled training examples:
- For every user $s$ have a list of others she will create links to $\left\{d_{1} \ldots d_{k}\right\}$ in the future
- Use FB network from May 2012 and $\left\{d_{1} \ldots d_{k}\right\}$ are the new friendships you created since then
" These are the "positive" training examples
- Use all other users as "negative" example
- Task:
- For a given node $s$, score nodes $\left\{d_{1} \ldots d_{k}\right\}$ higher than any other node in the network

- "positive" nodes
"negative" nodes
Green nodes are the nodes to which s creates links in the future


## Supervised Link Prediction

- How to combine node/edge features and the network structure?
- Estimate strength of each friendship ( $u, v$ ) using:
- Profile of user $u$, profile of user $v$
- Interaction history of users $u$ and $v$
- This creates a weighted graph
- Do Personalized PageRank from $s$ and measure the "proximity" (the visiting prob.) of any other node $w$ from $s$
- Sort nodes $w$ by decreasing

" "positive" nodes
"negative" nodes "proximity"


## Supervised Random Walks

- Let $s$ be the starting node
- Let $f_{\beta}(u, v)$ be a function that assigns strength $a_{u v}$ to edge ( $u, v$ ) $a_{u v}=f_{\beta}(u, v)=\exp \left(-\sum_{i} \beta_{i} \cdot \mathrm{x}_{u v}[i]\right)$
- $x_{u v}$ is a feature vector of $(u, v)$
- Features of node $u$
- "positive" nodes
"negative" nodes
- Features of node $v$
- Features of edge ( $u, v$ )
- Note: $\beta$ is the weight vector we will later estimate!
- Do Random Walk with Restarts from $s$ where transitions are according to edge strengths $a_{u v}$


## SRW: Prediction



Network


Set edge strengths
$a_{u v}=f_{\beta}(u, v)$

- How to estimate edge strengths?
- How to set parameters $\beta$ of $f_{\beta}(u, v)$ ?
- Idea: Set $\beta$ such that it (correctly) predicts the known future links

Random Walk with
Restarts on the
weighted graph.
Each node $w$ has a PageRank proximity $p_{w}$


Sort nodes $w$ by the decreasing PageRank score $p_{w}$


Recommend top $k$ nodes with the highest proximity $p_{w}$ to node $s$

## Personalized PageRank

- $a_{u v}$.... Strength of edge ( $u, v$ )
- Random walk transition matrix:

$$
Q_{u v}^{\prime}= \begin{cases}\frac{a_{u v}}{\sum_{w} a_{u w}} & \text { if }(u, v) \in E, \\ 0 & \text { otherwise }\end{cases}
$$



- PageRank transition matrix:
- "positive" nodes
"negative" nodes

$$
Q_{i j}=(1-\alpha) Q_{i j}^{\prime}+\alpha \mathbf{1}(j=s)
$$

- Where with prob. $\alpha$ we jump back to node $s$
- Compute PageRank vector: $p=p^{T} Q$
- Rank nodes $w$ by decreasing $p_{w}$


## The Optimization Problem

- Positive examples
$D=\left\{\boldsymbol{d}_{1}, \ldots, \boldsymbol{d}_{\boldsymbol{k}}\right\}$
- Negative examples $L=\{$ other nodes $\}$
- What do we want?

$$
\min _{\beta} F(\beta)=\|\beta\|^{2}
$$

We prefer small
weights $\beta$ to prevent overfitting


- "positive" nodes
"negative" nodes
such that

$$
\forall d \in D, l \in L: p_{l}<p_{d}
$$

- Note:

Every positive example has to have higher PageRank score than every negative example

- Exact solution to this problem may not exist
- So we make the constraints "soft" (i.e., optional)


## Making Constraints "Soft"

- Want to minimize:
$\min _{\beta} F(\beta)=\sum_{d \in D, l \in L} h\left(p_{l}-p_{d}\right)+\lambda\|\beta\|^{2}$

$$
d \in \overline{D, l \in L}
$$

- Loss: $h(x)=0$ if $x<0$, or $x^{2}$ else


$$
p_{l}<p_{d} \quad p_{l}=p_{d} \quad p_{l}>p_{d}
$$

Penalty for violating the constraint that $p_{d}>p_{l}$

## Solving the Problem: Intuition

- How to minimize F?

$$
\min _{\beta} F(\beta)=\sum_{d \in D, l \in L} h\left(p_{l}-p_{d}\right)+\lambda\|\beta\|^{2}
$$

- Both $p_{l}$ and $p_{d}$ depend on $\beta$
- Given $\beta$ assign edge weights $a_{u v}=f_{\beta}(u, v)$
- Using $Q=\left[a_{u v}\right]$ compute PageRank scores $p_{\beta}$
- Rank nodes by the decreasing score
- Goal: Want to find $\beta$ such that $p_{l}<p_{d}$


## Solving the Problem: Intuition

- How to minimize $\boldsymbol{F}(\boldsymbol{\beta})$ ?

$$
\min _{\beta} F(\beta)=\sum_{d \in D, l \in L} h\left(p_{l}-p_{d}\right)+\lambda\|\beta\|^{2}
$$

- Idea:
- Start with some random $\beta^{(0)}$
- Evaluate the derivative of $F(\beta)$ and do a small step in the opposite direction

$$
\beta^{(t+1)}=\beta^{(t)}-\eta \frac{\partial F\left(\beta^{(t)}\right)}{\partial \beta}
$$

- Repeat until convergence



## Gradient Descent

- What's the derivative $\frac{\partial F\left(\beta^{(t)}\right)}{\partial \beta}$ ?

$$
F(\beta)=\sum_{d \in D, l \in L} h\left(p_{l}-p_{d}\right)+\lambda\|\beta\|^{2}
$$

$$
\begin{aligned}
\frac{\partial F(\beta)}{\partial \beta} & = \\
& \quad \sum_{l, d} \frac{\partial h\left(p_{l}-p_{d}\right)}{\partial \beta}+2 \lambda \beta \\
& =\sum_{l, d} \frac{\partial h\left(p_{l}-p_{d}\right)}{\partial\left(p_{l}-p_{d}\right)}\left(\frac{\partial p_{l}}{\partial \beta}-\frac{\partial p_{d}}{\partial \beta}\right)+2 \lambda \beta
\end{aligned}
$$

$$
h(x)=\max \{x, 0\}^{2}
$$ Easy!

- We know:

$$
p=p^{T} Q \text { that is } p_{u}=\sum_{j} p_{j} Q_{j u}
$$

- So:

$$
\frac{\partial p_{u}}{\partial \beta}=\sum_{j} Q_{j u} \frac{\partial p_{j}}{\partial \beta}+p_{j} \frac{\partial Q_{j u}}{\partial \beta}
$$

## Gradient Descent

- We just got: $\frac{\partial p_{u}}{\partial \beta}=\sum_{j} Q_{j u} \frac{\partial p_{j}}{\partial \beta}+p_{j} \frac{\partial Q_{j u}}{\partial \beta}$
- Few details:
- Computing $\partial Q_{j u} / \partial \beta$ is easy. Remember: $Q_{u v}^{\prime}= \begin{cases}\frac{a_{w o}}{\sum_{w w_{u w}}} & \text { if }(u, v) \in E, \\ 0 & \text { otherwise }\end{cases}$
- We want $\frac{\partial p_{j}}{\partial \beta}$ but it appears on both sides of the equation. Notice the whole thing looks like a PageRank

$$
\begin{aligned}
& a_{u v}=f_{\beta}(u, v) \\
& =\exp \left(-\sum_{i} \beta_{i} \cdot \mathrm{x}_{u v}[i]\right)
\end{aligned}
$$ equation: $x=Q \cdot x+z$

- As with PageRank we can use the power-iteration to solve it:
- Start with a random $\frac{\partial p^{(0)}}{\partial \beta}$
- Then iterate: $\frac{\partial p}{\partial \beta}^{(t+1)}=Q \cdot \frac{\partial p}{\partial \beta}^{(t)}+\frac{\partial Q_{j u}}{\partial \beta} \cdot p$


## Optimizing $\boldsymbol{F}(\boldsymbol{\beta})$

- To optimize $\boldsymbol{F}(\boldsymbol{\beta})$, use gradient descent:
- Pick a random starting point $\beta^{(0)}$
- Using current $\beta^{(t)}$ compute edge strenghts and the transition matrix $Q$
- Compute PageRank scores $p$
- Compute the gradient with respect to weight vector $\beta^{(t)}$
- Update $\beta^{(t+1)}$



## Data: Facebook

- Facebook Iceland network
- 174,000 nodes (55\% of population)
- Avg. degree 168
- Avg. person added 26 friends/month
- For every node $s$ :

- Positive examples:
- $D=\{$ new friendships $s$ created in Nov '09 \}
- Negative examples:
- $L=\{$ other nodes $s$ did not create new links to $\}$
- Limit to friends of friends:
- On avg. there are 20,000 FoFs (maximum is 2 million)!


## Experimental Setting

- Node and Edge features for learning:
- Node: Age, Gender, Degree
- Edge: Age of an edge, Communication, Profile visits, Co-tagged photos
- Evaluation:
- Precision at top 20
- We produce a list of 20 candidates
- By taking top 20 nodes $x$ with highest PageRank score $p_{x}$
- Measure to what fraction of these nodes $s$ actually links to


## Results: Facebook Iceland

- Facebook: Predict future friends
- Adamic-Adar already works great
- Supervised Random Walks (SRW) gives slight improvement

| Learning Method | Prec@Top20 |
| :--- | ---: |
| Random Walk with Restart | 6.80 |
| Adamic-Adar | 7.35 |
| Common Friends | 7.35 |
| Degree | 3.25 |
| SRW: one edge type | 6.87 |
| SRW: multiple edge types | $\mathbf{7 . 5 7}$ |

## Results: Facebook

- 2.3x improvement over previous FB-PYMK (People You May Know)

Fraction of Friending from PYMK


## Results: Co-Authorship

- Arxiv Hep-Ph collaboration network:
- Poor performance of unsupervised methods
- SRW gives a boost of 25\%!

| Learning Method | Prec@Top20 |
| :--- | ---: |
| Random Walk with Restart | 3.41 |
| Adamic-Adar | 3.13 |
| Common Friends | 3.11 |
| Degree | 3.05 |
| SRW: one edge type | 4.24 |
| SRW: multiple edge types | $\mathbf{4 . 2 5}$ |

## Wtf: The Who to Follow Service at Twitter

## Introduction

## Ewitter

```
Home Profile Find People Settings Campaigns Help Sign out
```


## Users you may be interested in



Semantic differences between "interested in" and "similar to"

## WtF ("Who to Follow")

- WtF ("Who to Follow"): the Twitter user recommendation service
- help existing and new users to discover connections to sustain and grow
- used for search relevance, content discovery, promoted products, etc.
- Twitter Data:
- 200 million users
- 400 million tweets every day (as of early 2013)
- http://www.internetlivestats.com/twitter-statistics/


## The Twitter Graph

- Graph
- Node: user
- (Directed) Edge: follows
- Graph Statistics (Aug'12)
- Over 20 billion edges
- Power law of in- and out-degrees
- Over 1000 with more than 1 million followers
- 25 users with more than 10 million followers


## Algorithms: Circle of Trust

- Circle of Trust
- Based on an egocentric random walk (similar to personalized (rooted) PageRank)
- Computed in an online fashion (from scratch each time) given a set of parameters
- \# of random walk steps
- reset probability
- pruning settings to discard low probability vertices
- parameters to control sampling of outgoing edges at vertices with large out-degrees


## Algorithms

- Directed edge
- Asymmetric nature of the follow relationship
- Friendships in other social networks such as Facebook or LinkedIn are symmetric/reciprocal
- Similar to the user-item recommendations problem where the "item" is also a user


## Algorithms: SALSA

- SALSA (Stochastic Approach for Link-Structure Analysis)
- a variation of HITS
- HITS
- Intuition:
- Good hubs point to good authorities
- Good auth. are pointed by good hubs
- Recurs. comput. of hub score
- Recurs. comput. of auth. score

hubs authorities

$$
h_{i}=\sum_{j: i \rightarrow j} a_{j}
$$

$$
a_{i}=\sum_{j: j \rightarrow i} h_{j}
$$

## Algorithms: SALSA

- Random walks to rank hubs and authorities
- Two different random walks (Markov chains): a chain of hubs and a chain of authorities
- Each walk traverses nodes only in one side by traversing two links in each step $h \rightarrow a \rightarrow h, a \rightarrow h \rightarrow a$

Transition matrices of each chain: H and A W: the adjacency of the directed graph $\mathbf{W}_{\mathbf{r}}$ : divide each entry by the sum of its row $\mathbf{W}_{\mathbf{c}}$ : divide each entry by the sum of its column
$H=W_{r} W_{c}{ }^{\top}$
$A=W_{c}{ }^{\top} W_{r}$

hubs
authorities

## Algorithms: SALSA

- Reduces to the problem of HITS with tightly knit communities
- TKC effect
- Better for single-topic communities
- More efficient implementation


## HITS and the TKC effect

The HITS algorithm favors the most dense community of hubs and authorities

- Tightly Knit Community (TKC) effect



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after $\mathbf{n}$ iterations


## HITS and the TKC effect

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after normalization with the max
element as $\mathrm{n} \rightarrow \infty$


## Algorithms: SALSA

- Hubs:
- 500 top-ranked nodes from a user's circle of trust
- user similarity (based on homophily, also useful)
- Authorities:
- users that the hubs follow
- "interested in" user recommendations



## Algorithms: SALSA

- SALSA's recursive nature
- Two users are similar if they follow the same (or similar) users (LHS)
- A user u is likely to follow those who are followed by users that are similar to $\mathbf{u}$ (RHS)
- The random walk ensures fair distribution of scores in both directions
- Similar users are selected from the circle of trust of a user (via personalized PageRank)


## Evaluation

- Approaches
- Offline experiments on retrospective data
- Online $A / B$ testing on live traffic
- Various parameters may interfere:
- How the results are rendered
- Platform (mobile, etc.)
- New vs old users


## Extensions

- Add metadata to vertices (e.g., user profile information) and edges (e.g., edge weights, timestamp, etc.)
- Consider interaction graphs (e.g., graphs defined in terms of retweets, favorites, replies, etc.)


## Extensions

- Two phase algorithm
- $1^{\text {st }}$ - Candidate generation: produce a list of promising recommendations for each user, using any algorithm
- $\mathbf{2}^{\text {nd }}$ - Rescoring: apply a machine-learned model to the candidates, binary classification problem (logistic regression)
- Evaluation
${ }^{\text {- }}{ }^{\text {st }}$ Phase: recall + diversity
- $\mathbf{2}^{\text {nd }}$ Phase: precision + maintain diversity

