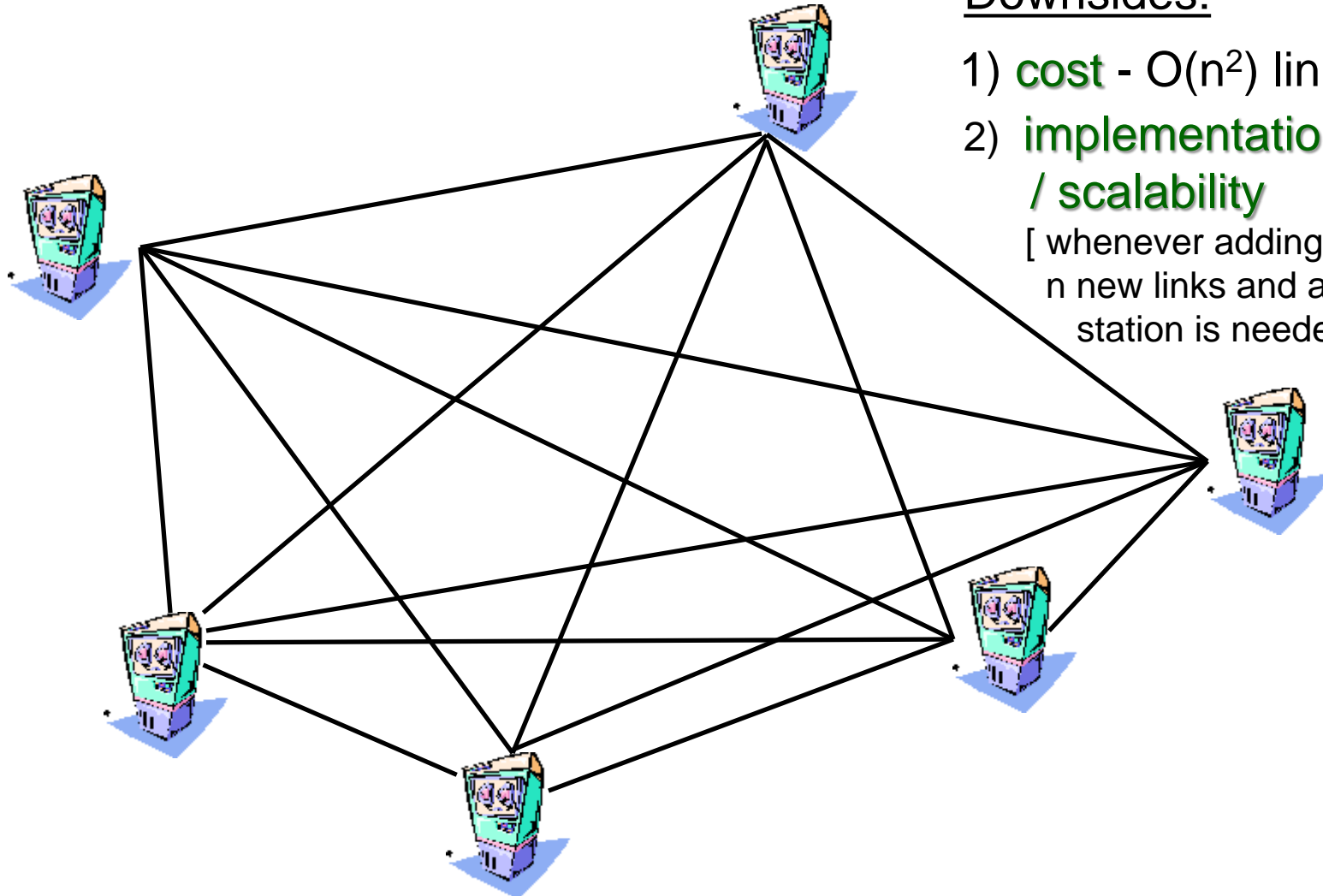


## Connecting Mainframes - Option 1: Full Mesh / Direct Link Infrastructure

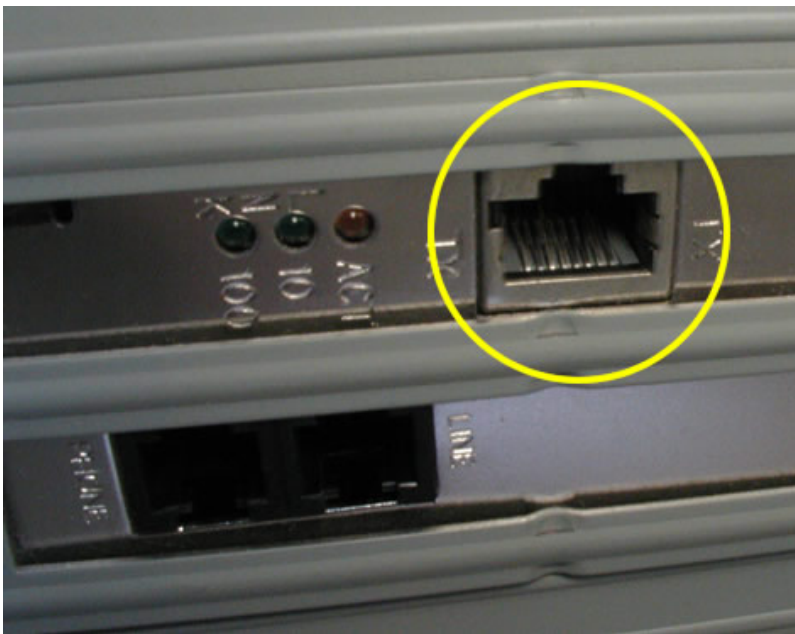


### Downsides:

- 1) **cost** -  $O(n^2)$  links and NICs
- 2) **implementation complexity / scalability**  
[ whenever adding a new station,  $n$  new links and a NICs at each station is needed ]

## NIC (Network Interface Card)

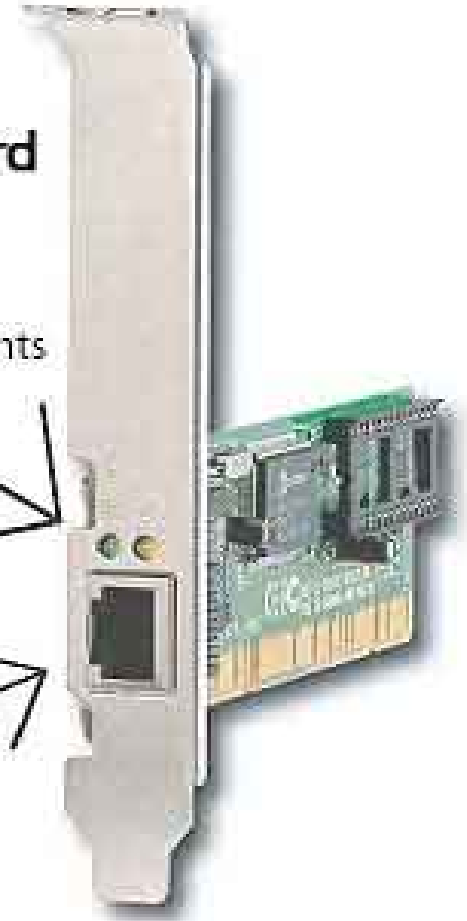
- piece of hardware that allows a computer to 'communicate' with other computers over a network



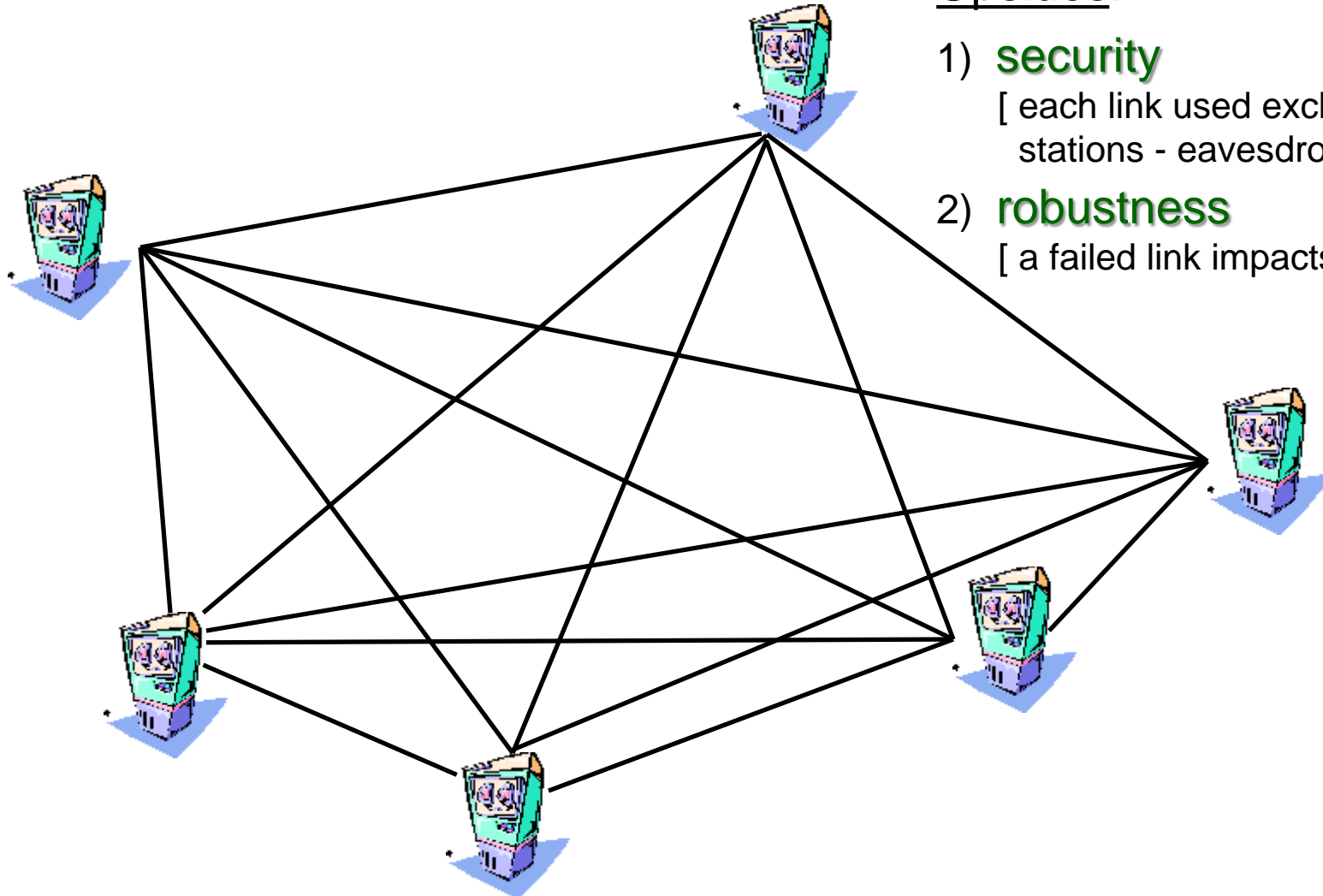
Network Card

Link & Activity Lights

RJ-45 Port



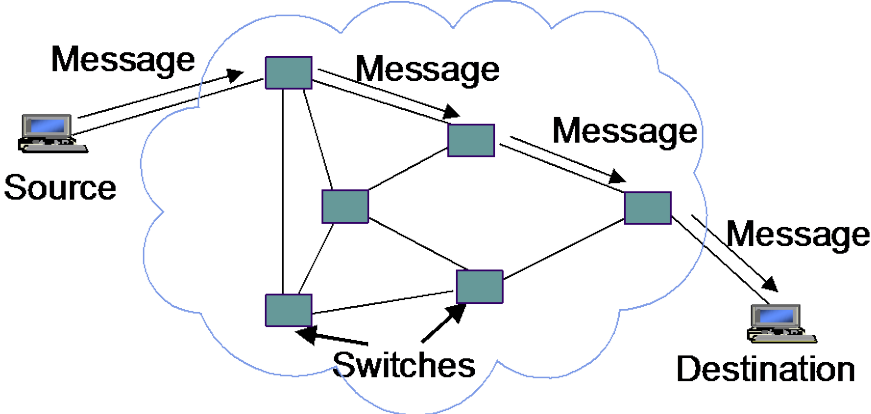
## Connecting Mainframes - Option 1: Full Mesh / Direct Link Infrastructure



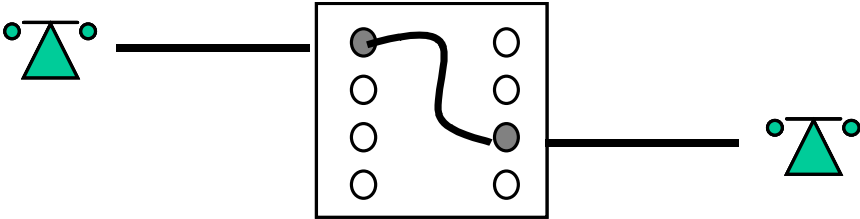
### Upsides:

- 1) **security**  
[ each link used exclusively by 2 stations - eavesdrop. impossible ]
- 2) **robustness**  
[ a failed link impacts only 2 stations ]

# Connecting Mainframes - Better Solution: Networked (Switched) Infrastructure!!!

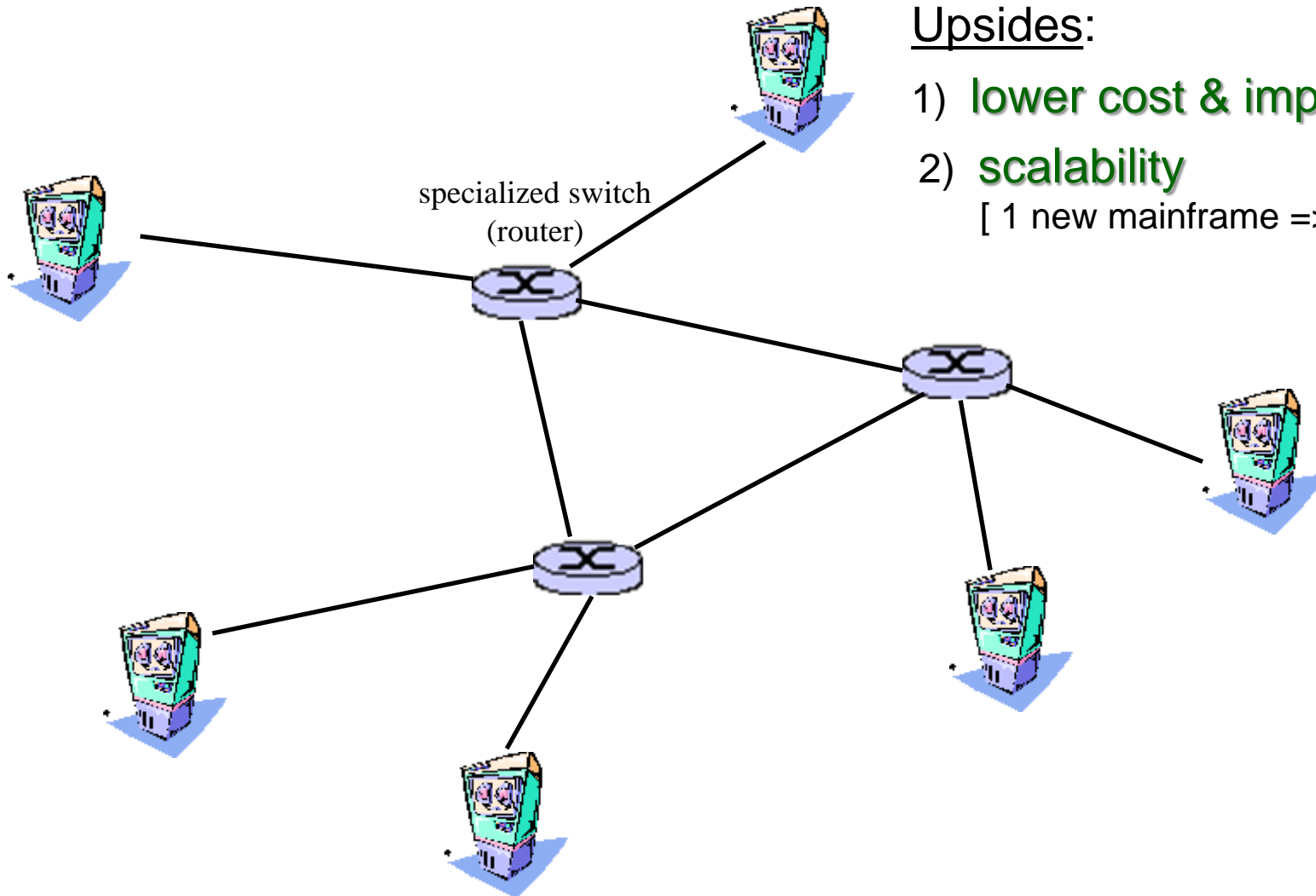


telegraph network



telephone network

## Connecting Mainframes - Option 1: Switched Network Infrastructure



### Upsides:

- 1) **lower cost & impl. complexity**
- 2) **scalability**  
[ 1 new mainframe => 1 new link ]

## Router / Packet Switch

- dedicated device/computer that forwards data packets between computer networks



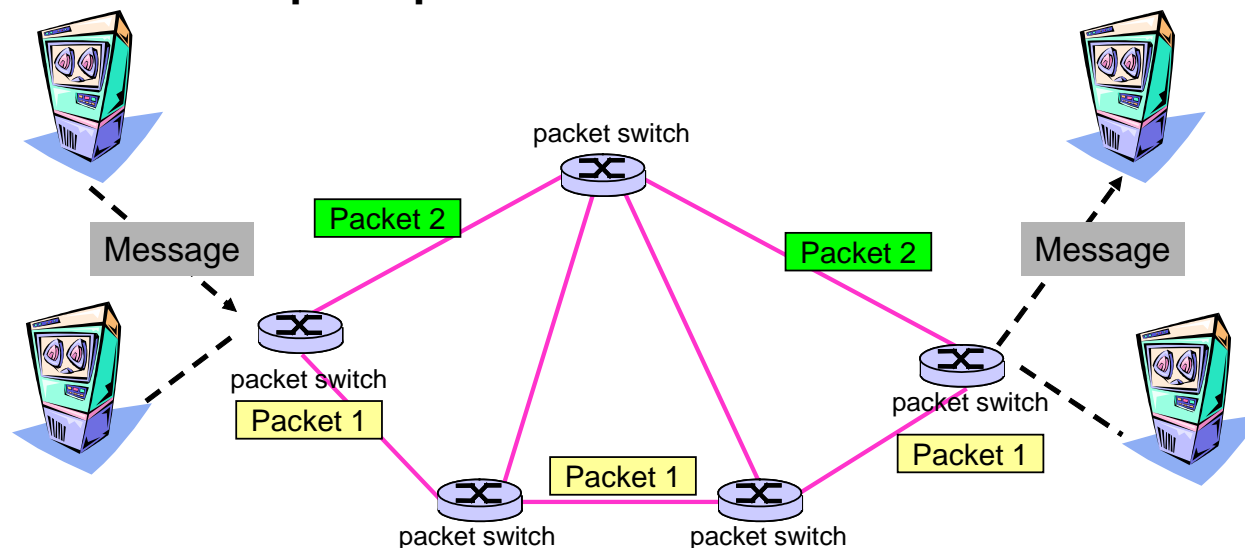
Cisco 1800



Cisco 7200

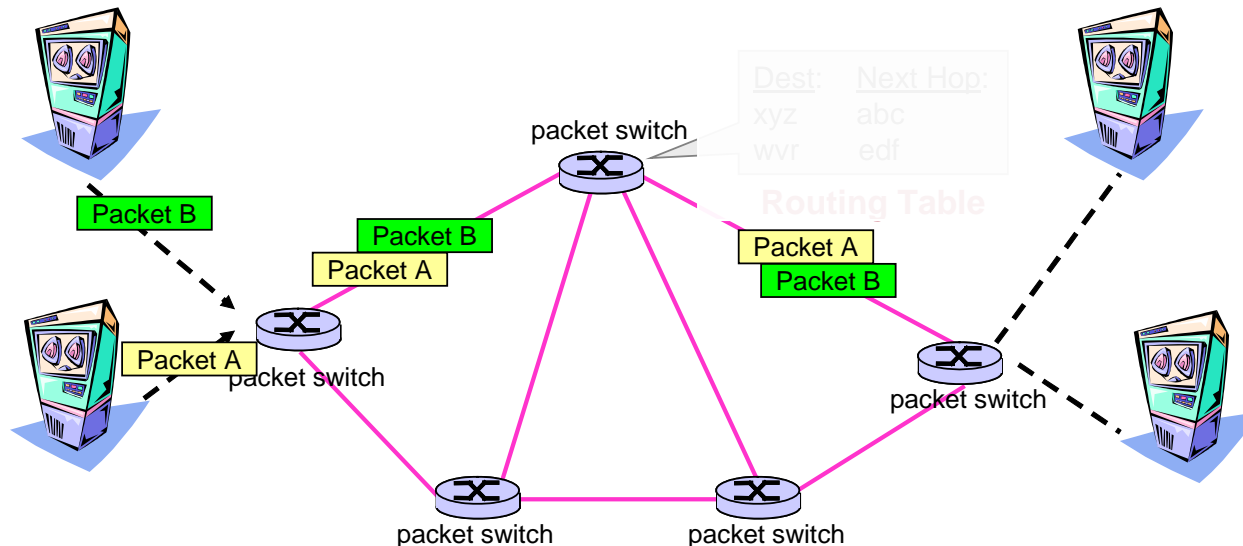
## ARPANET: Architecture

- network core consists of **packet switches** (dedicated minicomputers) to avoid costly full mesh topology
  - each packet switch connects to at least two other switches to provide alternative paths in case of failure
- network transfer messages by breaking them into **packets** of fixed size
  - long messages  $\Rightarrow$  long delays & higher prob. of error
  - each packet has a header with destination address – **packets are transmitted independently !!!**
- network transfers packets using “**store and forward**” principle



## ARPANET: Distinguishing Features

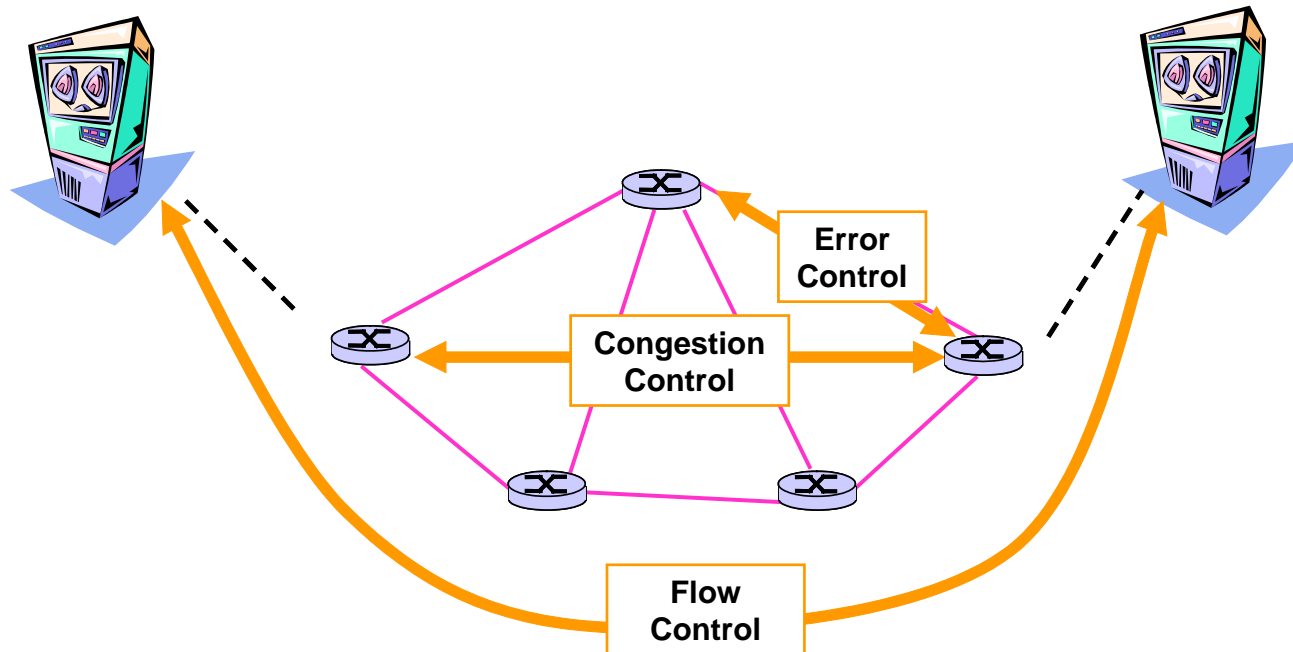
- each packet switch contains **routing / forwarding tables** ('next hop per destination' tables)
  - each packet contains destination address  $\Rightarrow$  packet switch looks at routing table and forwards packet in right direction
- **connectionless service (QoS not guaranteed)**
  - no connection setup is required prior to packet transmission
  - **packets are buffered at packet switches** to await transmission on appropriate link
  - **packets from different users are multiplexed on links** between packet switches





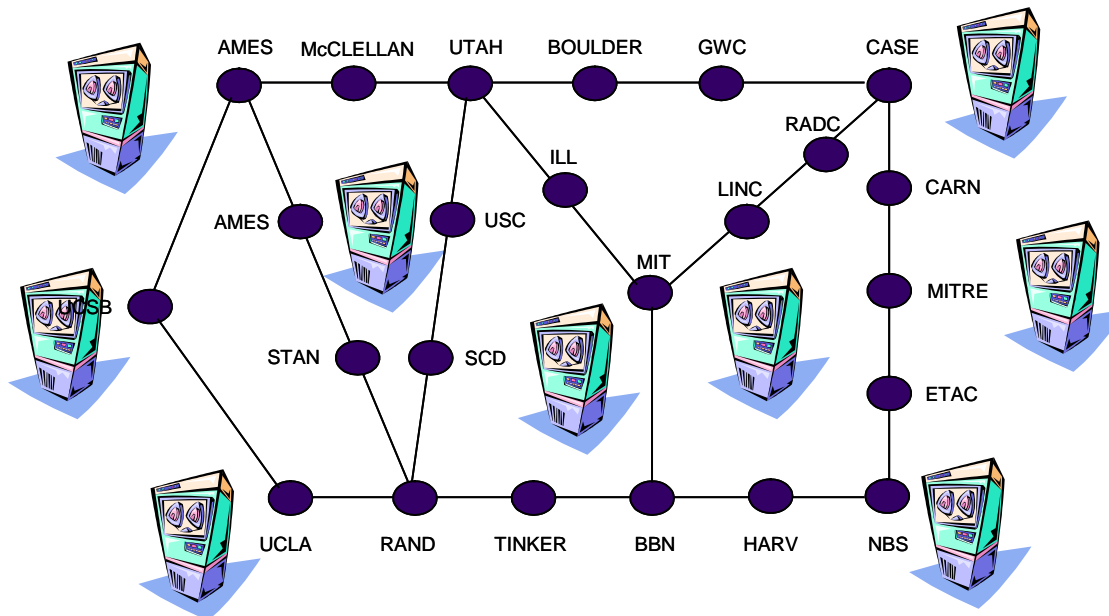
## ARPANET: Distinguishing Features

- **error control** between adjacent packet switches enables faster error recovery
  - partial responsibility of IP protocol
- **congestion control** inside the network prevents buffer overflow at core packet switches
- end-to-end **flow control** prevents buffer overflow at receiver / sender
  - responsibility of TCP protocol



## ARPANET: Applications

- “**dumb core, intelligent edges**” enabled development of many interesting and useful applications: e-mail, file transfer (FTP), remote login (Telnet)
  - **dumb core** – packet switches are only required / capable of packet forwarding
  - **intelligent edges** – end-devices have considerable CPU and memory capabilities



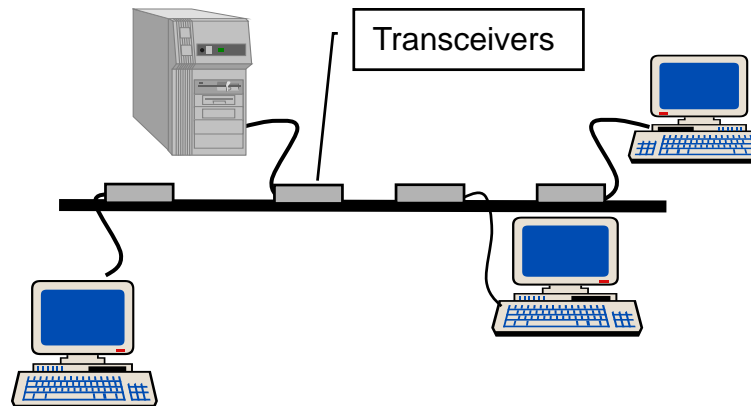
- 1950s - 1960s: Terminal-Oriented Computer Networks
- 1960s – 1970s: Computer-to-Computer Networks:  
the ARPANET – first Wide Area Network (WAN)
- 1980s:** **Local Area Networks (LANs)**
- 1980s: The Internet

# Local Area Networks

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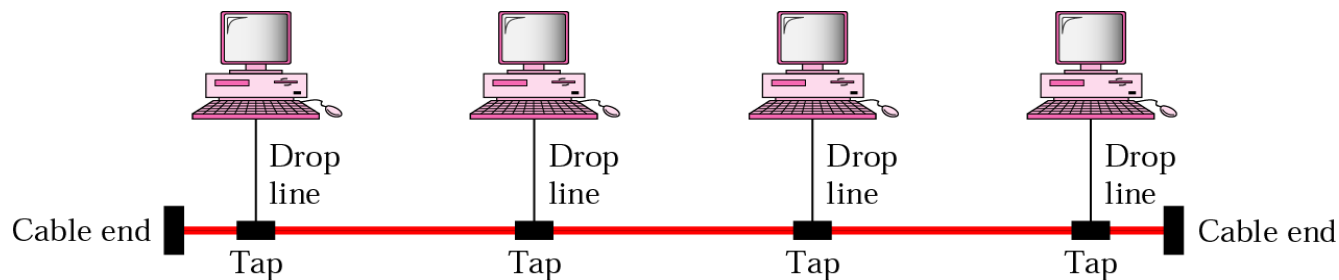
## LAN History

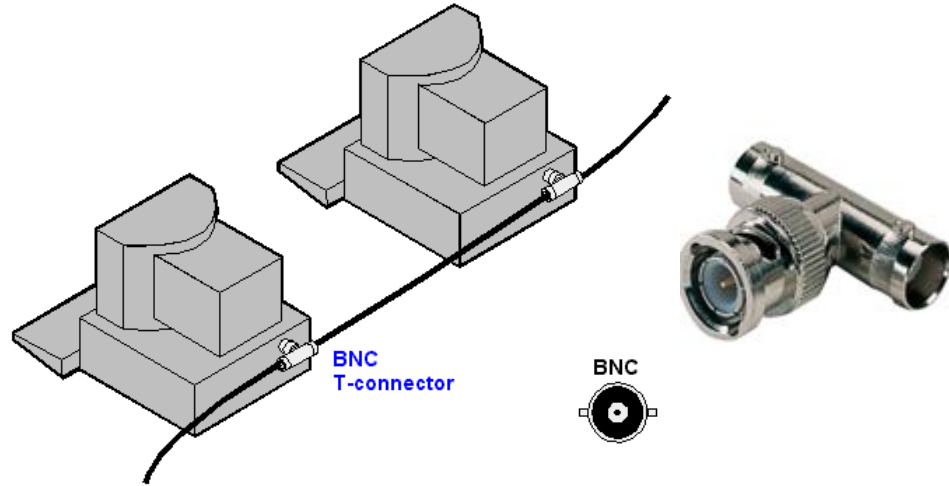
- in 1980s affordable computers become available
- subsequently, need for low-cost, high-speed, and low error-rate networks arose
  - to interconnect local workstations over small radius < 1km
  - to enable sharing of local resources (printers, servers, etc.)
- complex packet switching, congestion and flow control were unnecessary
- variety of LAN topologies emerged, including: bus, ring



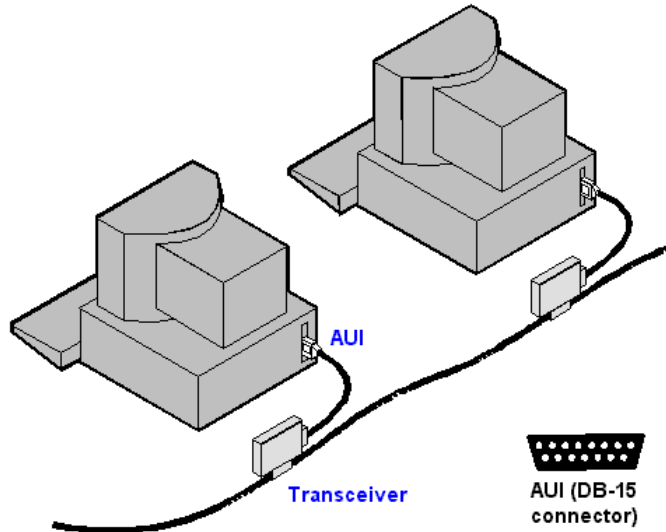
## Bus Topology (Ethernet)

- one long cable, so-called **backbone**, links all devices in the network – similar to single-line mainframe architecture
  - each workstation connects to backbone through Network Interface Card (NIC); each NIC has globally unique address
  - data frames are broadcast into coaxial cable
  - **receive**: NIC listens to medium for frames with its address
  - **send**: NIC listens to medium for presence of ongoing transmission – if no transmission is found, send frame
  - **collision**: if frame collides with somebody else's frame, abort transmission and retry later

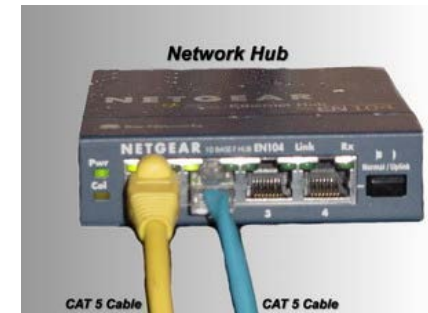
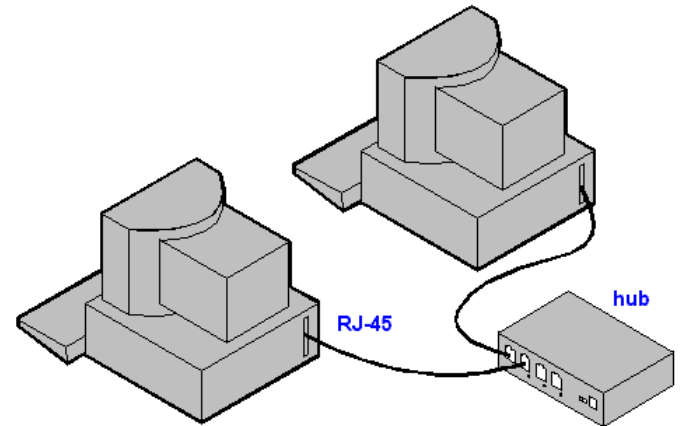




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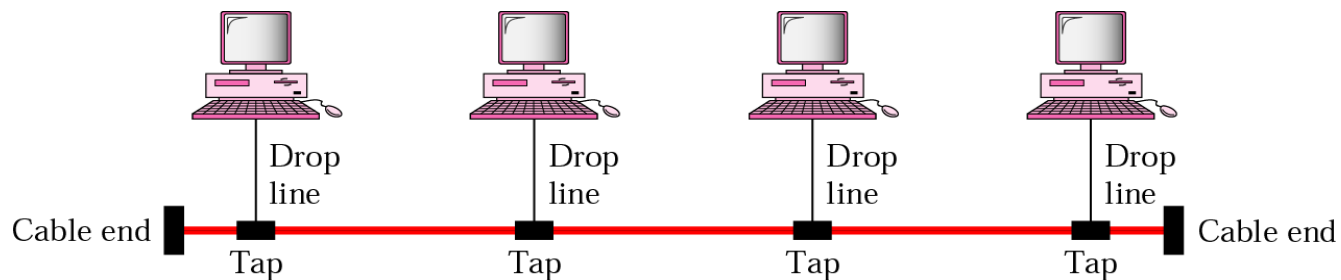


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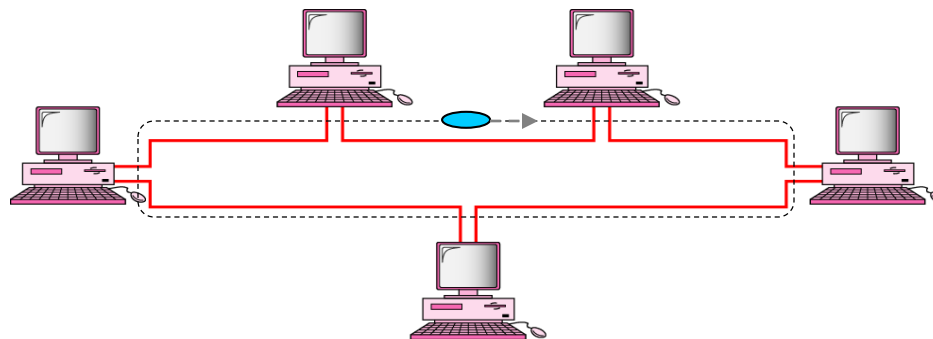


## Bus Topology (Ethernet)

- **advantages:** simple & inexpensive installation
- **disadvantages:** 1) backbone = single point of failure  
2) collisions  $\Rightarrow$  diminishing capacity
  - ◆ if two or more devices transmit simultaneously their signals will interfere



- Ring Topology** – each device has a dedicated point-to-point connection only with the two devices on either side of it
- a small frame – **token** – circulates around the ring; only the station that possesses the token is allowed to transmit at any given time
  - signal is passed along the ring in one direction, from device to device, until it reaches its destination
  - **advantages:** fairness in access / effective use of bandwidth – token-passing provides each station with a turn to transmit
  - **disadvantages:** entire network will fail if there is a failure in any transmission link or in the mechanism that relays the token





- 1950s - 1960s: Terminal-Oriented Computer Networks
- 1960s – 1970s: Computer-to-Computer Networks:  
the ARPANET – first Wide Area Network (WAN)
- 1980s: Local Area Networks (LANs)
- 1980s: The Internet**

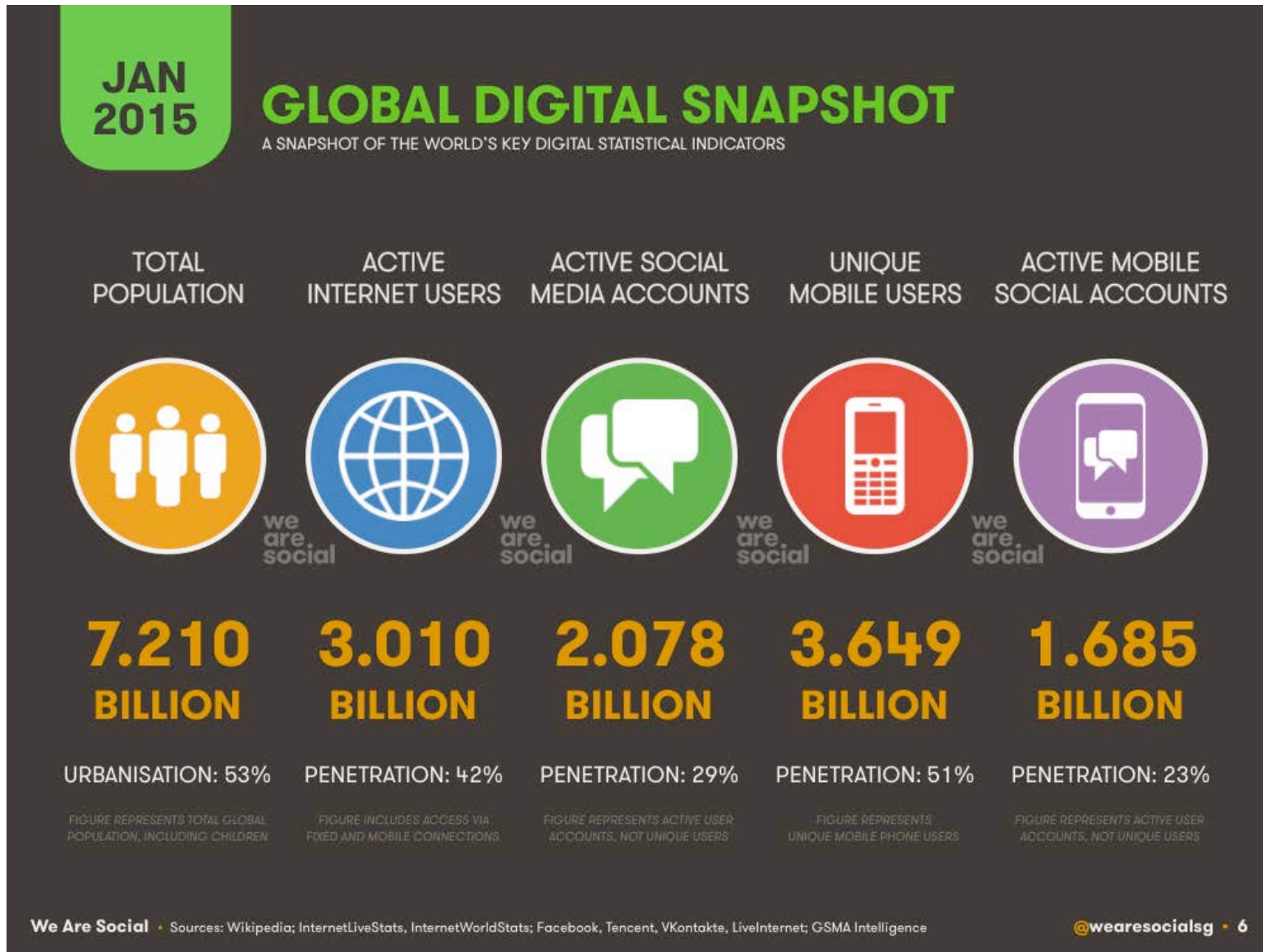
**Internet = Internetwork** – two or more interconnected networks –  
– network of networks

## The Internet: Past

- LANs that emerged in 1970s were different in terms of their underlying technology and operation
- a **protocol that would enable communication across multiple dissimilar networks** was needed
  - ◆ “higher level of abstraction” protocol – **IP Protocol**
- **Internet Protocol / Addressing** were soon developed and enabled creation of a single global internetwork

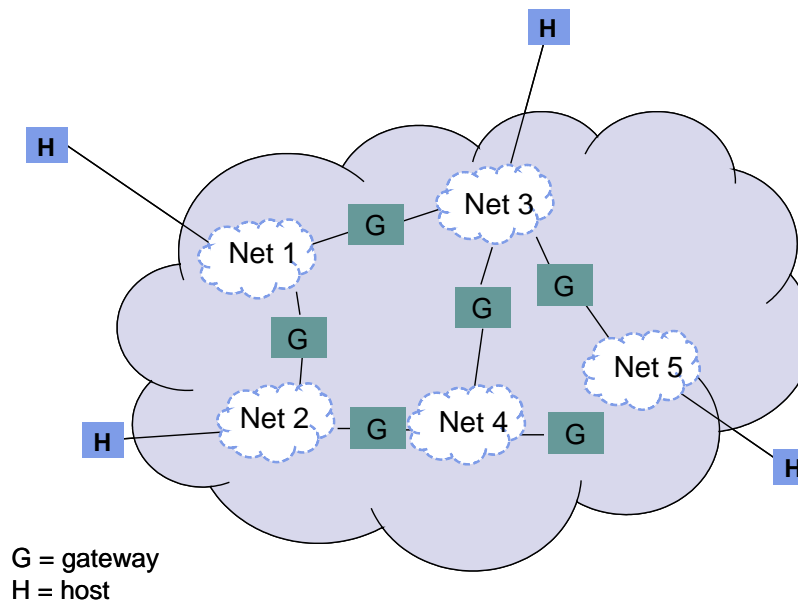
## The Internet: Present

- spread over 200 countries
- made up of 100,000s of interconnected networks, 10,000,000s of interconnected hosts, and 100,000,000s of users
- still grows exponentially ...



## The Internet = IP Network

- each component network must contain special packet switch, **gateway / router**, through which it interconnects with rest of the Internet
- host computers place data in **IP packets** (data + IP header) and deliver them to nearest router
- router, with help of other routers, attempts to forward packet across the Internet
- “**best effort service**” – IP provides no mechanism to deal with packet loss, corruption, reordering



## IP Addressing

- addressing scheme that fits (inter)network structure:  
**IP address = Net ID + Host ID**
- IP packets are routed only based on Net ID in destination IP address
  - routers have to know only major networks, not every single host  $\Rightarrow$  less memory / network update requirements
  - smaller routing tables  $\Rightarrow$  faster routing

