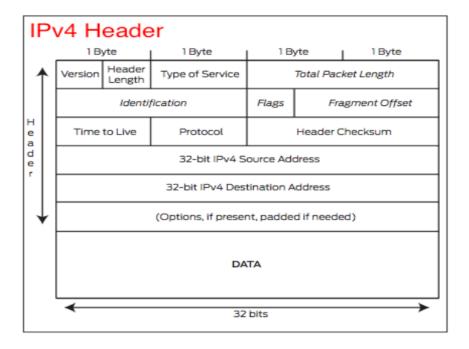


Syed Saad Siar Ahmad Khan

IPv4 VS IPv6



1 Byt		1 Byte	1 Byte	1 Byte
Version	Traffic Class		Flov	Label
	Payload Length		Next Header	Hop Limit
		128-bit IPv6	Source Address	
		128-bit IPv6	Source Address	
		128-bit IPv6	Source Address	
	 		Source Address	



Why switch to IPv6?

Population of Earth W / Population / World Population Current World Population 7,676,691,666

view all people on 1 page >

TODAY	THIS YEAR	
Births today 305,969	Births this year 4,547,984	
Deaths today 128,375	Deaths this year 1,908,189	
Population Growth today 177,594	Population Growth this year 2,639,795	

Number of unique IP's in IPv4

			2 ³² =		
	4294967296				
()	%	AC		
7	8	9	÷		
4	5	6	×		
1	2	3	-		
0	•	=	+		

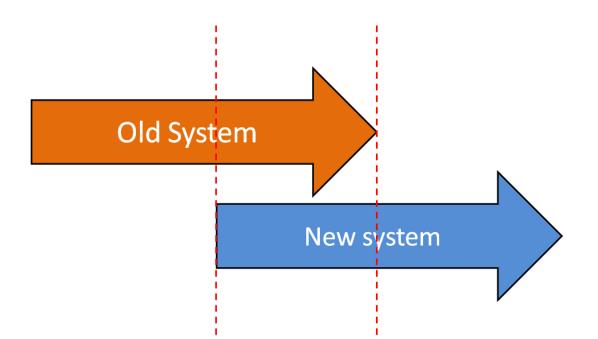






RELATIVELY NEW PROTOCOL

- IPv6 is not developed enough yet and may encounter unforeseen vulnerabilities in the future.
- Security products such as firewalls are not as available for the IPv6 protocol compared to IPv4.





COEXISTING LEADS TO GREATER RISK

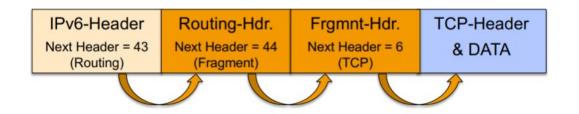
- Transition and Dual-stacked Technologies have been deployed which increases complexity in the networks.
- Lack of support from technical personnel





EXTENSION HEADERS

- Extension Headers having varying size.
- Headers are chained.
- Too many EH(s) can harm the network.
- Malicious EH(s) still have to be processed.





FRAGMENTATION ISSUES

Problems

Countermeasures

- Can avoid filtering and detection techniques (IDS/IPS evasion techniques).
- Can harm destination devices upon reassembly.

- Monitor Number of Fragments and discard if over a set amount.
- Block unequal fragments except for last one of a set.



- **1.Attacks against the IPv6 protocol**
- 2.Attacks against ICMPv6
- **3.Attacks against DHCPv6**



1.Attacks against the IPv6 protocol

2.Attacks against ICMPv6

3.Attacks against DHCPv6

1.1. Multicast:

Via certain multicast messages an attacker can very fast do a reconnaissance attack on a local network.

ping6 ff02::1%br0 PING ff02::1%br0(ff02::1) 56 data bytes 64 bytes from fe80::6e62:6dff:fed1:dfad: icmp_seq=1 ttl=64 time=0.052 ms 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=1 ttl=64 time=0.455 ms (DUP!) 64 bytes from fe80::5054:ff:fe90:de19: icmp_seq=1 ttl=64 time=0.650 ms (DUP!) 64 bytes from fe80::6e62:6dff:fed1:dfad: icmp_seq=2 ttl=64 time=0.046 ms 64 bytes from fe80::5054:ff:fe90:de19: icmp_seq=2 ttl=64 time=0.203 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=2 ttl=64 time=0.203 ms (DUP!) 64 bytes from fe80::5054:ff:fed1:dfad: icmp_seq=2 ttl=64 time=0.241 ms (DUP!) 64 bytes from fe80::6e62:6dff:fed1:dfad: icmp_seq=3 ttl=64 time=0.064 ms 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.237 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 64 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 65 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 65 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 65 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 65 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 65 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 66 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 66 bytes from fe80::5054:ff:fede:b69c: icmp_seq=3 ttl=64 time=0.254 ms (DUP!) 67 bytes from fe80::5054:ff:fede:b69c:b60/0000

Source: https://superuser.com/questions/840767/ipv6-multicast-address-for-all-nodes-on-network



1.Attacks against the IPv6 protocol

2. Attacks against ICMPv6

3.Attacks against DHCPv6

1.1. Multicast:

Via certain multicast messages an attacker can very fast do a reconnaissance attack on a local network.

1.2. Extension Headers

Inside extension headers an attacker can send information that remain undetected if the intermediary firewalls do not fully check the options of these headers. This kind of attack is called "covert channel".



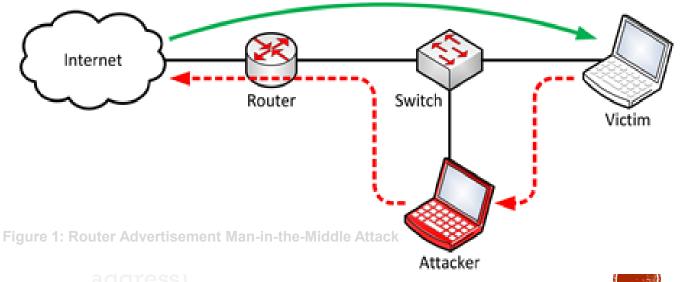
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If an attacker sends spoofed Router Advertisements inside a subnet, all IPv6 nodes will immediately change their routing tables and store the attacker as one of the default routers





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2.3. Neighbor Discovery Spoofing
When the attacker spoofs certain Neighbor
Advertisements, he can execute a MITM attack.

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A DoS attack is executed if the attacker answers to

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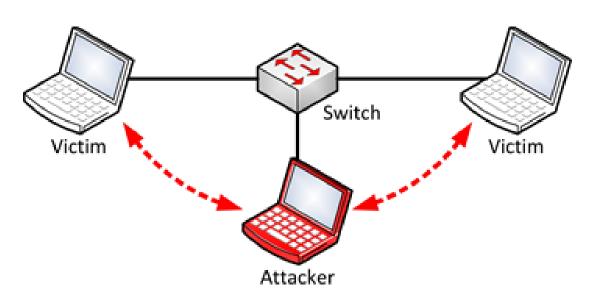


Figure 2: Neighbor Advertisement Man-in-the-Middle Attack

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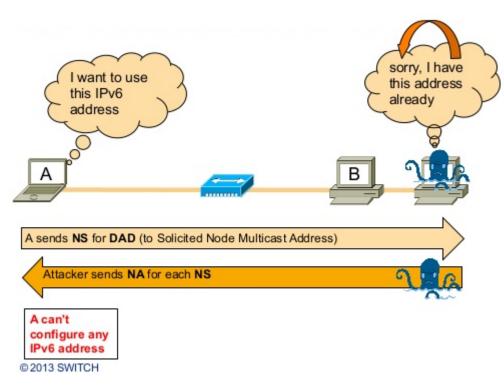
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3.1. Address Space Exhaustion

If the concept of stateful DHCPv6 is used, an attacker can exhaust the IPv6 address pool on the server, similar to a DHCPv4 server.

3.2. Rogue DHCPv6 Server

An attacker can also place his own DHCPv6 server inside a network and distribute falsified values, e.g. a spoofed DNSv6 server address.



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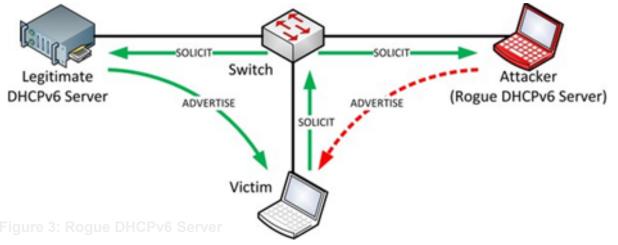
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MAC Address from IPv6

Step 1

Take your MAC address (52:74:f2:b1:a8:7f)

C:\Users\ABC>ipconfig/all

• Step 2

Add ff:fe to the middle of MAC address So, we have 52:74:f2:**ff:fe**:b1:a8:7f

Step 3

Convert 1st octet from Hexadecimal to Binary 52 equals 0101 0010

Step 4

Invert index number 6 0101 0010 becomes 0101 0000

• Step 5

Converting Binary to Hexadecimal 0101 0000 equals 50

hexadecimal	binary	
0	0000	
1	0001	
2	0010	
3	0011	
4	0100	
5	0101	
6	0110	
7	0111	
8	1000	
9	1001	
A	1010	
В	1011	
С	1100	
D	1101	
Е	1110	
F	1111	

• Step 6

New version after conversion **50**:74:f2:ff:fe:b1:a8:7f

• Step 7

Add the local prefix (fe80) fe80:: 50:74:f2:ff:fe:b1:a8:7f

• Step 8

Celebrate. You did it!



Conclusion

- IPv6 is slightly better than IPv4 in terms of security but IPv6 is not necessarily more secure than IPv4
- From a top view, the vulnerability concepts have not changed that much
- Take some time until all network and security devices have a full adoption of all IPv6 defense methods
- Some attacks are possible against the transition methods from IPv4 to IPv6
- IPv6 stack implementations in all operating systems will have errors. That means: Not only the generic IPv6 security issues but also application specific vulnerabilities that are new due to the usage of IPv6.



References

- https://www.whatismyip.com/ip-address-lookup/
- https://whatismyipaddress.com/ip-addresss
- <u>https://www.webopedia.com/DidYouKnow/Internet/ipv6_ipv4_difference.html</u>
- https://www.juniper.net/us/en/products-services/what-is/ipv4-vs-ipv6/
- https://dl.packetstormsecurity.net/papers/general/security-IPv6.pdf
- <u>https://www.first.org/resources/papers/conf2018/Herberg-</u> <u>Frank_FIRST_20180624.pdf</u>

